







# **Commodore 64 Assembler Workshop**





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An interface was used to produce this book from a microcomputer disc, which ensures direct reproduction of error-free program listings.

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# Introduction

The *Commodore 64 Assembler Workshop* is aimed at those of you who have been delving into the delights of programming at machine code level. It is a natural progression from *Commodore 64 Assembly Language*, but will be invaluable even if you learned assembler and machine code using any of the other relevant books available. It provides a bench full of useful assembly language routines and utilities programs and examines the techniques involved.

Extensive use of vectored addresses is made throughout the Commodore's operation, allowing modifications to be made to the manner in which the micro operates. Chapter 2 demonstrates how the CHRGET subroutine can be used to allow new RAM-based commands to be added to the already extensive facilities provided by the machine. A short 'wedge' interpreter is provided and the techniques for adding your own commands examined, and to get you going, three commands come supplied with the wedge interpreter: @CLS, @UP and @LOW.

Conversion between ASCII based numerical character strings and their two-byte binary equivalents and vice versa is not straightforward. Such conversions are fully described in Chapters 3 and 4, and working routines are listed.

Any program which handles strings of data must be able to manipulate the strings, whether it is an adventure game or the latest stock control reports. Routines for comparing, copying, deleting and inserting strings are included, and Chapter 6 goes on to show the various ways in which text can be printed to the screen.

Sorting data lists into order is a task which it is often necessary to perform within a program, so the technique of bubble sorting is investigated.

Many other processors provide operations that would be useful to have available when using the 6510. A software stack implementation similar to that found on the 6809 preprocessor is produced in Chapter 8, allowing up to eight selected registers to be pushed on to a memory-based stack.

Routines to move, fill and produce a hex and ASCII dump of memory are then examined and the final chapter provides a few hi-resolution graphics utilities to speed you along the way.

Many of the chapters suggest projects for you to undertake at your leisure, while every program has a detailed line-by-line description of its operation. Program listings are provided using BASIC loaders so that they can be used directly as they are. Included in each line is a REM statement giving the mnemonic representation of the instruction should you be using an assembler.

In fact, all the tools for using the *Assembler Workshop* are supplied—assuming of course you have the workbench!

Highbury, November 1984

Bruce Smith

# 1 Opening the Tool Box

The routines included in this book are designed to make your life that much easier when writing machine code. Quite often, after mastering the delights of the Commodore 64's microprocessor, programmers become frustrated because the techniques involved in, say, converting between ASCII characters and their equivalent binary values are not known. Nor are they readily available in a published form, so the painful process of sitting down armed with pencil and paper and working out the conversion through trial and error begins.

This is just one example of the type of assembler program you will find within these pages. Wherever possible, they are supplied in a form that will make them relocatable, the only addresses requiring alteration being those specified by JSR or JMP.

Each listing is in the form of a BASIC loader program, using a loop to READ and POKE decimal machine code data into memory. This will allow those of you who have not yet splashed out your hard earned cash on a suitable assembler program to get underway. For those lucky ones among you who do have an assembler, each data statement has been followed by a REM line containing the standard mnemonic representation of the instruction (see Appendix 1 for a summary). This can be entered directly and assembled as required.

Although the programs are typeset they have been spooled direct as ASCII files and loaded into my word processor so all should run as they are.

BASIC is used freely to demonstrate the machine code's operation—rather than repeating sections of assembler code, BASIC is often used to shorten the overall listing, and it is left to you to add further sections of assembler from other programs within the book or from your own resources. For example, many programs require you to input a decimal address. In the demonstrations, this is indicated by means of a one-line INPUT statement. In Chapter 3, however, there is a routine for inputting a string of five ASCII

decimal characters and converting it into a two-byte binary number. This can be inserted into the assembler text of the program, to go some way to making it a full machine code program available for use as a completely self-contained section of machine code.

## WRITING MACHINE CODE

You have an idea that you wish to convert into machine code—so what's the best way to go about it? Firstly, make some brief notes about its operation. Will it use the screen? If so, what mode? Will it require the user to input values from the keyboard? If so, what keys do you use? What will the screen presentation look like? Will you want to use sound?...and so on. Once you have decided on the effects you want, put them down in flowchart form. This need not be the normal flowchart convention of boxes and diamonds—I find it just as easy to write each operation I want the program to perform in a list and then join the flow of these up afterwards.

Quite often, the next step is to write the program in BASIC! This may sound crazy, but it allows you to examine various aspects of the program's operation in more detail. An obvious example of this is obtaining the correct screen layout—you might find after running the routine that the layout does not look particularly good. Finding this out at an early stage will save you a lot of time later, avoiding the need to rewrite the screen layout portion of your machine code—rewriting BASIC is much easier! If you write the BASIC tester as a series of subroutines, it will greatly simplify the process of conversion to machine code. Consider the main loop of such a BASIC tester, which takes the form:

```
10 GOSUB 200 : REM SET UP VARIABLES
20 GOSUB 300 : REM SET UP SCREEN
30          REM LOOP
40 GOSUB 400 : REM INPUT VALUES
50 GOSUB 500 : REM CONVERT AS NEEDED
60 GOSUB 600 : REM DISPLAY VALUES
70 GOSUB 700 : REM DO UPDATE
80 IF TEST=NOTDONE THEN GOTO 30
90 END
```

Each module can be taken in turn, converted into assembler and tested. Once performing correctly the next procedure can be examined. Debugging is made easier because the results of each module are known having used the BASIC tester. The final main loop of the assembler might then look something like this:

```

JSR $C200 : REM SET UP VARIABLES
JSR $C300 : REM SET UP SCREEN
           REM LOOP
JSR $C400 : REM INPUT VALUES
JSR $C500 : REM CONVERT AS NEEDED
JSR $C600 : REM DISPLAY VALUES
JSR $7700 : REM DO UPDATE
BNE LOOP

```

You might be surprised to learn that this technique of testing machine code programs by first using BASIC is employed by many software houses the world over.

## DEBUGGING

A word or two about debugging machine code programs that will not perform as you had hoped: if this happens to you, before pulling your hair out and throwing the latest copy of *Machine Code Nuclear Astrophysics Weekly* in the rubbish bin, a check of the following points may reveal the bug!

1. If you are using a commercial assembler, check that your labels have all been declared and correctly assigned. If you are assembling 'by hand', double-check all your branch displacements and JMP and JSR destination addresses. You can normally ascertain exactly where the problem is by examining how much of the program works before the error occurs, rather than checking it all.
2. If your program uses immediate addressing, ensure you have prefixed the mnemonic with a hash (#) to inform the assembler or, if compiling by hand, check that you have used the correct opcode. It is all too easy to assemble the coding for LDA \$41 when you really want the coding for LDA #\$41.
3. Check that you have set or cleared the Carry flag before subtraction or addition.
4. My favourite now—ensure that you save the result of a subtraction or an addition. The sequence:

```

CLC
LDA $FB
ADC #1
BCC OVER
INC $FC
OVER
RTS

```

is not much good if you don't save the result of the addition with:

```
STA $FB
```

*before* the RTS!

5. Does the screen clear to the READY prompt whenever you perform a SYS call, seemingly without executing any of the machine code? The bug that often causes this is due to an extra comma being inserted into a series of DATA statements. For example the DATA line:

```
DATA 169,Ø, , 162, 255
```

with an extra comma between the Ø and 162, would assemble the following:

```
LDA #$$ØØ  
BRK  
LDX #$$FF
```

as the machine has interpreted ‘.,’ as ‘,Ø,’ and assembled the command which has zero as its opcode—BRK!

6. Does the program ‘hang up’ every time you run it, when you are quite certain that the data statements are correct? This is often caused by a full stop instead of a comma being used between DATA statements, e.g.

```
DATA 169, 6, 162.5, 96
```

Here, if a full stop has been used instead of a comma between the 162 and the 5, the READ command interprets this as a single number, 162.5, rounds it down to 162, and assembles this ignoring the 5 and using the 96 (RTS) as the operand, as follows:

```
LDA #$$Ø6  
LDX #$$6Ø  
XXX
```

When executed, the garbage after the last executable instruction results in the system hanging up. This error should not occur if you calculate your loop count correctly, so always double-check this value before running your program.

If none of these errors is the cause of the problem, then I'm afraid you must put your thinking cap on. Well-commented assembler will make debugging very much easier.



## 2 Commodore Command

One of the disadvantages of using random access memory-based machine code routines as utilities within a BASIC program is that it is left to you, the programmer, to remember just where they are stored, and to use the appropriate SYS call to implement them. This doesn't usually pose any problems if only one or two machine code utilities are present; the problems occur when several are being used. Normally you would need to keep a written list of these next to you, looking up the address of each routine as you need it. Great care must be taken to ensure that the SYS call is made to the correct address, as a mis-typed or wrongly called address can send the machine into an infinite internal loop, for which the only cure is a hard reset, which would destroy all your hard work.

The program offered here provides a useful and exciting solution to the problem, enabling you to add new commands to your Commodore 64's vocabulary. Each of your routines can be given a command name, and the machine code comprising any command will be executed by simply entering its command name. The routine is written so that these new commands can be used either directly from the keyboard or from within programs.

The trick in 'teaching' the Commodore 64 new commands is to get the machine to recognize them. If an unrecognized command is entered at the keyboard, the almost immediate response from the 64 is '?SYNTAX ERROR'. If you have any expansion cartridges you'll know that it is possible to expand the command set, and the *Programmer's Reference Guide* gives a few hints on how to do this, on pages 307 and 308—the method pursued here is by resetting the system CHRGET subroutine.

### CHRGET

The CHRGET routine is, in fact, a subroutine which is called by the main BASIC Interpreter. You can think of it as a loop of code, protruding from the machine, into which we can wedge our own

bits of code, thereby allowing fundamental changes to be made to the manner in which the Commodore operates. Let's have a look at how the normal CHRGET subroutine (which is located in zero page from \$73) operates:

**Table 2.1**

Address	Machine code	Assembler
\$0073	E6 7A	INC \$7A
\$0075	D0 02	BNE \$0079
\$0077	E6 7B	INC \$7B
\$0079	AD xx xx	LDA \$xxxx
\$007C	C9 3A	CMP #\$3A
\$007E	B0 0A	BCS \$008A
\$0080	C9 20	CMP #\$20
\$0082	F0 EF	BEQ \$0073
\$0084	38	SEC
\$0085	E9 30	SBC #\$30
\$0087	38	SEC
\$0088	E9 D0	SBC #\$D0
\$008A	60	RTS

The subroutine begins by incrementing the byte located at \$7A. This address forms a vector which holds the address of the interpreter within the BASIC program that is currently being run. If there is no carry over into the high byte, which must therefore itself be incremented, a branch occurs to location \$0079. You will notice that the bytes which have just been incremented lie within the subroutine itself. These are signified in the above listing by 'xx xx', because they are being updated continually by the routine. The reason for this should be fairly self-evident: looking at the opcode, we can see that it is LDA, therefore each byte is, in turn, being extracted from the program.

The next two bytes at \$007C perform a compare, CMP #\$3A. The operand here, \$3A, is the ASCII code for a colon, so CHRGET is checking for a command delimiter. The BCS \$008A will occur if the accumulator contents are greater than \$3A, effectively returning control back to within the BASIC Interpreter ROM. The next line, CMP #\$20, checks whether a space has been encountered within the program. If it has, the branch is executed back to \$0073 and the code rerun.

The rest of the coding is checking that the byte is a legitimate

one—it should be an ASCII character code in the range \$30 to \$39, that is, a numeric code. If it is, the coding will return to the main interpreter with the Carry flag clear. If the accumulator contains less than \$30 (it could, of course, have ASCII \$20 in it, as we have already checked for this) then the Carry flag is set.

It is important to understand what is happening here, as we will need to overwrite part of this code to point it in the direction of our own ‘wedge’ interpreter. This has to perform the ‘deleted’ tasks before returning to the main interpreter to ensure the smooth and correct running of the Commodore 64.

## THE WEDGE OPERATING SYSTEM

To distinguish the Wedge Operating System (WOS) commands from normal commands (and illegal ones!), we must prefix them with a special character—one which is not used by the Commodore 64. The *Programmer’s Reference Guide* suggests the use of the ‘@’ sign, so that’s what we will use.

Program 1a lists the coding for the WOS. I have chosen to place it well out of the way, in the free RAM area from 49666 (3C202) onwards. As we shall see the memory below (bis to 49152 (\$C000)) is also used by the WOS.

### Program 1a

```
10 REM *** WEDGE OPERATING SYSTEM - WOS ***
20 REM *** WOS INTERPRETER FOR COMMODORE 64 ***
30 :
40 CODE=49666
50 FOR LOOP=0 TO 188
60 READ BYTE
70 POKE CODE+LOOP,BYTE
80 NEXT LOOP
90 :
100 REM ** M/C DATA **
110 DATA 169,0 : REM LDA #00
120 DATA 160,192 : REM LDY #C0
130 DATA 32,30,171 : REM JSR $AB1E
140 DATA 169,76 : REM LDA #$4C
150 DATA 133,124 : REM STA $7C
160 DATA 169,24 : REM LDA #$18
170 DATA 133,125 : REM STA $7D
180 DATA 169,194 : REM LDA #$C2
190 DATA 133,126 : REM STA $7E
```

```

200 DATA 108,2,3 : REM JMP ($0302)
205 :: REM WOS STARTS HERE
210 DATA 201,64 : REM CMP #40
220 DATA 208,68 : REM BNE $44
230 DATA 165,157 : REM LDA $9D
240 DATA 240,40 : REM BEQ $28
250 DATA 173,0,2 : REM LDA $0200
260 DATA 201,64 : REM CMP #40
270 DATA 208,28 : REM BNE $1C
280 DATA 32,114,194 : REM JSR $C272
290 DATA 160,0 : REM LDY #00
300 DATA 177,122 : REM LDA ($7A),Y
310 DATA 201,32 : REM CMP #20
320 DATA 240,9 : REM BEQ $09
330 DATA 230,122 : REM INC $7A
340 DATA 208,246 : REM BNE $F6
350 DATA 230,123 : REM INC $7B
360 DATA 56 : REM SEC
370 DATA 176,241 : REM BCS $F1
380 DATA 32,116,164 : REM JSR $A474
390 DATA 169,0 : REM LDA #00
400 DATA 56 : REM SEC
410 DATA 176,29 : REM BCS $1D
420 DATA 169,64 : REM LDA #40
430 DATA 56 : REM SEC
440 DATA 176,24 : REM BCS $18
445 :: REM PROGRAM-MODE
450 DATA 32,114,194 : REM JSR $C272
460 DATA 160,0 : REM LDY #00
470 DATA 177,122 : REM LDA ($7A),Y
480 DATA 201,0 : REM CMP #00
490 DATA 240,13 : REM BEQ $0D
500 DATA 201,58 : REM CMP #3A
510 DATA 240,9 : REM BEQ $09
520 DATA 230,122 : REM INC $7A
530 DATA 208,242 : REM BNE $F2
540 DATA 230,123 : REM INC $7B
550 DATA 56 : REM SEC
560 DATA 176,237 : REM BCS $ED

```

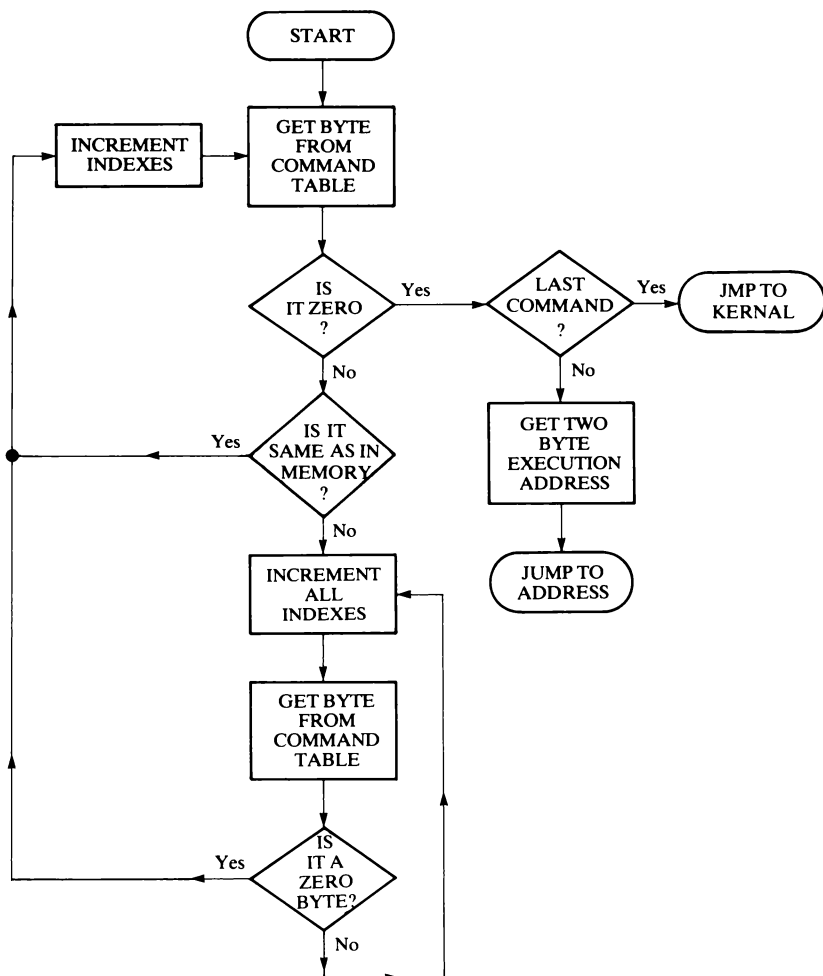


Figure 2.1 The wedge operating system flowchart

```

570 DATA 201,58 : REM CMP # $3A
580 DATA 176,10 : REM BCS $0A
590 DATA 201,32 : REM CMP # $20
600 DATA 240,7 : REM BEQ $07
610 DATA 56 : REM SEC
620 DATA 233,48 : REM SBC # $30
630 DATA 56 : REM SEC
640 DATA 233,208 : REM SBC # $D0
650 DATA 96 : REM RTS
660 DATA 76,115,0 : REM JMP $0073
665 :: REM FIND-EXECUTE
  
```

```

670 DATA 169,0 : REM LDA #000
680 DATA 133,127 : REM STA $7F
690 DATA 169,193 : REM LDA #$C1
700 DATA 133,128 : REM STA $80
710 DATA 230,122 : REM INC $7A
720 DATA 208,2 : REM BNE $02
730 DATA 230,133 : REM INC $7B
740 DATA 160,0 : REM LDY #000
750 DATA 162,0 : REM LDX #000
760 DATA 177,127 : REM LDA ($7F),Y
770 DATA 240,36 : REM BEQ $24
780 DATA 209,122 : REM CMP ($7A),Y
790 DATA 208,4 : REM BNE $04
800 DATA 200 : REM INY
810 DATA 56 : REM SEC
820 DATA 176,244 : REM BCS $F4
830 DATA 177,127 : REM LDA ($7F),Y
840 DATA 240,4 : REM BEQ $04
850 DATA 200 : REM INY
860 DATA 56 : REM SEC
870 DATA 176,248 : REM BCS $F8
880 DATA 200 : REM INY
890 DATA 152 : REM TYA
900 DATA 24 : REM CLC
910 DATA 101,127 : REM ADC $7F
920 DATA 133,127 : REM STA $7F
930 DATA 169,0 : REM LDA #000
940 DATA 101,128 : REM ADC $80
950 DATA 133,128 : REM STA $80
960 DATA 160,0 : REM LDY #000
970 DATA 232 : REM INX
980 DATA 232 : REM INX
990 DATA 56 : REM SEC
1000 DATA 176,216 : REM BCS $D8
1010 DATA 189,80,192 : REM LDA $C050,X
1020 DATA 133,128 : REM STA $80
1030 DATA 232 : REM INX
1040 DATA 189,80,192 : REM LDA $C050,X
1050 DATA 133,129 : REM STA $81
1060 DATA 108,128,0 : REM JMP ($0080)

```

```

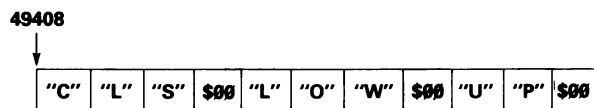
1065 ::                                REM ILLEGAL
1070 DATA 162,11                       : REM LDX #0B
1080 DATA 108,0,3                       : REM JMP ($300)
1090 :
1100 REM ** SET UP COMMAND TABLE **
1110 TABLE=49408
1120 FOR LOOP=0 TO 10
1130 READ BYTE
1140 POKE TABLE+LOOP,BYTE
1150 NEXT LOOP
1160 :
1170 REM ** ASCII COMMAND DATA **
1180 DATA 67,76,83,0                   : REM CLS
1190 DATA 76,79,87,0                   : REM LOW
1200 DATA 85,80,0                       : REM UP

```

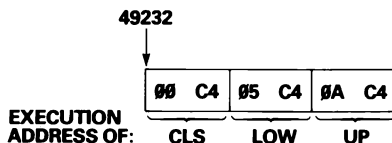
To enable the WOS to identify a wedge command, it needs a complete list to which it can compare the one it is interpreting in the program—this is done with the aid of a command table, which is formed by the program lines from 1100 to 1200. This ASCII table is based at 49408 (\$C100) and, as you can see from the listing, three commands are provided: @CLS, @LOW and @UP. Note that the @ is omitted from the front of each command in the table—it is unnecessary at the comparison stage, as by this time it has already been established that it is a WOS command—and that each command is terminated by a zero. A table listing the execution address of each command must also be constructed, but more of this later.

The main program consists of two parts, an initialization routine and the interpreter proper.

The initialization routine is embodied in lines 110 to 200. Its function is to reset the CHRGET subroutine investigated earlier. Lines 110 to 130 issue a heading on the screen indicating that the



The Command Table



The Address Table

*Figure 2.2 The Command and Address Tables.*

WOS has been initialized. The subroutine at \$AB1E, called by line 130, prints out an ASCII string located at the address given by the index registers. In this instance it is located at \$C000 (49152), and is assembled into memory by the second part of the listing. Lines 140 to 190 poke three bytes into the CHRGET subroutine which effectively assembles the code:

```
JMP $C218
```

The address \$C218 is the address of the start of the WOS interpreter at line 210. Finally, line 200 does an indirect jump through the IMAN vector at \$0302 to perform a warm BASIC start. The CHRGET subroutine, complete with wedge jump, now looks like this:

**Table 2.2**

Address	Machine code	Assembler
\$0073	E6 7A	INC \$7A
\$0075	D0 02	BNE \$0079
\$0077	E6 7B	INC \$7B
\$0079	AD xx xx	LDA \$xxxx
\$007C	4C 18 C2	JMP \$C218

When the WOS is entered, the byte in the accumulator is checked to see if it is an @ (line 210), signifying a wedge command. If it is not, then a branch to line 570 is performed. As you can see, the code from line 570 to 650 performs the normal function of the CHRGET routine, with control returning to the BASIC Interpreter.

If the byte is an @, the interpretation continues. The byte at \$9D is located, to detect whether the command is within a program or has been issued in direct mode. A zero indicates that the command has been called from within a program and the branch of line 240 to line 450 is performed. In both instances the interpretation follows similar lines—for descriptive purposes, we will assume program mode and resume the commentary from line 450.

The subroutine at \$C272 is the interpreter proper. Starting at line 665 it locates the command and executes it. The first eight bytes (lines 670 to 700) set up a zero page vector to point to the command table at \$C100. Lines 710 to 730 update the zero page bytes at \$7A and \$7B, which hold the address of the current point within the program. After initializing both index registers, the first



byte within the command table is located (lines 740 to 760), and compared to the byte within the program, immediately after the @ (line 780). If the comparison fails, the branch to line 830 is performed, locating the zero and therefore the next command in the command table. When a comparison is successful (the command is identified) and the terminating zero located by line 770, the branch to line 1010 is performed. Lines 1010 to 1060 locate the execution address of the command from the address table located at \$C050. The X register is used as an offset into this, being incremented by two each time a command table comparison fails (lines 970 and 980). The two address bytes are loaded to form a zero page vector and the machine code is executed via an indirect jump.

On completion of the routine, its terminating RTS returns control to line 460, and the next byte after the command is sought out. When a zero is found, the branch of line 490 is performed and the CHRGET routine is completed, control being returned to the BASIC Interpreter.

## THE NEW COMMANDS

Program 1b provides the assembly routines to construct the initialization prompts, the machine code for the new commands and the address table:

### Program 1b

```

1210 REM ** TITLE MESSAGE DISPLAYED ON SYS
      49666 **
1220 HEAD=49152
1230 FOR LOOP=0 TO 40
1240 READ BYTE
1250 POKE HEAD+LOOP, BYTE
1260 NEXT LOOP
1270 :
1280 REM ** ASCII CHARACTER DATA **
1290 DATA 147, 13, 32, 32, 42, 42, 32, 67, 54, 52, 32
1300 DATA 69, 88, 84, 69, 78, 68, 69, 68, 32, 83, 85
1200 DATA 80, 69, 82, 32, 66, 65, 83, 73, 67, 32, 86, 49
1310 DATA 46, 48, 32, 42, 42, 13, 0
1320 ::
1360 REM ** SET UP M/C FOR COMMANDS **
1370 MC=50176
1380 FOR LOOP=0 TO 14
1390 READ BYTE

```

```

1400 POKE MC+LOOP, BYTE
1410 NEXT LOOP
1420 :
1430 REM ** COMMAND M/C **
1440 :: REM CLS
1450 DATA 169, 147 : REM LDA #93
1460 DATA 76, 210, 255 : REM JMP $FFD2
1470 :: REM LOW
1480 DATA 169, 14 : REM LDA #0E
1490 DATA 76, 210, 255 : REM JMF $FFD2
1500 :: REM UP
1510 DATA 169, 142 : REM LDA #8D
1520 DATA 76, 210, 255 : REM JMP $FFD2
1530 ::
1540 REM ** SET UP ADDRESS TABLE **
1550 ADDR=49232
1560 FOR LOOP=0 TO 5
1570 READ BYTE
1580 POKE ADDR+LOOP, BYTE
1590 NEXT LOOP
1600 :
1610 REM ** ADDRESS DATA **
1620 DATA 0, 196 : REM CLS $C400
1630 DATA 5, 196 : REM LOW $C405
1640 DATA 10, 196 : REM UP $C40A

```

Each command's machine code is located from 50176 (\$C400). The three new commands and their functions are:

```

CLS : clear screen and home cursor
LOW : select lower case character set
UP  : select upper case character set

```

Nothing to set the house alight, admittedly, but the techniques involved are more important at present. These are simple to implement and, once understood, enable more useful and complex commands to be added. The code associated with each command is responsible simply for printing its ASCII code. The final section of listing (lines 1540 to 1650) pokes the execution address of each command into memory. The final address points to the code at line 170, and the program jumps to this position if the command is not found within the command table. This code performs an indirect jump to the BASIC Interpreter's error handler.

## USING THE WOS

Using the Wedge Operating System is easy: enter the program as shown, run it to assemble the code into memory, and if all goes well, save the program. To initialize the WOS enter:

```
SYS 49666
```

The screen will clear, and the following message be printed across the top of the screen:

```
** C64 EXTENDED SUPER BASIC V1.0 **
```

The wedge commands are now available for immediate use. Remember that pressing RUN/STOP and RESTORE together will reset the CHRGET routine to its default value making the WOS invisible. To relink it, simply execute the SYS 49666 call again.

### Line-by-line

A line-by-line description of the WOS now follows, to enable you to examine its operation in more detail:

```
line 110 : load accumulator with low byte message address
line 120 : load accumulator with high byte message address
line 130 : print start up message
line 140 : reset CHRGET subroutine
line 200 : do a BASIC warm start
line 205 : main entry for WOS
line 210 : is it an '@' and therefore a WOS command?
line 220 : no, so branch to line 570 to update
line 230 : yes, check for direct or program mode
line 240 : if zero, then WOS command is within program,
           so branch to line 450
line 250 : else direct mode so get byte from buffer
line 260 : recheck that it is a WOS command
line 270 : if error, branch to line 410
line 280 : find and execute the command else issue
           appropriate error message
line 290 : initialize index
line 300 : get byte from buffer
line 310 : is it a space?
line 320 : yes, so branch to line 380
```

line 330 : increment low byte of address  
line 340 : branch back to line 300 if high byte does not need  
to be updated  
line 350 : else increment high byte of address  
line 360 : set Carry flag and do a forced branch back to  
line 300  
line 380 : print 'READY' prompt  
line 390 : clear accumulator  
line 400 : set Carry flag and force a branch to line 500 to  
update and return  
line 420 : get '@' into accumulator  
line 430 : set Carry flag and force a branch to line 570  
line 445 : entry point for PROGRAM-MODE  
line 450 : locate and execute command or print appropriate  
error message  
line 460 : clear indexing register  
line 470 : get byte from program  
line 480 : is it a 0 and therefore end of line?  
line 490 : yes, branch to line 500  
line 500 : no, is it the command delimiter ':'?  
line 510 : yes, branch to line 570  
line 520 : no, increment low byte of address  
line 530 : if not zero, branch back to line 470 to redo loop  
line 540 : increment high byte of address  
line 550 : set Carry flag and force a branch back to line 470.  
line 570 : is it a command delimiter ':'?  
line 580 : if greater than or equal to ':' then branch to line 650  
line 590 : is it a space?  
line 600 : yes, so branch to line 650  
line 610 : set Carry flag  
line 620 : subtract ASCII base code  
line 630 : set Carry flag  
line 640 : subtract token and ASCII set bits  
line 650 : return to BASIC Interpreter  
line 660 : jump to CHRGET  
line 665 : entry for FIND-EXECUTE subroutine  
line 670 : seed address of command table (\$C100) into vector  
at \$7F  
line 710 : increment low byte of command address  
line 720 : branch over if no carry into high byte

line 730 : else increment high byte of address  
 line 740 : back together, initialize Y register  
 line 750 : and X register  
 line 760 : get byte from the command table  
 line 770 : if zero byte, then command is identified, branch to  
           line 1010  
 line 780 : is it the same as the byte pointed to in the command  
           table?  
 line 790 : no, branch to line 830  
 line 800 : increment index  
 line 820 : set Carry flag and force a branch back to line 760  
 line 830 : command not identified—seek out zero byte. Get  
           byte from command table  
 line 840 : if zero, branch to line 880  
 line 850 : increment index  
 line 860 : set Carry flag and force a branch to line 830  
 line 880 : increment index  
 line 890 : transfer into accumulator  
 line 900 : clear Carry flag  
 line 910 : add to low byte of vector address  
 line 920 : save result  
 line 930 : clear accumulator  
 line 940 : add carry to high byte of vectored address  
 line 950 : and save the result  
 line 960 : initialize index  
 line 970 : add two to X to move onto next address in the  
 line 980 : command address table  
 line 990 : set Carry flag and force a branch to line 760  
 line 1010 : get low byte of command execution address  
 line 1020 : save it in a vector  
 line 1030 : increment index  
 line 1040 : get high byte of command execution address  
 line 1050 : save it in vector  
 line 1060 : jump to vectored address to execute machine code  
           of identified command  
 line 1065 : entry for ILLEGAL—unrecognized WOS command  
 line 1070 : get error code into X register  
 line 1080 : and jump to error handling routine

# 3 ASCII to Binary Conversions

An important aspect of interactive machine code is the ability to convert strings of ASCII characters into their hexadecimal equivalents, so that they may be manipulated by the processor. In this chapter we shall examine, with program examples, how this is performed. The routines provide the following conversions:

1. Single ASCII hex characters into binary.
2. Four ASCII hex digits into two hex bytes.
3. Signed ASCII decimal string into two signed hex bytes.

## ASCII HEX TO BINARY CONVERSION

This routine converts a hexadecimal ASCII character in the accumulator into its four-bit binary equivalent. For example, if the accumulator contains \$37 (that is, ASC“7”), the routine will result in the accumulator holding \$7, or 00001111 binary. Similarly, if the accumulator holds \$46 (ASC“F”) the routine will return \$F, or 00001111, in the accumulator.

Conversion is quite simple, and Table 3.1 gives some indication of what is required.

**Table 3.1**

Hex	Binary value	ASCII value	ASCII binary
0	00000000	\$30	00110000
1	00000001	\$31	00110001
2	00000010	\$32	00110010
3	00000011	\$33	00110011
4	00000100	\$34	00110100

**Table 3.1 (contd.)**

5	0000101	\$35	0110101
6	0000110	\$36	0110110
7	0000111	\$37	0110111
8	0001000	\$38	0111000
9	0001001	\$39	0111001
A	0001010	\$41	0100001
B	0001011	\$42	0100010
C	0001100	\$43	0100011
D	0001101	\$44	0100010
E	0001110	\$45	0100101
F	0001111	\$46	0100110

The conversion of ASCII characters 0 to 9 is straightforward. All we need to do is mask off the high nibble of the character's ASCII code. For example ASC "1" is \$31 or 00110001 binary—masking the high nibble with AND \$0F results in 00000001. Converting ASCII characters A and F is a little less obvious, however. If the high nibble of the code is masked off, then the remaining bits are 9 less than the hex required. For example, the ASCII for the letter 'D' is \$44 or 01000100. Masking the high nibble with AND \$0F gives 4, or 00000100, and adding 9 to this gives:

$$\begin{array}{r}
 0000100 \\
 + 0000100 \\
 \hline
 0000110
 \end{array}$$

the binary value for \$D.

### Program 2

```

10 REM ** CONVERT ASCII CHARACTER IN **
20 REM ** ACCUMULATOR TO BINARY **
30 REM ** REQUIRES 20 BYTES OF MEMORY **
40 :
50 CODE=49152
60 FOR LOOP=0 TO 20
70 READ BYTE
80 POKE CODE+LOOP,BYTE
90 NEXT LOOP
100 :
```

```

11Ø REM ** M/C DATA **
12Ø DATA 2Ø1,48      REM  CMP  #Ø3Ø
13Ø DATA 144,15     REM  BCC  ØØF
14Ø DATA 2Ø1,58     REM  CMP  #Ø3A
15Ø DATA 144,8      REM  BCC  ØØ8
16Ø DATA 233,7      REM  SBC  #ØØ7
17Ø DATA 144,7      REM  BCC  ØØ7
18Ø DATA 2Ø1,64     REM  CMP  #Ø4Ø
19Ø DATA 176,2      REM  BCS  ØØ2
2ØØ ::              REM  ZERO-NINE
21Ø DATA 41,15      REM  AND  #ØØF
22Ø ::              REM  RETURN
23Ø DATA 96         REM  RTS
24Ø ::              REM  ILLEGAL
25Ø DATA 56         REM  SEC
26Ø DATA 96         REM  RTS
27Ø :
28Ø :
29Ø REM ** TESTING ROUTINE **
3ØØ TEST=49184
31Ø FOR LOOP=Ø TO 14
32Ø  READ BYTE
33Ø  POKE TEST+LOOP,BYTE
34Ø NEXT LOOP
35Ø :
36Ø REM ** M/C TEST DATA **
37Ø ::              REM  TEST
38Ø DATA 32,228,255 REM  JSR  ØFFE4
39Ø DATA 24Ø,251    REM  BEQ  ØFB
4ØØ DATA 32,Ø,192  REM  JSR  ØCØØØ
41Ø DATA 144,2     REM  BCC  ØØ2
42Ø DATA 169,255   REM  LDA  #ØFF
43Ø ::              REM  OVER
44Ø DATA 133,251   REM  STA  ØFB
45Ø DATA 96        REM  RTS
46Ø :
47Ø PRINT CHR$(147)
48Ø PRINT"HIT A HEX CHARACTER KEY, AND ITS
      BINARY"
49Ø PRINT"EQUIVALENT VALUE WILL BE PRINTED"

```



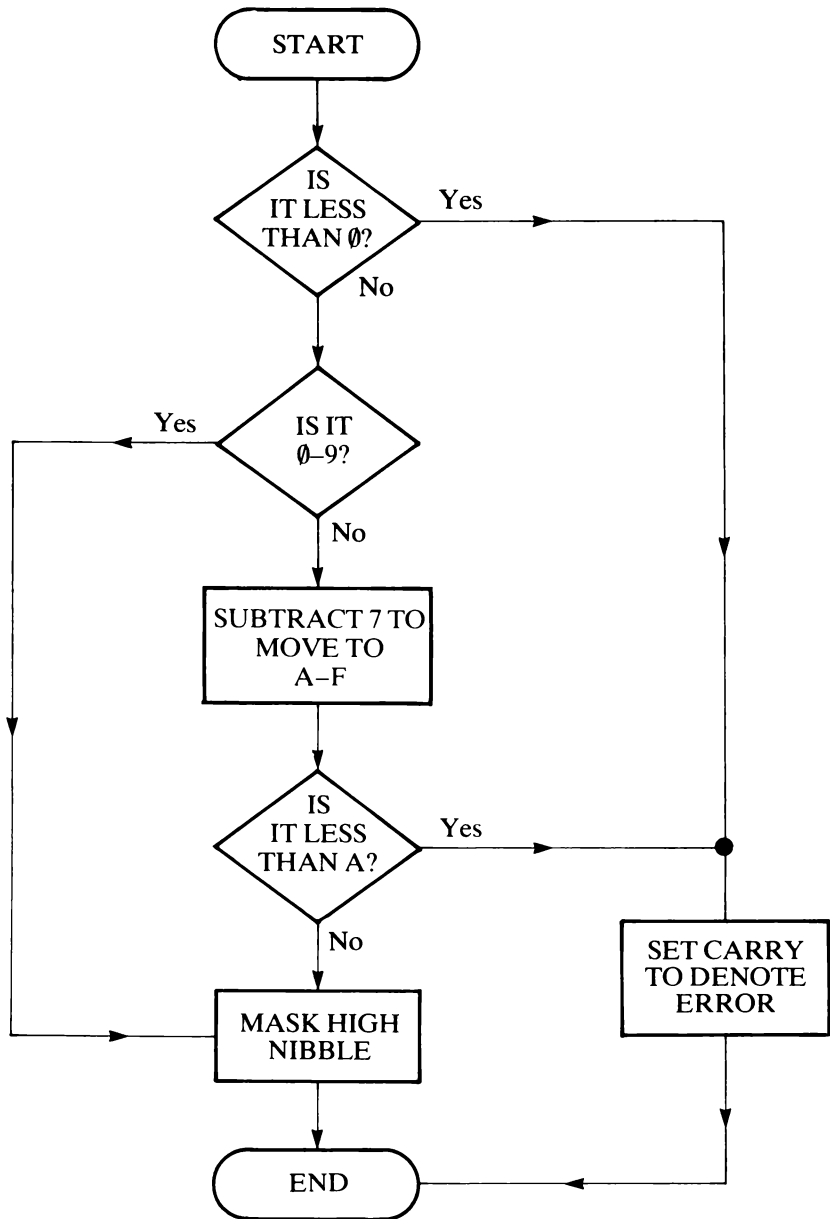


Figure 3.1 Conversion flowchart

```

500 :
510 SYS TEST
520 :
530 PRINT "RESULT = "PEEK(251)
  
```

Program 2 contains a short demonstration, prompting for a hexadecimal value key to be pressed (i.e. 0 to F) and returning its hexadecimal code. Thus, pressing the 'A' key will produce a result of 41.

The ASCII-BINARY routine begins by checking for the legality of the character, by comparing it with 48 (\$30). If the value in the accumulator is less than ASC"0", the Carry flag will be cleared, signalling an error. If the character is legal, the contents are then compared with 58 (\$3A), which is one greater than the ASCII code for 9. This part of the routine ascertains whether the accumulator's contents are in the range \$30 to \$39. If they are, the Carry flag will be cleared and the branch to ZERO-NINE (lines 150 and 120) performed. The high nibble is then masked off to complete the conversion.

If the branch of line 150 fails, a legality check for the hex characters A to F is performed. This is done by subtracting 7 from the accumulator's contents, which should bring the value it holds down below 64 (\$40), or one less than the ASCII code for the letter 'A'. At this point the Carry flag is set (it was previously set as the branch of the previous line was not performed), and the CMP#\$40 of line 180 clears it if the contents are higher than 64. The routine then masks off the high nibble, leaving the correct binary.

The following example shows how the conversion of ASC"F" to \$F works:

Mnemonic	Accumulator	Carry flag
	\$46 (ASC"F")	
CMP #30	\$46	1
BCC ILLEGAL		
CMP #3A	\$46	1
BCC ZERO-NINE		
SBC #7	\$3F	1
BCC ILLEGAL		
CMP #40	\$3F	0
BCS RETURN		
AND 0F	0F	0
RTS		

Note that this routine indicates an error by returning with the Carry flag set, so any calls to the conversion routine should always check for this on return. The short test routine does this, and loads the accumulator with \$FF to signal the fact.

Using two calls to this routine would allow two-byte hex values to be input and converted into a full eight-byte value. On completion of the first call, the accumulator's contents would need to be shifted into the high nibble.

The coding might look like this:

```

: REM WAIT
JSR GETIN          : REM GET FIRST CHARACTER
BEQ WAIT1
JSR ASCII-BINARY  : REM CONVERT TO BINARY
BCS REPORT-ERROR  : REM NON-HEX IF C=1
ASL A              : REM MOVE INTO HIGHER
                  NIBBLE

ASL A
ASL A
ASL A
STA HIGH-NIBBLE   : REM SAVE RESULT
                  : REM WAIT2

JSR GETIN          : REM GET SECOND CHARACTER
BEQ WAIT2
JSR ASCII-BINARY  : REM CONVERT TO BINARY
BCS REPORT-ERROR  : REM NON-HEX IF C=1
ORA HIGH-NIBBLE   : REM ADD HIGH NIBBLE
                  : REM ALL BINARY NOW IN
                  ACCUMULATOR
```

Using this routine and entering, say, \$FE will return 11111110 in the accumulator.

### Line-by-line

A line-by-line description of Program 2 follows:

```

line 120 : is it >= than ASC“0”?
line 130 : no, branch to ILLEGAL
line 140 : is it in range 0-9?
line 150 : yes, branch to ZERO-NINE to skip A-F
          translation.

line 160 : move onto ASCII codes for A-F
line 170 : branch to ILLEGAL if Carry flag clear
line 180  is it higher than ASC“@”?
line 190 : no, branch to ILLEGAL
line 200 : entry for ZERO-NINE
line 210 : clear high nibble
line 220 : entry for RETURN
```

line 23Ø : return with binary in accumulator  
 line 24Ø : entry for ILLEGAL  
 line 25Ø : set Carry flag to denote an error  
 line 26Ø : return to BASIC  
 line 37Ø : entry for TEST  
 line 38Ø : read keyboard  
 line 39Ø : if null string, branch to TEST  
 line 4ØØ : call conversion at \$CØØØ  
 line 41Ø : if no errors, branch OVER  
 line 42Ø : else error, place 255 in accumulator  
 line 43Ø : entry for OVER  
 line 44Ø : save accumulator in \$FB  
 line 45Ø : and return to BASIC

## FOUR ASCII DIGITS TO HEX

We can use the ASCII-BINARY routine as the main subroutine in a piece of coding which will convert four ASCII digits into a two-byte hexadecimal number, making the routine most useful for inputting two-byte hexadecimal addresses. For example, the routine would convert the ASCII string "CAFE" into a two-byte binary number 11ØØ1Ø1Ø 1111111Ø or \$CAFE. Program 3 lists the entire coding:

### Program 3

```

1Ø REM  ** CONVERT FOUR ASCII DIGITS INTO **
2Ø REM  ** A TWO-BYTE HEXADECIMAL NUMBER **
3Ø CODE=49152
4Ø FOR LOOP=0 TO 62
5Ø  READ BYTE
6Ø  POKE CODE+LOOP,BYTE
7Ø NEXT LOOP
8Ø  :
9Ø REM  ** M/C DATA **
1ØØ DATA 16Ø,Ø      : REM  LDY #Ø
11Ø DATA 162,251   : REM  LDX # $FB
12Ø DATA 148,Ø     : REM  STY $ØØ,X
13Ø DATA 148,1     : REM  STY $Ø1,X
14Ø DATA 148,2     : REM  STY $Ø2,X
15Ø ::              REM  NEXT-CHARACTER
  
```

```

160 DATA 185,60,3      : REM LDA $33C,Y
170 DATA 32,42,192    : REM JSR $C02A
180 DATA 176,21       : REM BCS $15
190 DATA 10,10        : REM ASL A : ASL A
200 DATA 10,10        : REM ASL A : ASL A
210 DATA 148,2        : REM STY $02,X
220 DATA 160,4        : REM LDY #$04
225 ::                : REM AGAIN
230 DATA 10           : REM ASL A
240 DATA 54,0         : REM ROL $00,X
250 DATA 54,1         : REM ROL $01,X
260 DATA 136          : REM DEY
270 DATA 208,248     : REM BNE $F8
280 DATA 180,2        : REM LDY $02,Y
290 DATA 200          : REM INY
300 DATA 208,227     : REM BNE $E3
310 ::                : REM ERROR
320 DATA 181,2        : REM LDA $02,X
330 DATA 96           : REM RTS
340 :
350 REM *** ASCII-BINARY CONVERSION ***
360 DATA 201,48       : REM CMP #$30
370 DATA 144,15       : REM BCC $0F
380 DATA 201,58       : REM CMP #$3A
390 DATA 144,8        : REM BCC $08
400 DATA 233,7        : REM SBC $07
410 DATA 144,7        : REM BCC $07
420 DATA 201,64       : REM CMP #$40
430 DATA 176,2        : REM BCS $02
440 ::                : REM ZERO-NINE
450 DATA 41,15        : REM AND $0F
460 ::                : REM RETURN
470 DATA 96           : REM RTS
480 ::                : REM ILLEGAL
490 DATA 56           : REM SEC
500 DATA 96           : REM RTS
510 :
520 REM *** SET UP A TEST PROCEDURE ***
530 TEST=49232
540 FOR LOOP=0 TO 34

```

```

550 READ BYTE
560 POKE TEST+LOOP, BYTE
570 NEXT LOOP
580 :
590 REM ** TEST M/C DATA **
600 DATA 160,0 : REM LDY #000
610 DATA 162,4 : REM LDX #004
620 :: REM OVER
630 DATA 142,52,3 : REM STX $334
640 DATA 140,53,3 : REM STY $335
650 :: REM INNER
660 DATA 32,228,255 : REM JSR $FFE4
670 DATA 240,251 : REM BEQ $FB
680 DATA 174,52,3 : REM LDX $334
690 DATA 172,53,3 : REM LDY $335
700 DATA 153,60,3 : REM STA $33C,Y
710 DATA 32,210,255 : REM JSR $FFD2
720 DATA 200 : REM INY
730 DATA 202 : REM DEX
740 DATA 208,229 : REM BNE $E5
750 DATA 32,0,192 : REM JSR $C000
760 DATA 96 : REM RTS
770 :
780 PRINT CHR$(147)
790 PRINT "INPUT A FOUR DIGIT HEX NUMBER : $";
800 SYS TEST
810 PRINT
820 PRINT "THE FIRST BYTE WAS :";PEEK(251)
830 PRINT "THE SECOND BYTE WAS :";PEEK(252)

```

The machine code begins by clearing three bytes of zero page RAM pointed to by the contents of the X register (lines 100 to 140). The ASCII characters are accessed one by one from a buffer which may be resident anywhere in memory (line 160), though in this case it is the four bytes at the start of the cassette buffer. Conversion and error-detection are performed (lines 170 and 180) and the four returned bits shifted into the high four bits of the accumulator. The buffer index, which keeps track of the character position in the buffer, is saved in the third of the three bytes cleared.

The loop between lines 250 and 300 is responsible for moving the four bits through the two zero page bytes which hold the final result. In fact, with the accumulator, the whole process of the loop is to perform the operation of a 24-bit shift register. Figure 3.2

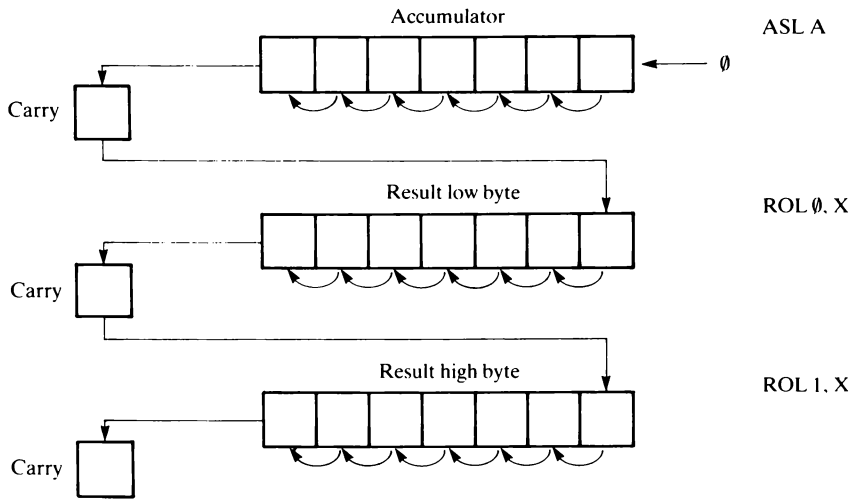


Figure 3.2 Movement of bits through a 3-byte shift register

illustrates the procedure.

The ASL A instruction shuffles the bits in the accumulator one bit to the left, with the dislodged bit 7 moving across into the Carry flag bit. This carry bit is then rotated into bit 0 of the result address low byte, which in turn rotates its bit 7 into the Carry flag. The next ROL instruction repeats this movement on the high byte. The net effect of all this is that as the process is executed four times, the returned conversions are shifted through the result address to reside in the correct place, as Figure 3.3 illustrates.

	1, X	0, X	Accumulator
Entry	00000000	00000000	11110000
1st pass	00000000	00001111	00000000
2nd pass	00000000	11110000	00000000
3rd pass	00001111	00000000	00000000
4th pass	11110000	00000000	00000000

Figure 3.3 A 24-bit shift register, showing passage of the bits in the number \$F000

Error-checking is provided for, the routine aborting when it encounters an illegal hex character, leaving the accumulator containing the index into the buffer, pointing to the illicit value. In fact, this method is used to complete the execution of the conversion-rotate loop, using a RETURN character placed at the end of the

ASCII hex string.

The test routine (lines 590 to 800) prompts for four hex-based characters to be input. These are placed in the buffer (line 610) and printed to the VDU. On completion of the input, the address-binary routine is called, and the result placed in the first two bytes of the user area, for printing or manipulation purposes.

### **Line-by-line**

A line-by-line description of Program 3 follows:

```
line 100 : clear indexing register
line 110 : get byte destination
line 120 : clear three bytes
line 150 : entry for NEXT-CHARACTER
line 160 : get character from buffer
line 170 : call ASCII-BINARY to convert
line 180 : branch to ERROR if Carry flag is set
line 190 : move low nibble into high nibble
line 210 : save index into buffer
line 220 : moving four bits
line 225 : entry for AGAIN
line 230 : move bit 7 into Carry flag
line 240 : move carry into bit 0 and bit 7 into Carry flag
line 250 : move carry into bit 0 and bit 7 into Carry flag
line 260 : decrement bit count
line 270 : and do until four bits done
line 280 : restore index into buffer
line 290 : increment it to point to next character
line 300 : do branch to NEXT-CHARACTER
line 310 : entry for ERROR
line 320 : get illegal character
line 330 : return to calling routine
```

### **CONVERT DECIMAL ASCII STRING TO BINARY**

This routine takes a signed decimal string of ASCII characters and transforms it into a two-byte hexadecimal number. For example, entering -32,678 will return the value \$8000, where \$8000 is its signed binary equivalent. Entry requirements to the conversion routine are obtained by the BASIC text in lines 880 to 940. Note



that in addition to obtaining the characters for insertion into the string buffer, the number of characters for conversion is required, this being placed in the first byte of the buffer.

#### Program 4

```

10 REM ** DECIMAL ASCII TO BINARY **
20 REM ** READ & POKE M/C DATA **
30 CODE=49152
40 FOR LOOP=0 TO 155
50 READ BYTE
60 POKE CODE+LOOP,BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 :
110 DATA 174,60,3 REM LDX $33C
120 DATA 208,3 REM BEQ $03
125 DATA 76,154,192 : REM JMP $C09A
130 DATA 160,0 REM LDY #0
140 DATA 140,55,3 REM STY $337
150 DATA 140,53,3 REM STY $335
160 DATA 150,54,3 REM STY $336
170 DATA 200 REM INY
180 DATA 140,52,3 REM STY $334
190 DATA 185,60,3 REM LDA $33C,Y
200 DATA 201,45 REM CMP #$2D
210 DATA 208,14 REM BNE $0E
220 DATA 169,255 REM LDA #$FF
230 DATA 141,55,3 REM STA $337
240 DATA 238,52,3 REM INC $334
250 DATA 202 REM DEX
260 DATA 240,113 REM BEQ $71
270 DATA 76,54,192 REM JMP $C036
280 :: REM POSITIVE
290 DATA 201,43 REM CMP #$2B
300 DATA 208,12 REM BNE $06
310 DATA 238,52,3 REM INC $334
320 DATA 202 REM DEX
330 DATA 240,100 REM BEQ $64

```

340	::		REM CONVERT-CHARACTER
350	DATA	172, 52, 3	REM LDY \$334
360	DATA	185, 60, 3	REM LDA \$33C, Y
370	::		REM CHECK-LEGALITY
380	DATA	201, 58	REM CMP #\$3A
390	DATA	16, 90	REM BPL \$5A
400	DATA	201, 48	REM CMP #\$30
410	DATA	48, 86	REM BMI \$56
420	DATA	72	REM PHA
430	DATA	14, 53, 3	REM ASL \$335
440	DATA	46, 54, 3	REM ROL \$336
450	DATA	173, 53, 3	REM LDA \$335
460	DATA	172, 53, 3	REM LDY \$336
470	DATA	14, 53, 3	REM ASL \$335
480	DATA	46, 54, 3	REM ROL \$336
490	DATA	14, 53, 3	REM ASL \$335
500	DATA	46, 54, 3	REM ROL \$336
510	DATA	24	REM CLC
520	DATA	109, 53, 3	REM ADC \$335
530	DATA	141, 53, 3	REM STA \$335
540	DATA	152	REM TYA
550	DATA	109, 54, 3	REM ADC \$336
560	DATA	141, 54, 3	REM STA \$336
570	DATA	56	REM SEC
580	DATA	104	REM PLA
590	DATA	233, 48	REM SBC #\$30
600	DATA	24	REM CLC
610	DATA	109, 53, 3	REM ADC \$335
620	DATA	141, 53, 3	REM STA \$335
630	DATA	144, 3	REM BCC \$03
640	DATA	238, 54, 3	REM INC \$336
650	::		REM NO-CARRY
660	DATA	238, 52, 3	REM INC \$334
670	DATA	202	REM DEX
680	DATA	208, 181	REM BNE \$B5
690	DATA	173, 55, 3	REM LDA \$337
700	DATA	16, 17	REM BPL \$11
710	DATA	56	REM SEC
720	DATA	169, 0	REM LDA #0
730	DATA	237, 53, 3	REM SBC \$335

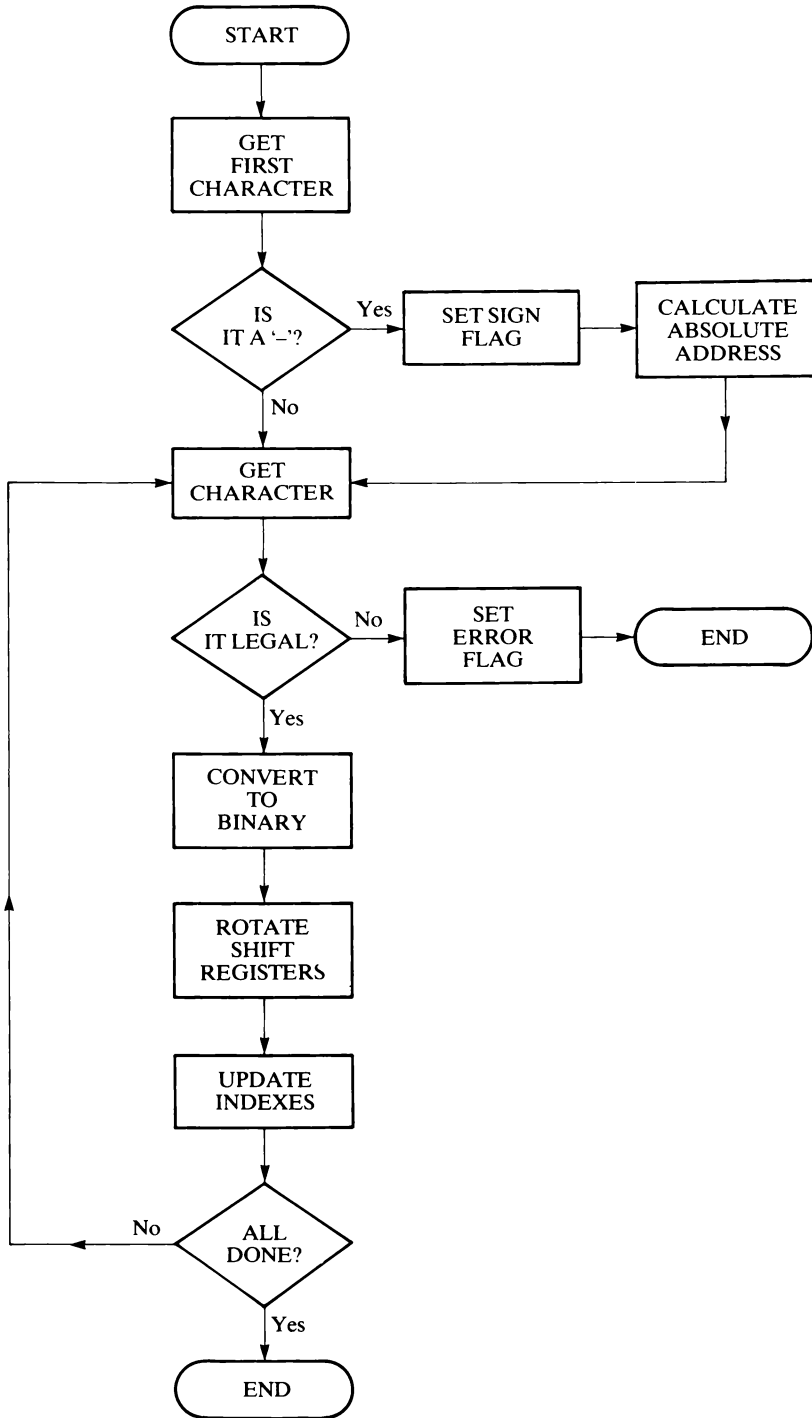


Figure 3.4 ASCII string to binary conversion flowchart

```

74Ø DATA 141,53,3      REM STA $335
75Ø DATA 169,Ø        REM LDA #Ø
76Ø DATA 237,54,3     REM SBC $336
77Ø DATA 141,54,3     REM STA $336
78Ø ::                REM NO-COMPLEMENT
79Ø DATA 24           REM CLC
80Ø DATA 144,1        REM BCC $1
81Ø ::                REM ERROR
82Ø DATA 56           REM SEC
83Ø ::                REM FINISH
84Ø DATA 96           REM RTS
85Ø :
86Ø REM ** SET UP SCREEN AND GET NUMBER **
87Ø PRINT CHR$(147)
88Ø INPUT"NUMBER FOR CONVERSION";A$
89Ø FOR LOOP=1 TO LEN(A$)
90Ø   TEMP$=MID$(A$,LOOP,1)
91Ø   B=ASC(TEMP$)
92Ø   POKE 828+LOOP,B
93Ø NEXT LOOP
94Ø POKE 828,LEN(A$)
95Ø :
96Ø SYS CODE
97Ø :
98Ø PRINT"THE TWO BYTES ARE AS FOLLOWS"
99Ø PRINT"LOW  BYTE ";PEEK(821)
100Ø PRINT"HIGH BYTE ";PEEK(822)

```

Bytes are designated as follows:

```

82Ø ($334) : string index
821 ($335)   current count
823 ($336) : sign flag
828 ($33C)   length of string
829 ($33D) : start of character string

```

The machine code begins by obtaining the character count from the X register. An error is signalled if this count is zero, otherwise the

program progresses, clearing the sign flag (used to signal positive or negative values) and result destination bytes at 'current' (lines 130 to 160). Location \$70 is used to hold the string index, pointing to the next character for conversion. This byte is initially loaded with 1 so that it skips over the count byte in the buffer.

The first byte of the string is tested for a '+' or '-' sign, the former being an optional item in the string, and the sign flag is set accordingly (lines 190 to 230). The CONVERT-CHARACTER loop starts by testing the character about to be manipulated to ensure it is a decimal value, i.e. 0 to 9 inclusive. Converting the byte into binary form is achieved by multiplying the byte by 10. This multiplication is readily available using four arithmetic shifts and an addition:  $2 * 2 * 2 + 2 = 10$ .

Because we are dealing with a two-byte result, the arithmetic shift must be performed on the two bytes, allowing bits to be transferred from one byte to the other. This is performed by using an ASL followed by a ROL. As figure 3.5 illustrates, this acts exactly like a 16-bit ASL. The first pass through this character-conversion loop has little effect, as it is operating on characters already converted, of which there are none first time round!

Lines 570 to 620 carry out the conversion of ASCII to binary and store the result. This is performed, as we know from earlier examples, by masking off the high nibble. Another technique for doing this is simply to subtract the ASCII code for '0': \$30.

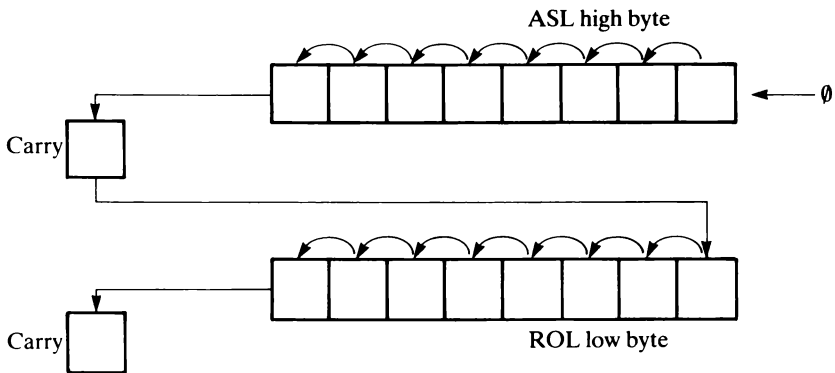


Figure 3.5 A16-bit arithmetic shift

Once all the characters have been processed, the sign flag at \$334 (820) is checked for a negative value. If this is indicated (lines 690 and 700), the value of current is subtracted from zero, thereby converting the absolute value into a signed negative byte (lines 710 to 770). The Carry flag is used to indicate any error conditions—if it is set an error occurred, and the string index at \$334 points to the illegal character.

## Line-by-line

A line-by-line description of Program 4 now follows:

```
line 110 : get length of string
line 120 : branch if not zero
line 125 : else jump to ERROR
line 130 : clear Y register
line 140 : sign flag
line 150 : and store bytes
line 170 : increment Y
line 180 : set index to first ASCII character
line 190 : get first character
line 200 : is it a minus sign?
line 210 : no, branch to POSITIVE
line 220 : yes, get negative byte
line 230 : and set the sign flag
line 240 : move to next character
line 250 : decrement length counter
line 260 : branch to ERROR if zero
line 270 : else jump to CONVERT-CHARACTER
line 280 : entry for POSITIVE
line 290 : is first character a +?
line 300 : no, branch to CHECK-LEGALITY
line 310 : yes, move to next character
line 320 : decrement length counter
line 330 : branch to ERROR if zero
line 340 : entry for CONVERT-CHARACTER
line 350 : restore index
line 360 : get character from buffer
line 370 : entry for CHECK-LEGALITY
line 380 : is it <= ASC"9"?
line 390 : no, it's bigger, branch to ERROR
line 400 : is it >= ASC"0"?
line 410 : no, branch to ERROR
line 420 : save code on stack
line 430 : multiply both bytes by two
line 450 : save low byte
```

line 46Ø : save high byte  
line 47Ø : multiply by two again (now \*4)  
line 49Ø : and again (now \*8)  
line 51Ø : clear Carry flag  
line 52Ø : add low byte \*2  
line 53Ø : and save result  
line 54Ø : transfer high byte \*2  
line 55Ø : and add to \*8 high byte  
line 56Ø : save it. Now \*1Ø  
line 57Ø : set Carry flag  
line 58Ø : restore ASCII code from stack  
line 59Ø : convert ASCII to binary  
line 6ØØ : clear Carry flag  
line 61Ø : add it to low byte current  
line 62Ø : save result  
line 63Ø : branch if NO-CARRY  
line 64Ø : else increment high byte  
line 65Ø : entry for NO-CARRY  
line 66Ø : move index on to next byte  
line 67Ø : decrement length counter  
line 68Ø : branch to CONVERT-CHARACTER if not finished  
line 69Ø : completed so get sign flag  
line 7ØØ : if clear branch to NO-COMPLEMENT  
line 71Ø : else set Carry flag  
line 72Ø : clear accumulator  
line 73Ø : and obtain two's complement  
line 74Ø : save low byte result  
line 75Ø : clear accumulator  
line 76Ø : subtract high byte from Ø  
line 77Ø : and save result  
line 78Ø : entry for NO-COMPLEMENT  
line 79Ø : clear Carry flag  
line 8ØØ : and force branch to FINISH  
line 81Ø : entry for ERROR  
line 82Ø : set Carry flag to denote error  
line 83Ø : entry for FINISH  
line 84Ø : return to BASIC

## 4 Binary to Hex ASCII

This chapter complements the previous one and illustrates how memory-based hex values can be converted into their ASCII representation. The routines provide the following conversions:

1. Print accumulator as two ASCII hex characters.
2. Print two hex bytes as four ASCII hex characters.
3. Print two-byte signed binary number as signed decimal number.

### PRINT ACCUMULATOR

To convert an eight-bit binary number into its ASCII hex equivalent characters, the procedure described in Chapter 3 must be reversed. However, because text is printed on the screen from left to right, we must deal with the high nibble of the byte first. Program 5 uses the hexprint routine to print the hexadecimal value of any key pressed at the keyboard.

#### Program 5

```
1Ø REM ** PRINT ACCUMULATOR AS A HEX NUMBER **
2Ø :
3Ø CODE=49152
4Ø FOR LOOP=Ø TO 21
5Ø READ BYTE
6Ø POKE CODE+LOOP,BYTE
7Ø NEXT LOOP
8Ø :
9Ø REM ** M/C DATA **
1ØØ :
```



```

11Ø DATA 72                REM PHA
12Ø DATA 74, 74            REM ASL A : ASL A
13Ø DATA 74, 74            REM ASL A : ASL A
14Ø DATA 32, 9, 192        REM JSR $CØØ9
15Ø DATA 1Ø4                REM PLA
16Ø ::                      REM FIRST $CØØ9
17Ø DATA 41, 15            REM AND #$ØF
18Ø DATA 2Ø1, 1Ø           REM CMP #$ØA
19Ø DATA 144, Ø2           REM BCC $Ø2
2ØØ DATA 1Ø5, 6            REM ADC #$Ø6
21Ø ::                      REM OVER
22Ø DATA 1Ø5, 48           REM ADC #$3Ø
23Ø DATA 76, 21Ø, 255      REM JMP $FFD2
24Ø :
25Ø REM ** SET UP DEMO AT 828 **
26Ø REM LDA $FB : JMP $CØØØ
27Ø POKE 828, 165 : POKE 829, 251
28Ø POKE 83Ø, 76 : POKE 831, Ø : POKE 832, 192
29Ø `PRINT CHR$(147)
3ØØ PRINT "HIT ANY KEY AND ITS HEX VALUE IN"
31Ø PRINT "ASCII WILL BE DISPLAYED"
32Ø GET A$: IF A$="" THEN GOTO 32Ø
33Ø A=ASC(A$)
34Ø POKE 251, A
35Ø :
36Ø SYS 828
37Ø REM CALL 'SYS CODE' TO USE DIRECTLY

```

The hexprint routine is embedded between lines 11Ø and 23Ø. The accumulator's contents are first pushed on to the hardware stack. This procedure is necessary as it will have to be restored before the second pass, which calculates the ASCII code for the second character. The first pass through the routine sets about moving the upper nibble of the accumulator byte into the lower nibble (lines 12Ø and 13Ø). The FIRST subroutine ensures that the high nibble is cleared by logically ANDing it with \$ØF. This is, of course, surplus to requirement on the first pass, but is needed on the second pass to isolate the low nibble. Comparing the accumulator's contents with 1Ø will ascertain whether the value is in the range Ø to 9 or A to F. If the Carry flag is clear, it falls in the lower range (Ø to 9) and simply setting bits 4 and 5, by adding \$3Ø, will give the required ASCII code. A further 7 must be added to skip non-hex ASCII codes to arrive at the ASCII codes for A to F (\$41 to \$46). You may have

noticed that line 200 does not add 7 but in fact adds one less, 6. This is because, for this section of coding to be executed, the carry must have been set, and the 6510 addition opcode references the Carry flag in addition. Therefore, the addition performed is: accumulator + 6 + 1.

The JMP of line 230 will return the program back to line 150. Remember, FIRST was called with a JSR, so the RTS from completion of the CHROUT call returns control here. The accumulator is restored and the process repeated for the second ASCII digit.

A short test routine is established in lines 250 to 340. This requests you to hit a key, the value of which is placed in a free zero page byte. The 'hand-POKed' routine at 828 is called by line 360, and puts the key's value into the accumulator before performing a jump to the main routine.

The following example illustrates the program's operation, assuming the accumulator holds the value 01001111, \$4F:

Mnemonic	Accumulator	Carry flag
	\$4F	
LSR A	\$27	1
LSR A	\$13	1
LSR A	\$09	1
LSR A	\$04	1
JSR FIRST		
AND #\$0F	\$04	1
CMP #\$0A	\$04	0
BCC OVER		
OVER		
ADC #\$30	\$34 (ASC"4")	0
JMP CHROUT		
PLA	\$4F	0
AND #\$0F	\$0F	0
CMP #\$0A		

### Line-by-line

A line-by-line description of Program 5 follows:

line 110 : save accumulator on stack

line 120 : move high nibble into low nibble

```

line 140 : call FIRST subroutine
line 150 : restore accumulator
line 160 : entry for FIRST
line 170 : ensure only low nibble set
line 180 : is it < 10?
line 190 : yes, branch to OVER
line 200 : no, add 7, value $A to $F
line 210 : entry for OVER
line 220 : add 48 to convert to ASCII code
line 230 : and print, returning to line 140 or BASIC

```

## PRINT A HEXADECIMAL ADDRESS

The hexprint routine can be extended to enable two zero page bytes to be printed out in hexadecimal form. This is an especially important procedure when writing machine based utilities, such as a hex dump or disassembler. The revamped program is listed below:

### Program 6

```

10 REM ** PRINT TWO HEX BYTES AS **
20 REM ** A TWO-BYTE ADDRESS **
30 CODE=49152
40 FOR LOOP=0 TO 34
50 READ BYTE
60 POKE CODE+LOOP, BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 REM ** CALL WITH $FB,$FC HOLDING BYTES **
110 :: REM ADDRESS-PRINT
120 DATA 162,251 : REM LDX #$FB
130 DATA 181,1 : REM LDA $01,X
140 DATA 32,13,192 : REM JSR $C00D
150 DATA 181,0 : REM LDA $00,X
160 DATA 32,13,192 : REM JSR $C00D
170 DATA 96 : REM RTS
180 :: REM HEXPRINT
190 DATA 72 : REM PHA
200 DATA 74,74 : REM LSR A : LSR A

```

```

21Ø DATA 74,74 : REM LSR A : LSR A
22Ø DATA 32,22,192 : REM JSR $CØ16
23Ø DATA 1Ø4 : REM PLA
24Ø :: REM FIRST
25Ø DATA 41,15 : REM AND # $ØF
26Ø DATA 2Ø1,1Ø : REM CMP # $ØA
27Ø DATA 144,2 : REM BCC $Ø2
28Ø DATA 1Ø5,6 : REM ADC # $Ø6
29Ø :: REM OVER
3ØØ DATA 1Ø5,48 : REM ADC # $3Ø
31Ø DATA 76,21Ø,255 : REM JMP $FFD2

```

Zero paged indexed addressing is used to access the two bytes, the crucial location being given in the X register, which acts as the index for the high byte, LDA \$Ø1,X (line13Ø), and the low byte, LDA \$ØØ,X (line 15Ø). The all-important address in this instance is \$FB (line 13Ø), so the bytes accessed by ADDRESS-PRINT are \$FB (\$FB+Ø) and \$FC (\$FB+1). Using this method, various addresses can be housed within zero page and any one reached simply by seeding the X register with the location value.

## Project

Adapt Program 6 to accept a five character decimal number from the keyboard, printing its hexadecimal value on the screen. Remember—no BASIC, and the input routine must be able to accept numbers in the range Ø to 65!

## BINARY SIGNED NUMBER TO SIGNED ASCII DECIMAL STRING

This conversion utility takes a two-byte hexadecimal number and converts it into its equivalent decimal based ASCII character string. For example, if the two-byte value is \$7FFF, the decimal string is 32,767, \$7FFF being 32,767 in decimal. The coding uses signed binary values so that if the most significant bit is set, a negative value is interpreted. This is relayed in the string with a minus sign. This means that the routine can handle values in the range 32,767 to -32,768. When using the routine, remember that the two's complement representation is used, so that a hex value of \$FFFF is converted to the string -1, and \$8ØØØ returns the character string -32,767.

The two address bytes are located at \$334 and \$335 and the string buffer from \$FB onwards. The length of the string buffer will vary, but its maximum length will not exceed six digits, so this number of bytes should be reserved.

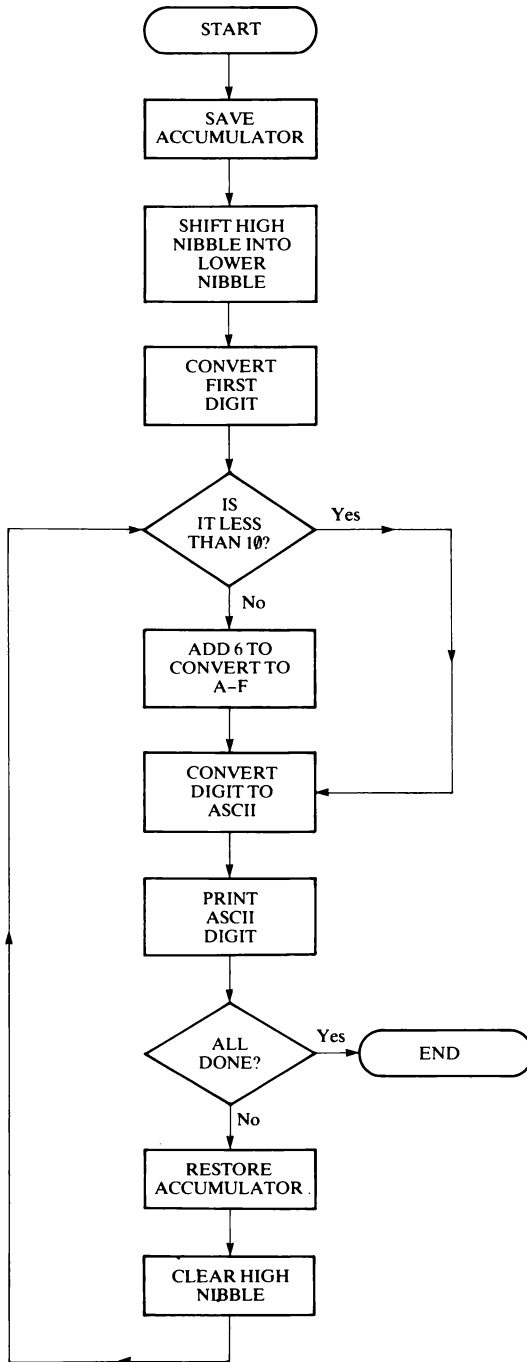


Figure 4.1 Hex to ASCII conversion flowchart

## Program 7

```
10 REM ** BINARY SIGNED NUMBER CONVERSION **
20 REM ** INTO SIGNED DECIMAL ASCII STRING **
30 CODE=49152 : OUTPUT=49301
40 FOR LOOP=0 TO 163
50 READ BYTE
60 POKE CODE+LOOP,BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 DATA 160,0 : REM LDY #00
110 DATA 152 : REM TYA
120 DATA 133,251 : REM STA $FB
130 DATA 133,252 : REM STA $FC
140 DATA 133,253 : REM STA $FD
150 DATA 133,254 : REM STA $FE
160 DATA 133,255 : REM STA $FF
170 DATA 173,53,3 : REM LDA $335
180 DATA 141,56,3 : REM STA $338
190 DATA 16,15 : REM BPL $0F
200 DATA 56 : REM SEC
210 DATA 152 : REM TYA
220 DATA 237,52,3 : REM SBC $334
230 DATA 141,52,3 : REM STA $334
240 DATA 152 : REM TYA
250 DATA 237,53,3 : REM SBC $335
260 DATA 141,53,3 : REM STA $335
270 :: REM CONVERSION
280 DATA 169,0 : REM LDA #00
290 DATA 141,54,3 : REM STA $336
300 DATA 141,55,3 : REM STA $337
310 DATA 24 : REM CLC
320 DATA 162,16 : REM LDX #10
330 :: REM LOOP
340 DATA 46,52,3 : REM ROL $334
350 DATA 46,53,3 : REM ROL $335
360 DATA 46,54,3 : REM ROL $336
370 DATA 46,55,3 : REM ROL $337
380 DATA 56 : REM SEC
```

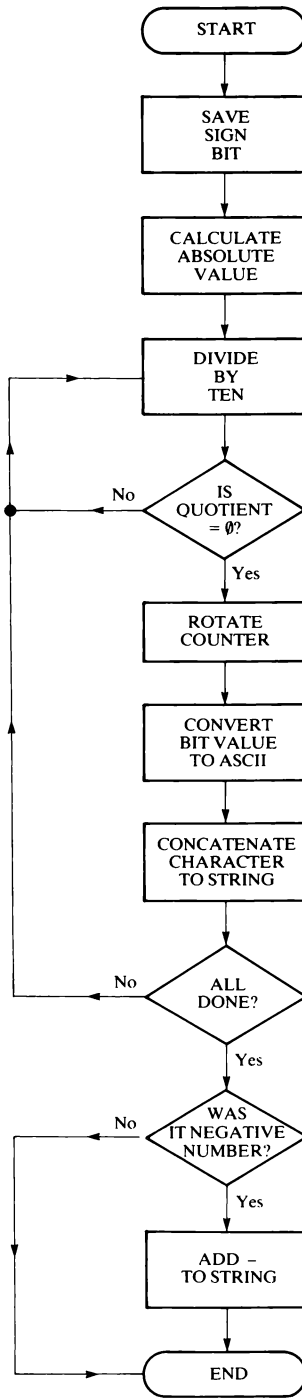


Figure 4.2 Binary to ASCII string conversion flowchart

```

39Ø DATA 173,54,3 : REM LDA $336
40Ø DATA 233,1Ø : REM SBC #ØØA
41Ø DATA 168 : REM TAY
42Ø DATA 173,55,3 : REM LDA $337
43Ø DATA 233,Ø : REM SBC #ØØØ
44Ø DATA 144,6 : REM BCC ØØ6
45Ø DATA 14Ø,54,3 : REM STY $336
46Ø DATA 141,55,3 : REM STA $337
47Ø :: REM LESS-THAN
48Ø DATA 2Ø2 : REM DEX
49Ø DATA 2Ø8,221 : REM BNE $DD
50Ø DATA 46,52,3 : REM ROL $334
51Ø DATA 46,53,3 : REM ROL $335
52Ø :: REM ADD-ASCII
53Ø DATA 24 : REM CLC
54Ø DATA 173,54,3 : REM LDA $336
55Ø DATA 1Ø5,48 : REM ADC #Ø3Ø
56Ø DATA 32,116,192 : REM JSR $CØ74
57Ø DATA 173,52,3 : REM LDA $334
58Ø DATA 13,53,3 : REM ORA $335
59Ø DATA 2Ø8,187 : REM BNE $BB
60Ø :: REM FINISHED
61Ø DATA 173,56,3 : REM LDA $338
62Ø DATA 16,5 : REM BPL ØØ5
63Ø DATA 169,45 : REM LDA #Ø2D
64Ø DATA 32,116,192 : REM JSR $CØ74
65Ø :: REM POSITIVE
66Ø DATA 96 : REM RTS
67Ø REM SUBROUTINE TO FORM ASCII CHARACTER
STRING IN $FB
68Ø :: REM CONCATENATE
69Ø DATA 72 : REM PHA
70Ø DATA 16Ø,Ø : REM LDY #ØØØ
71Ø DATA 185,251,Ø : REM LDA $ØØFB,Y
72Ø DATA 168 : REM TAY
73Ø DATA 24Ø,11 : REM BEQ ØØB
74Ø :: REM SHUFFLE-ALONG
75Ø DATA 185,251,Ø : REM LDA $ØØFB,Y
76Ø DATA 2ØØ : REM INY
77Ø DATA 153,251,Ø : REM STA $ØØFB,Y

```



```

78Ø DATA 136,136      : REM  DEY : DEY
79Ø DATA 2Ø8,245     : REM  BNE $F5
8ØØ  ::              REM  ZERO-FINISH
81Ø DATA 1Ø4        : REM  PLA
82Ø DATA 16Ø,1      : REM  LDY #ØØ1
83Ø DATA 153,251,Ø  : REM  STA $ØØFB,Y
84Ø DATA 136        : REM  DEY
85Ø DATA 182,251    : REM  LDX $FB,Y
86Ø DATA 232        : REM  INX
87Ø DATA 15Ø,251    : REM  STX $FB,Y
88Ø DATA 96         : REM  RTS
89Ø REM STRING PRINTING ROUTINE
9ØØ  ::              REM  STRING-PRINT
91Ø DATA 166,251    : REM  LDX $FB
92Ø DATA 16Ø,1      : REM  LDY #ØØ1
93Ø  ::              REM  PRINT-LOOP
94Ø DATA 185,251,Ø  : REM  LDA $FB,Y
95Ø DATA 32,21Ø,255 : REM  JSR $FFD2
96Ø DATA 2ØØ        : REM  INY
97Ø DATA 2Ø2        : REM  DEX
98Ø DATA 2Ø8,246    : REM  BNE $F6
99Ø DATA 96         : REM  RTS
1ØØØ :
1Ø1Ø REM ** GET IN A HEX NUMBER **
1Ø2Ø PRINT CHR$(147) : PRINT
1Ø3Ø PRINT"INPUT A HEX NUMBER :$";
1Ø4Ø GOSUB 2ØØØ
1Ø5Ø POKE 82Ø,LOW      : REM  LOW BYTE HEX
      NUMBER
1Ø6Ø GOSUB 2ØØØ
1Ø7Ø POKE 821,HIGH     : REM  HIGH BYTE HEX
      NUMBER
1Ø8Ø :
1Ø9Ø SYS CODE         : REM  CALL CONVERSION
11ØØ :
111Ø PRINT"ITS DECIMAL EQUIVALENT IS :";
112Ø SYS OUTPUT
113Ø END
114Ø :
1999 REM ** HEX INPUT CONTROL **

```

```

2000 GOSUB 2500
2010 F=NUM : PRINT Z$;
2020 GOSUB 2500
2030 S=NUM : PRINT Z$;
2040 HIGH=F*16+S
2050 GOSUB 2500
2060 F=NUM : PRINT Z$;
2070 GOSUB 2500
2080 S=NUM : PRINT Z$
2090 LOW=F*16+S
2100 RETURN
2200 :
2499 REM ** GET HEX ROUTINE **
2500 GET Z$
2510 IF Z$="" THEN GOTO 2500
2520 IF Z$>"F" THEN GOTO 2500
2530 IF Z$="A" THEN NUM=10: RETURN
2540 IF Z$="B" THEN NUM=11: RETURN
2550 IF Z$="C" THEN NUM=12: RETURN
2560 IF Z$="D" THEN NUM=13: RETURN
2570 IF Z$="E" THEN NUM=14: RETURN
2580 IF Z$="F" THEN NUM=15: RETURN
2590 NUM=VAL(Z$) : RETURN

```

### Functional bytes:

251–255	(\$FB–\$FF)	: ASCII string buffer
820–821	(\$334–\$335)	: binary address for conversion
822–823	(\$336–\$337)	: temporary storage
824	(\$338)	: sign flag

To demonstrate the routine's workings, the program first prompts for a hexadecimal number using the BASIC hex loader subroutine at line 2000. This is evaluated and placed at BINARY-ADDRESS by lines 1050 and 1070.

The program proper begins by clearing the string buffer area (lines 100 to 160), an important procedure which ensures no illicit characters find their way into the ASCII string. The sign of the number is tested by loading the high byte of the address byte into the accumulator and saving its value in the sign flag byte. This process will condition the Negative flag. If it is set, a negative number is interpreted and the plus branch to CONVERSION (line

190) fails. The next seven operations obtain the absolute value of the two-byte number by subtracting it from itself and the set carry bit. Thus \$FFFF will result in an absolute value of 1 and \$8000 an absolute value of 32,678.

The two flows of the program rejoin at line 280, where the two temporary bytes are cleared. These bytes are used in conjunction with the binary address bytes to form a 32-bit shift register, allowing bits to flow from the low byte address to the high byte of temporary.

The loop of lines 340 to 510 performs the conversion, by successively dividing through by ten until the quotient has a value of zero. By this time the binary equivalent of this ASCII character being processed will have been placed in the temporary byte. To produce this, the loop needs sixteen iterations so the X register is used to count these out. Converting the binary to hex involves simply adding \$30 or ASC"0" to it (lines 530 to 550).

Because it may not be immediately clear what is happening, Table 4.1 shows the values of the accumulator and four associated bytes after each of the 16 passes of the loop, when converting \$FFFF into its absolute ASCII value of 1. It should be clear from this how the bits shuffle their way through the four byte 'register'.

**Table 4.1**

Iteration	Accumulator	\$334	\$335	\$336	\$337
1	00	01	00	00	00
2	FF	02	00	00	00
3	FF	04	00	00	00
4	FF	08	00	00	00
5	FF	10	00	00	00
6	FF	20	00	00	00
7	FF	40	00	00	00
8	FF	80	00	00	00
9	FF	00	00	01	00
10	FF	00	00	01	00
11	FF	00	00	01	00
12	FF	00	00	01	00
13	FF	00	00	01	00
14	FF	00	00	01	00
15	FF	00	00	01	00
16	FF	00	00	01	00

All that is now required is for this character to be added to the string buffer. This concatenation is completed by the code of lines 690 to 880. This began by obtaining the buffer index, which contains the current number of characters already concatenated. This is stored in the first byte of the buffer, \$FB in this instance. It is then moved across into the accumulator. Next, lines 750 to 790 move any characters present in the buffer up memory one byte, thereby opening up a gap of one byte into which the newly formed character can be placed (lines 810 to 870). The buffer index is also incremented and restored at this point, before an RTS is made back to the main body of the program.

End of program operation is tested for by logically ORing the contents of the high and low bytes of the address. If the result is zero, all bits have been rotated and dealt with, in which case the sign flag byte is tested to ascertain whether a minus sign need be placed at the start of the ASCII string (lines 600 to 660).

## Line-by-line

A line-by-line description of Program 7 follows:

```
line 100 : clear Y register
line 110 : and accumulator
line 120 : and then the five buffer bytes
line 170 : get high byte for conversion
line 180 : save in sign flag
line 190 : if positive branch to CONVERSION
line 200 : else set Carry flag
line 210 : clear accumulator
line 220 : obtain absolute value of low byte
line 230 : and save
line 240 : clear accumulator
line 250 : obtain absolute value of high byte
line 260 : and save
line 270 : entry for CONVERSION
line 280 : clear accumulator
line 290 : clear temporary storage bytes
line 310 : clear Carry flag
line 320 : sixteen bits to process
line 330 : entry for LOOP
line 340 : move bit 7 into Carry flag
line 350 : and on into bit 0
line 360 : move bit 7 into Carry flag
```

line 370 : and on into bit 0  
line 380 : set Carry flag  
line 390 : get low byte of temp  
line 400 : subtract 10  
line 410 : save result in Y  
line 420 : get high byte of temporary  
line 430 : subtract carry bit  
line 440 : branch to LESS-THAN if divisor>dividend  
line 450 : else save result of operation in temporary  
line 470 : entry for LESS-THAN  
line 480 : decrement bit count  
line 490 : branch to LOOP until 16 bits done  
line 500 : rotate bit 7 into Carry flag  
line 510 : and on into bit 0  
line 520 : entry for ADD-ASCII  
line 530 : clear Carry flag  
line 540 : get low byte from temporary  
line 550 : convert into ASCII character  
line 560 : concatenate on to string in buffer  
line 570 : get low byte of binary number  
line 580 : OR with high byte. If 0 then all done  
line 590 : if not finished branch to CONVERSION  
line 600 : entry for FINISHED  
line 610 : get sign  
line 620 : if N = 0 branch to POSITIVE  
line 630 : otherwise get ASC“-”  
line 640 : and add it to final string  
line 650 : entry for POSITIVE  
line 660 : back to BASIC  
line 680 : entry for CONCATENATE, \$C074  
line 690 : save accumulator  
line 700 : initialize index  
line 710 : and get buffer length  
line 720 : move it into Y for indexing  
line 730 : if 0 branch to ZERO-LENGTH  
line 740 : entry for SHUFFLE-ALONG  
line 750 : get character from buffer  
line 760 : increment index  
line 770 : save character one byte along  
line 780 : restore original address minus one

line 790 : branch to SHUFFLE-ALONG until completed  
line 800 : entry for ZERO-FINISH  
line 810 : restore accumulator  
line 820 : index past length byte  
line 830 : add character to buffer  
line 840 : decrement index  
line 850 : get length byte  
line 860 : increment it  
line 870 : save it  
line 880 : back to calling routine  
line 900 : entry for OUTPUT  
line 910 : get length of string as counter  
line 920 : set index to first character  
line 930 : entry for PRINT-LOOP  
line 940 : get character  
line 950 : print it  
line 960 : increment index  
line 970 : decrement count  
line 980 : branch to PRINT-LOOP until all done  
line 990 : back to BASIC

## 5 String Manipulation

In this chapter we will look at how ASCII character strings can be manipulated using machine code routines to perform the following operations:

1. Compare two strings.
2. Concatenate one string onto another.
3. Copy a substring from within a main string.
4. Insert a substring into a main string.

These types of routines are essential if you intend to write any programs that manipulate data and information. Adventure games are a typical example of this kind of program.

### COMPARING STRINGS

String comparison is normally performed after the computer user has input some information from the keyboard. In BASIC this might be written as:

```
100 A$="MOVE LEFT"
110 INPUT"WHICH DIRECTION?"; B$
120 IF A$=B$ THEN PRINT "CORRECT!"
```

We do not always wish to test for equality, however. In BASIC, we are able to test for unlike items using the NOT operators '<>'. Thus, line 120 could have been written as:

```
120 IF A$ <> B$ PRINT "WRONG!"
```

At other times, we may wish to test which of two strings has a greater length, and this is possible in BASIC using the LEN statement:

```
21Ø IF LEN(A$) > LEN(B$) THEN PRINT "FIRST"
```

Program 8 gives the assembler and BASIC listing for the string comparison routine, which puts all the functions described above at your disposal whenever the program is used. The Status register holds these answers in the Zero and Carry flags. The Zero flag is used to signal equality: if it is set ( $Z=1$ ), the two strings compared were identical; if it is cleared ( $Z=0$ ) they were dissimilar.

The Carry flag returns information as to which of the two strings was the longer: if it is set ( $C=1$ ), they were identical in length or the first string was the larger. The actual indication required here is evaluated in conjunction with the Zero flag. If  $Z=0$  and  $C=1$ , then a longer string rather than an equal-length string is indicated, but if the Carry flag is returned clear ( $C=0$ ), then the second string was longer than the first.

### Program 8

```
1Ø REM ** STRING COMPARISON ROUTINE **
2Ø CODE=49152
3Ø TEST=49184
4Ø FOR LOOP=Ø TO 41
5Ø READ BYTE
6Ø POKE CODE+LOOP,BYTE
7Ø NEXT LOOP
8Ø :
9Ø REM ** M/C DATA **
1ØØ DATA 173,52,3 : REM LDA $334
11Ø DATA 2Ø5,53,3 : REM CMP $335
12Ø DATA 144,3 : REM BCC $Ø3
13Ø DATA 174,53,3 : REM LDX $335
14Ø :: REM COMPARE-STRING
15Ø DATA 24Ø,12 : REM BEQ $ØC
16Ø DATA 16Ø,Ø : REM LDY #$$ØØ
17Ø :: REM COMPARE-BYTES
18Ø DATA 177,251 : REM LDA ($FB),Y
19Ø DATA 2Ø9,253 : REM CMP ($FD),Y
2ØØ DATA 2Ø8,1Ø : REM BNE $ØA
21Ø DATA 2ØØ : REM INY
22Ø DATA 2Ø2 : REM DEX
23Ø DATA 2Ø8,246 : REM BNE $F6
24Ø :: REM CONDITION-FLAGS
25Ø DATA 173,52,3 : REM LDA $334
```



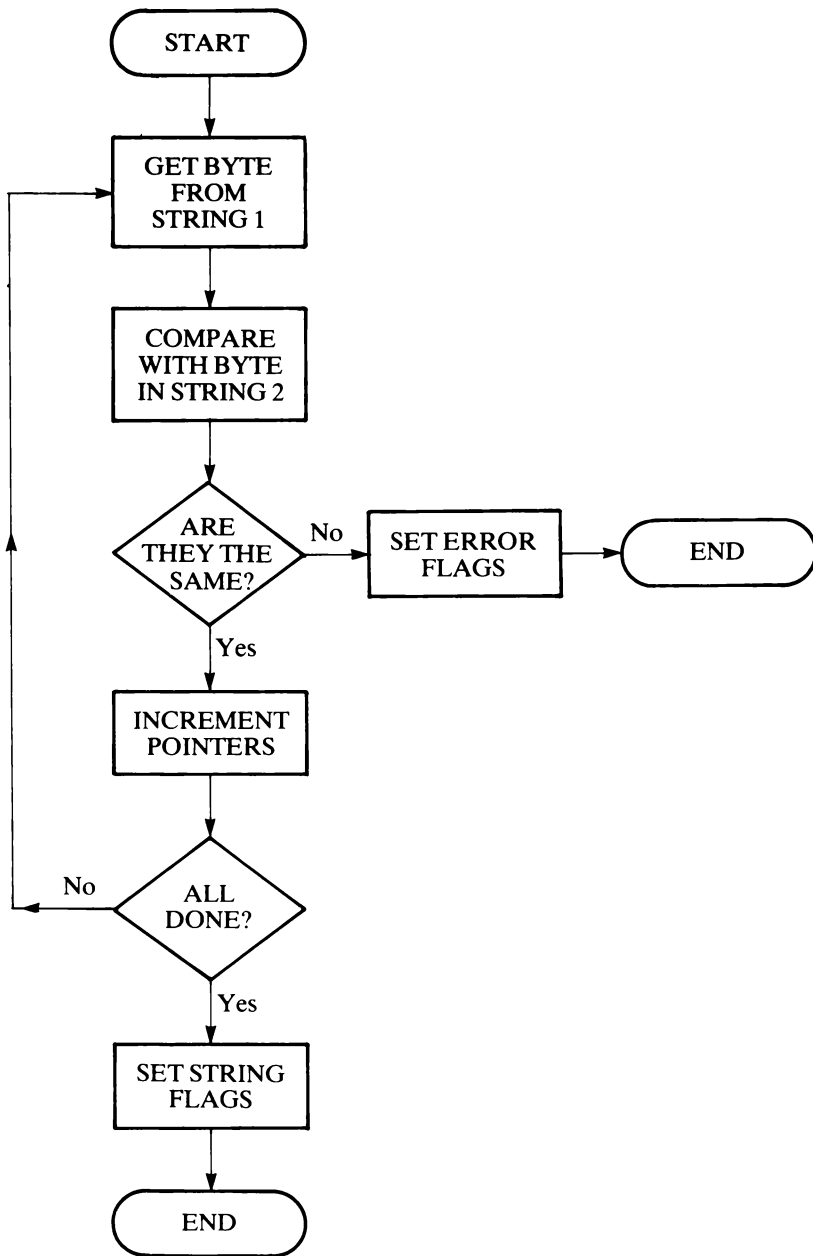


Figure 5.1 Compare strings flowchart

```

260 DATA 205,53,3 : REM CMP $335
270 :: REM FINISH
280 DATA 96 : REM RTS
290 :
300 :: REM TEST ROUTINE
310 DATA 32,0,192 : REM JSR $C000
  
```

```

320 DATA 8 : REM PHP
330 DATA 104 : REM PLA
340 DATA 41,3 REM AND #03
350 DATA 133,251 : REM STA $FB
360 DATA 96 : REM RTS
370 :
380 REM ** SET UP STRINGS FOR COMPARISON **
390 PRINT CHR$(147)
400 INPUT "FIRST STRING : ";A$
410 FOR LOOP=1 TO LEN(A$)
420 TEMP$=MID$(A$,LOOP,1)
430 A=ASC(TEMP$)
440 POKE 50432+LOOP-1,A
450 NEXT LOOP
460 :
470 INPUT "SECOND STRING : ";B$
480 FOR LOOP=1 TO LEN(B$)
490 TEMP$=MID$(B$,LOOP,1)
500 B=ASC(TEMP$)
510 POKE 50688+LOOP-1,B
520 NEXT LOOP
530 :
540 POKE 251,0 : POKE 252,197
550 POKE 253,0 : POKE 254,198
560 POKE 820,LEN(A$) POKE 821,LEN(B$)
570 :
580 SYS TEST
590 :
600 PRINT "RESULT IS : ";PEEK(251)

```

Bytes reserved:

251-252	(\$FB-\$FC)	: address of first string
253-254	(\$FD-\$FE)	: address of second string
820	(\$334)	: length of first string
821	(\$335)	: length of second string

Once run, the BASIC text of lines 380 to 520 calls for two strings to be input. These are stored in memory from \$C500 and \$C600. Note that the routine cannot handle strings greater than 256 characters in length (though it could of course be expanded to do so). The length

of each string is also required by the routine, so this is ascertained and stored in the appropriate zero page bytes at \$334 and \$335 (line 560).

To allow the string buffers to be fully relocatable, the string addresses are held in two zero page vectors (lines 540 and 550).

String comparison proper starts by evaluating the length bytes to find out if they are the same length. If they are not equal, then the strings cannot be identical. However, as the routine returns information about the lengths of the strings it is still completed—in this case the program compares bytes through the length of the smaller of the two strings.

Byte comparison is performed by lines 170 to 190, using post-indexed indirect addressing. On the first non-equal characters the main loop is exited to FINISH. Assuming the entire comparison works, and the X register, which holds the working string length, has been decremented to zero, the length bytes (lines 250 and 260) are compared to condition the Zero and Carry flags before the routine completes.

The short test routine returns the Zero and Carry flag values and prints them out, indicating the following results:

Returned	Z	C	Result
0	0	0	Strings <> and string 1 larger
1	0	1	Strings <> and string 2 larger
3	1	1	Strings =

### Line-by-line

A line-by-line description of Program 8 follows:

```

line 100 : get length of first string
line 110 : is it the same length as the second string?
line 120 : no, it's longer, so branch to COMPARE-STRING
line 130 : yes, so get length of second string
line 140 : entry for COMPARE-STRING
line 150 : if zero, branch to CONDITION-FLAGS
line 160 : initialize indexing register
line 170 : entry for COMPARE-BYTES
line 180 : get character from first string
line 190 : compare to same character in second string
line 200 : if dissimilar, branch to FINISH
line 210 : increment index

```

```

line 22Ø : decrement string counter
line 23Ø : branch back to COMPARE-BYTES until zero
line 24Ø : entry for CONDITION-FLAG
line 25Ø : get length of first string
line 26Ø : compare with length of the second string
line 27Ø : entry for finish
line 28Ø : back to calling routine
line 30Ø : entry for TEST routine
line 31Ø : push status onto stack
line 32Ø : pull into accumulator
line 33Ø : save Z and C
line 34Ø : save at location $FB
line 35Ø : back to BASIC

```

## STRINGS UNITE

Strings may be joined together by a process called 'concatenation'. In BASIC the addition operator '+' performs this function. Thus the program:

```

10Ø A$="REM"
11Ø B$="ARK"
12Ø C$=A$+B$

```

assigns the string 'REMARK' to the string C\$. If line 12Ø were rewritten as:

```

12Ø C$=B$+A$

```

the resultant value assigned to C\$ would be 'ARKREM'. We can see from this that one string is simply tagged on to the end of the other, overwriting the former's RETURN character, but preserving the latter's.

This process of concatenation can be performed quite readily as Program 9 illustrates. However, the actual BASIC equivalent of the operation we are performing here is:

```

A$=A$+B$

```

In other words, we are adding the second string on to the first string, rather than summing the two to give a separate final string, although this is possible with slight modifications to the assembler text.

## Program 9

```
10 REM ** STRING CONCATENATION **
20 CODE=49152
30 FOR LOOP=0 TO 96
40 READ BYTE
50 POKE CODE+LOOP,BYTE
60 NEXT LOOP
70 :
80 REM ** M/C DATA **
90 :: REM STRING-CONCATENATION
100 DATA 173,52,3 : REM LDA $334
110 DATA 141,54,3 : REM STA $336
120 DATA 169,0 : REM LDA #$00
130 DATA 141,55,3 : REM STA $337
140 DATA 24 : REM CLC
150 DATA 173,53,3 : REM LDA $335
160 DATA 109,52,3 : REM ADC $334
170 DATA 176,3 : REM BCS $03
180 DATA 76,45,192 : REM JMP $C02D
190 :: REM TOO-LONG
200 DATA 169,255 : REM LDA #$FF
210 DATA 141,57,3 : REM STA $339
220 DATA 56 : REM SEC
230 DATA 237,52,3 : REM SBC $334
240 DATA 144,51 : REM BCC $33
250 DATA 141,56,3 : REM STA $338
260 DATA 169,255 : REM LDA #$FF
270 DATA 141,52,3 : REM STA $334
280 DATA 76,59,192 : REM JMP $C03B
290 :: REM GOOD-LENGTH
300 DATA 141,52,3 : REM STA $334
310 DATA 169,0 : REM LDA #$00
320 DATA 141,57,3 : REM STA $339
330 DATA 173,53,3 : REM LDA $335
340 DATA 141,56,3 : REM STA $338
350 :: REM CONCATENATION
360 DATA 173,56,3 : REM LDA $338
370 DATA 240,21 : REM BEQ $15
380 :: REM LOOP
```

```

39Ø DATA 172,55,3 : REM LDY $337
40Ø DATA 177,253 : REM LDA ($FD),Y
41Ø DATA 172,54,3 : REM LDY $336
42Ø DATA 145,251 : REM STA ($FB),Y
43Ø DATA 238,54,3 : REM INC $336
44Ø DATA 238,55,3 : REM INC $337
45Ø DATA 2Ø6,56,3 : REM DEC $338
46Ø DATA 2Ø8,235 : REM BNE $EB
47Ø :: REM FINISHED
48Ø DATA 172,52,3 : REM LDY $334
49Ø DATA 169,13 : REM LDA #$ØD
50Ø DATA 145,251 : REM STA ($FB),Y
51Ø DATA 173,57,3 : REM LDA $339
52Ø DATA 1Ø6 : REM ROR A
53Ø DATA 96 : REM RTS
54Ø :
60Ø PRINT CHR$(147)
61Ø INPUT "FIRST STRING ";A$
62Ø INPUT "SECOND STRING ";B$
63Ø :
64Ø F=49664 : REM $C2ØØ
65Ø S=4992Ø : REM $C3ØØ
66Ø :
67Ø FOR LOOP=1 TO LEN(A$)
68Ø TEMP$=MID$(A$,LOOP,1)
69Ø A=ASC(TEMP$)
70Ø POKE F+LOOP-1,A
71Ø NEXT LOOP
72Ø :
73Ø FOR LOOP=1 TO LEN(B$)
74Ø TEMP$=MID$(B$,LOOP,1)
75Ø B=ASC(TEMP$)
76Ø POKE S+LOOP-1,B
77Ø NEXT LOOP
78Ø :
79Ø POKE 251,Ø POKE 252,194
80Ø POKE 253,Ø : POKE 254,195
81Ø POKE 82Ø,LEN(A$)

```

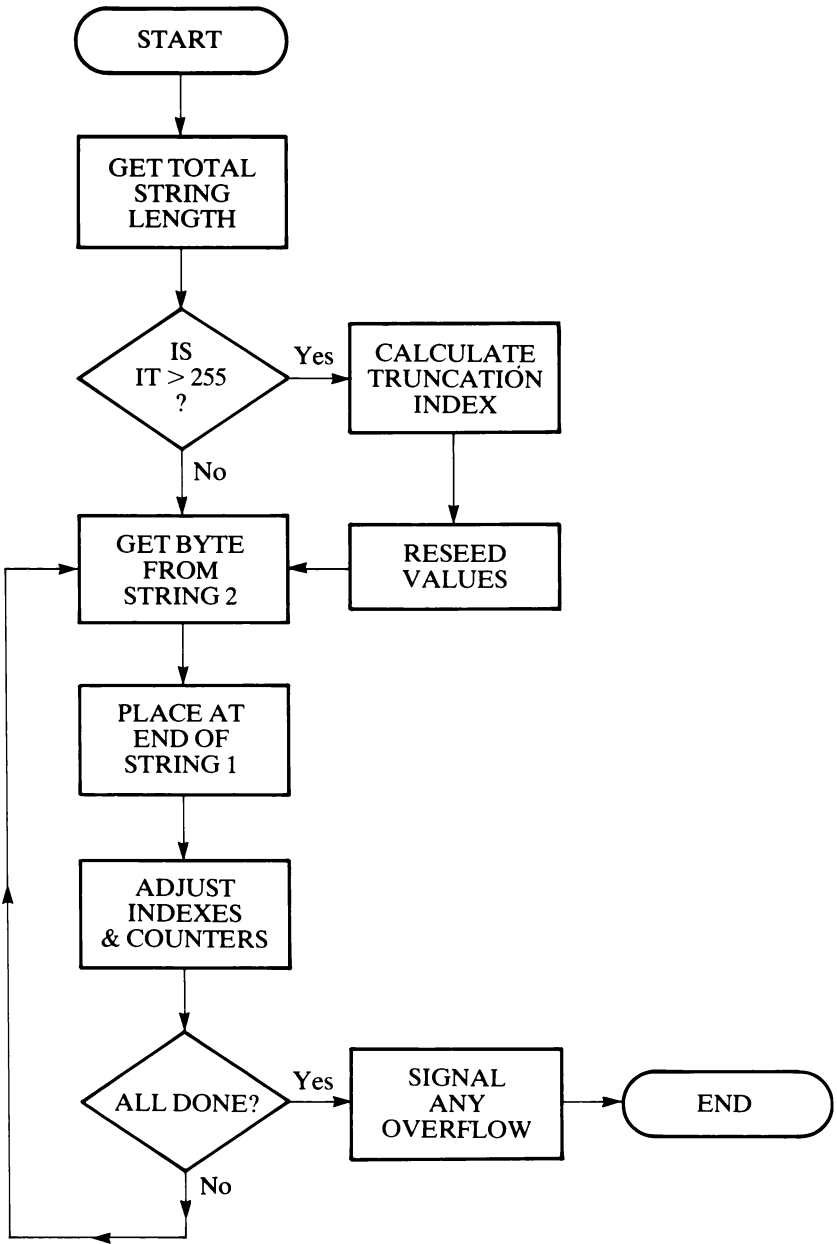


Figure 5.2 Concatenate strings flowchart

82Ø POKE 821, LEN(B\$)

83Ø :

84Ø SYS CODE

85Ø :

```

86Ø REM *** PRINT OUT FINAL STRING ***
87Ø PRINT "FINAL STRING IS :";
88Ø LOOP=Ø
89Ø REM ** REPEAT **
90Ø BYTE=PEEK(F+LOOP)
91Ø PRINT CHR$(BYTE);
92Ø LOOP=LOOP+1
93Ø IF BYTE=13 THEN END
94Ø GOTO 90Ø

```

This program allows a final string of 256 characters in length to be manipulated. Therefore, as the program stands, the combined lengths of the two strings should not exceed this length. If they do, then only as many characters as space allows will be concatenated on to the first string, leaving the second string truncated. The Carry flag is used to signal whether any truncation has taken place, being set if it has and cleared otherwise. As with the string comparison routine, the string buffers are accessed via two zero page vectors (lines 790 and 800) and two bytes are reserved to hold the length of each string. A further two bytes are used to save index values.

The first nine machine code operations (lines 100 to 180) determine the final length of the string, by adding the length of the first string to that of the second string. A sum greater than 256 is signalled in the Carry flag and the branch of line 170 is performed, in which case the number of characters which can be inserted into the first string buffer is ascertained. The overflow indicator is loaded with \$FF if a truncation occurs; otherwise it is cleared with \$00.

The concatenating loop is held between lines 350 and 460. This simply moves a byte from the vectored address plus the index of the second string and places it at the end of the first string, as pointed to by the first string index byte. This process is reiterated until the value of 'count' has reached zero. Lines 480 and 500 place a RETURN character at the end of the string to facilitate printing from BASIC or machine code. The Overflow flag is loaded into the accumulator and bit 7 rotated across into the Carry flag, thereby signalling whether truncation has occurred. Lines 610 to 770 hold the BASIC test routine that reads in and then pokes the character strings into memory at \$C200 and \$C300. After the SYS call (line 840), the final BASIC routine prints the concatenated string from memory.

## Project

Adapt the program to perform the BASIC equivalent of  $C\$ = A\$ + B\$$  or  $C\$ = B\$ + A\$$  on request.



## Line-by-line

A line-by-line description of Program 9 now follows:

line 100 : get first string's length  
line 110 : string one's index  
line 120 : clear accumulator  
line 130 : set string two's index to zero  
line 140 : clear Carry flag  
line 150 : get second string's length  
line 160 : and add to length of first string  
line 170 : branch to TOO-LONG if total greater than 256 bytes  
line 180 : otherwise jump to GOOD-LENGTH  
line 190 : entry for TOO-LONG  
line 200 : load accumulator with 255  
line 210 : and store to indicate overflow  
line 220 : set Carry flag and subtract  
line 230 : string one's length from maximum length  
line 240 : branch to FINISH if first string is greater than  
256 bytes in length  
line 250 : save current count  
line 260 : restore maximum length  
line 270 : store in string one's length  
line 280 : jump to concatenation routine  
line 290 : entry for GOOD-LENGTH  
line 300 : save accumulator in string one's length  
line 310 : load with 0 to clear  
line 320 : overflow indicator  
line 330 : get string two's length  
line 340 : save in count  
line 350 : entry for CONCATENATION  
line 360 : get count value  
line 370 : if zero, then finish  
line 380 : entry for LOOP  
line 390 : get index for string two  
line 400 : and get character from second string  
line 410 : get string one's index  
line 420 : and place character into first string  
line 430 : increment first string's index  
line 440 : increment second string's index  
line 450 : decrement count

```

line 46Ø : branch to LOOP until count=Ø
line 47Ø : entry for FINISHED
line 48Ø : get final length of first string
line 49Ø : load accumulator with ASCII return
line 50ØØ : place at end of string
line 51Ø : get overflow indicator
line 52Ø : and move it into Carry flag
line 53Ø : back to calling routine

```

## COPY-CAT

String manipulation routines must include a method of copying substrings of characters from anywhere within a string of characters. In BASIC, three such commands are provided. They are MID\$, LEFT\$ and RIGHT\$, although with the first of these, any point in a string can be accessed. The following shows the sort of thing possible in BASIC:

```

1ØØ A$="CONCATENATE"
11Ø B$=MID$(A$,Ø,3)
12Ø PRINT B$

```

Running this will output the string 'CON'. What the code has done is to take the three characters from the first character in the Main\$. Program 10 produces the same type of operation from machine code.

### Program 10

```

1Ø REM ** COPY A SUBSTRING FROM WITHIN **
2Ø REM ** A MAIN ASCII STRING **
3Ø CODE=49152
4Ø MAIN=5Ø432      : REM $C5ØØ
5Ø SUB=5Ø688      : REM $C6ØØ
6Ø REM ** READ AND POKE M/C DATA **
7Ø FOR LOOP=Ø TO 123
8Ø   READ BYTE
9Ø   POKE CODE+LOOP,BYTE
1ØØ NEXT LOOP
11Ø :
12Ø REM ** M/C DATA **
13Ø DATA 16Ø,Ø      : REM LDY #$ØØ
14Ø DATA 14Ø,52,3  : REM STY $334

```

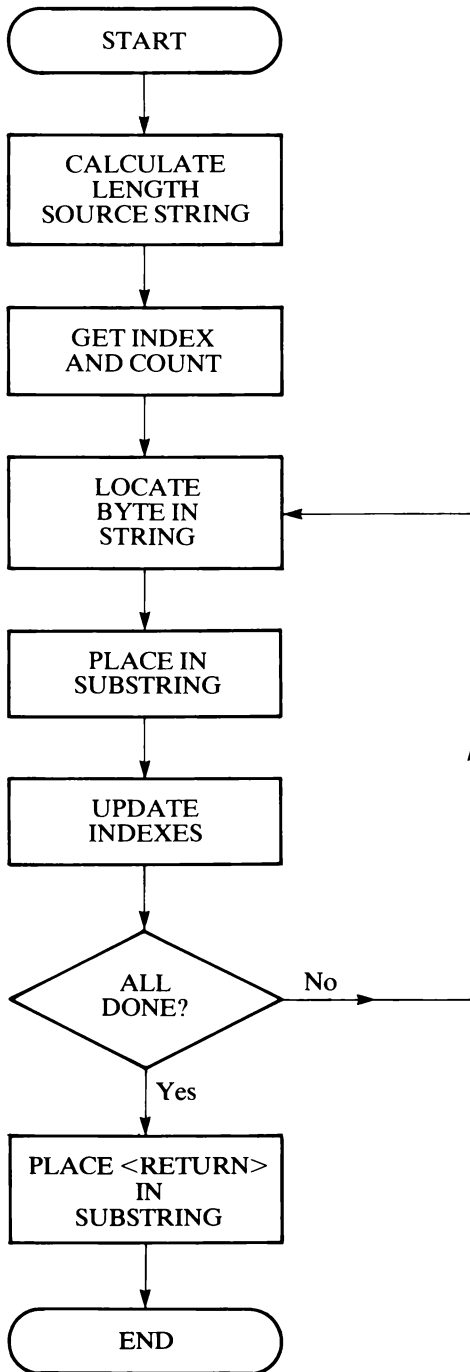


Figure 5.3 Copy string flowchart

```

15Ø DATA 14Ø,56,3 : REM STY $338
16Ø DATA 173,54,3 : REM LDA $336
17Ø DATA 24Ø,98 : REM BEQ $62
  
```

```

180 DATA 173,53,3 : REM LDA $335
190 DATA 205,55,3 : REM CMP $337
200 DATA 144,93 : REM BCC $5D
210 DATA 24 : REM CLC
220 DATA 173,55,3 : REM LDA $337
230 DATA 109,54,3 : REM ADC $336
240 DATA 176,9 : REM BCS $09
250 DATA 170 : REM TAX
260 DATA 202 : REM DEX
270 DATA 236,53,3 : REM CPX $335
280 DATA 144,20 : REM BCC $14
290 DATA 240,18 : REM BEQ $12
300 :: REM TRUNCATION
310 DATA 56 : REM SEC
320 DATA 173,53,3 : REM LDA $335
330 DATA 237,55,3 : REM SBC $337
340 DATA 141,54,3 : REM STA $336
350 DATA 238,54,3 : REM INC $336
360 DATA 169,255 : REM LDA #$FF
370 DATA 141,56,3 : REM STA $338
380 :: REM GREATER-EQUAL
390 DATA 173,54,3 : REM LDA $336
400 DATA 201,255 : REM CMP #$FF
410 DATA 144,10 : REM BCC $0A
420 DATA 240,8 : REM BEQ $08
430 DATA 169,255 : REM LDA #$FF
440 DATA 141,54,3 : REM STA $336
450 DATA 141,56,3 : REM STA $338
460 :: REM COPY-SUBSTRING
470 DATA 174,54,3 : REM LDX $336
480 DATA 240,35 : REM BEQ $23
490 DATA 169,0 : REM LDA #$00
500 DATA 141,52,3 : REM STA $334
510 :: REM LOOP
520 DATA 172,55,3 : REM LDY $337
530 DATA 177,251 : REM LDA ($FB),Y
540 DATA 172,52,3 : REM LDY $334
550 DATA 145,253 : REM STA ($FD),Y
560 DATA 238,55,3 : REM INC $337
570 DATA 238,52,3 : REM INC $334

```

```

58Ø DATA 2Ø2      : REM  DEX
59Ø DATA 2Ø8,237  : REM  BNE $ØD
6ØØ DATA 2Ø6,52,3 : REM  DEC $334
61Ø DATA 173,56,3 : REM  LDA $338
62Ø DATA 2Ø8,3    : REM  BNE $Ø3
63Ø ::            REM  FINISH
64Ø DATA 24       : REM  CLC
65Ø DATA 144,1    : REM  BCC $Ø1
655 ::            REM  ERROR
66Ø DATA 56       : REM  SEC
67Ø ::            REM  OUT
68Ø DATA 169,13   : REM  LDA #$$ØD
69Ø DATA 172,52,3 : REM  LDY $334
7ØØ DATA 2ØØ      : REM  INY
71Ø DATA 145,253  : REM  STA ($FD),Y
72Ø DATA 96       : REM  RTS
73Ø :
74Ø REM ** SET UP MAIN STRING **
75Ø PRINT CHR$(147)
76Ø ::            REM  ERROR
77Ø INPUT "MAIN STRING ";B$
78Ø FOR LOOP=1 TO LEN(B$)
79Ø   TEMP$=MID$(B$,LOOP,1)
8ØØ   B=ASC(TEMP$)
81Ø   POKE MAIN+LOOP-1,B
82Ø NEXT LOOP
83Ø :
84Ø INPUT"INDEX INTO STRING ";X
85Ø INPUT"NUMBER OF BYTES TO COPY ";Y
86Ø :
87Ø REM ** SET UP BYTES FOR M/C **
88Ø POKE 251,Ø : POKE 252,197
                        : REM $C5ØØ VECTOR
89Ø POKE 253,Ø : POKE 254,198
                        : REM $C6ØØ VECTOR
9ØØ POKE 821,LEN(B$)
91Ø POKE 822,Y
92Ø POKE 823,X
93Ø :
94Ø SYS CODE

```

```

95Ø :
96Ø REM ** READ COPIED SUBSTRING **
97Ø FOR LOOP=1 TO Y
98Ø Z=PEEK(SUB+LOOP-1)
99Ø PRINT CHR$(Z);
1ØØØ NEXT LOOP

```

Bytes are designated as follows:

251–252	(\$FB–\$FC)	:	main string vector
253–254	(\$FD–\$FE)	:	substring vector
82Ø	(\$334)	:	length of substring
821	(\$335)	:	length of main string
822	(\$336)	:	number of bytes to be copied
823	(\$337)	:	index into main string
824	(\$338)	:	error flag

Once again, a few lines of BASIC demonstrate the operation of the routine, requesting the source string, starting index and length of substring, or rather the number of bytes to be copied into the substring from the starting index. The main string is in a buffer located at \$C5ØØ and the substring is copied into its own buffer at \$C6ØØ. As always, these addresses may be changed to suit user needs, as they are vectored through zero page (lines 88Ø and 89Ø).

Error-checking is allowed, as the Carry flag is set on exit if an error has occurred. Normally, an error will occur only if the starting index is beyond the length of the source string, or the number of bytes to be copied from the main string is zero. If the number of bytes requested in the length exceeds the number left from the indexed position to the end of the main string, then only the bytes available will be copied to the substring buffer.

On entry to the routine, error-checking is performed (lines 16Ø to 24Ø) and if any are found, the program exits. Lines 3ØØ to 37Ø perform a truncation if the number of bytes to be copied exceeds those available. The COPY-SUBSTRING loop (lines 46Ø to 59Ø) copies each string byte from the vectored address in the main string to the substring buffer. Each time a character is copied, the substring length byte is incremented. On completion of this loop, controlled by the X register, the error flag is restored and the Carry flag conditioned accordingly (lines 61Ø to 66Ø). Finally (lines 69Ø to 73Ø), an ASCII RETURN character is placed at the end of the substring.

The following example shows the resultant substrings produced from the main string 'CONCATENATE' for different indexes. Figure 5.4 illustrates the index value for each of the main string's characters.

Index	Length	Substring
∅	3	CON
3	3	CAT
4	3	ATE

String	C	O	N	C	A	T	E	N	A	T	E
Index	∅	1	2	3	4	5	6	7	8	9	1∅

Figure 5.4 String Index

### Line-by-line

A line-by-line description of Program 10 follows:

line 13∅ : initialize Y register  
line 14∅ : clear substring length  
line 15∅ : and error flag  
line 16∅ : get substring length  
line 17∅ : if a null string, branch to FINISH  
line 18∅ : get main string's length  
line 19∅ : compare it with index byte  
line 20∅ : branch to ERROR if index is greater  
line 21∅ : clear the Carry flag  
line 22∅ : get index  
line 23∅ : add it to substring length  
line 24∅ : branch to TRUNCATION if result is greater than 255  
line 25∅ : move index across into X register  
line 26∅ : decrement it by one  
line 27∅ : compare result with string length  
line 28∅ : branch to GREATER-EQUAL if result is  
line 29∅ : greater than or equal to string length  
line 30∅ : entry for TRUNCATION  
line 31∅ : set the Carry flag  
line 32∅ : get string length  
line 33∅ : subtract the index from it  
line 34∅ : save the new length  
line 35∅ : and increment it by one

line 360 : denote an error by  
line 370 : setting the error flag  
line 380 : entry for GREATER-EQUAL  
line 390 : get length into accumulator  
line 400 : compare with maximum length  
line 410 : branch if count is  
line 420 : greater or equal to maximum length  
line 430 : put maximum length in accumulator  
line 440 : store in bytes to copy  
line 450 : and also in error flag  
line 460 : entry for COPY-SUBSTRING  
line 470 : get the index position  
line 480 : branch to ERROR if zero  
line 490 : clear accumulator  
line 500 : and substring length  
line 510 : entry for LOOP  
line 520 : get main string index into Y register  
line 530 : get character from main string  
line 540 : get substring index  
line 550 : copy character into substring  
line 560 : increment main string index  
line 570 : increment substring index  
line 580 : decrement bytes to move counter  
line 590 : branch to LOOP if still bytes to be copied  
line 600 : decrement final substring count  
line 610 : get error flag into accumulator  
line 620 : branch to ERROR if not zero  
line 630 : FINISH entry  
line 640 : clear Carry flag as no error  
line 650 : branch to OUT  
line 655 : entry for ERROR  
line 660 : set Carry flag to indicate error  
line 670 : entry for OUT  
line 680 : place RETURN in accumulator  
line 690 : get substring index into Y  
line 700 : increment Y  
line 710 : place RETURN at end of substring  
line 720 : return to BASIC.



## INSERTION

This final routine provides the facility for inserting a string within the body of another string, allowing textual material—for example, in word processing applications—to be manipulated. If the main string held 'ELIZABETH OKAY', this routine could be called to insert the string 'RULES', so that the final string would read 'ELIZABETH RULES OKAY'. As with the COPY routine, the position of the insertion is pointed to by an index byte, and the Carry flag is set if an error is detected—that is, if an index of 0 or a null substring is specified.

The maximum length of the final string is 256 characters. If the insertion of the substring would cause this length to be exceeded, the substring is truncated to the length given by (256 minus length of main string) and only these characters are inserted.

As always, a BASIC primer demonstrates the routine's use. The string buffers are held at \$C500 and \$C600 and in this instance they are accessed directly, although there is no reason why vectored addresses could not be used.

### Program 11

```
10 REM ** INSERT ONE ASCII STRING **
20 REM ** INTO ANOTHER ASCII STRING **
30 MAIN=50432          : REM $C500
40 SUB=50688           : REM $C600
50 CODE=49152
60 REM ** READ AND POKE DATA **
70 FOR LOOP=0 TO 141
80  READ BYTE
90  POKE LOOP+CODE,BYTE
100 NEXT LOOP
110 :
120 REM ** M/C DATA **
130 DATA 160,0        : REM LDY #0
140 DATA 140,53,3    : REM STY $335
150 DATA 165,252     : REM LDA $FC
160 DATA 208,3       : REM BNE $03
170 DATA 76,137,192  : REM JMP $C089
180 ::               : REM ZERO-LENGTH
190 DATA 165,253     : REM LDA $FD
200 DATA 240,124     : REM BEQ $7C
210 DATA ::         : REM CHECK
```

```

22Ø DATA 24 : REM CLC
23Ø DATA 165,252 : REM LDA $FC
24Ø DATA 1Ø1,251 : REM ADC $FB
25Ø DATA 176,6 : REM BCS $Ø6
26Ø DATA 2Ø1,255 : REM CMP #FFF
27Ø DATA 24Ø,18 : REM BEQ $12
28Ø DATA 144,16 : REM BCC $1Ø
29Ø :: REM CUT-OFF
3ØØ DATA 169,255 : REM LDA #FFF
31Ø DATA 56 : REM SEC
32Ø DATA 229,251 : REM SBC $FB
33Ø DATA 24Ø,1Ø4 : REM BEQ $68
34Ø DATA 144,1Ø2 : REM BCC $66
35Ø DATA 133,252 : REM STA $FC
36Ø DATA 169,255 : REM LDA #FFF
37Ø DATA 141,53,3 : REM STA $335
38Ø :: REM CALC-LENGTH
39Ø DATA 165,251 : REM LDA $FB
4ØØ DATA 197,253 : REM CMP $FD
41Ø DATA 176,2Ø : REM BCS $14
42Ø DATA 166,251 : REM LDX $FB
43Ø DATA 232 : REM INX
44Ø DATA 134,253 : REM STX $FD
45Ø DATA 169,255 : REM LDA #FFF
46Ø DATA 141,53,3 : REM STA $335
47Ø DATA 24 : REM CLC
48Ø DATA 165,251 : REM LDA $FB
49Ø DATA 1Ø1,252 : REM ADC $FC
5ØØ DATA 133,251 : REM STA $FB
51Ø DATA 76,1Ø9,192 : REM JMP $CØ6D
52Ø :: REM NO-PROBLEMS
53Ø DATA 56 : REM SEC
54Ø DATA 165,251 : REM LDA $FB
55Ø DATA 229,253 : REM SBC $FD
56Ø DATA 17Ø : REM TAX
57Ø DATA 232 : REM INX
58Ø DATA 165,251 : REM LDA $FB
59Ø DATA 133,254 : REM STA $FE
6ØØ DATA 24 : REM CLC
61Ø DATA 1Ø1,252 : REM ADC $FB

```

```

62Ø DATA 133,251 : REM STA $FB
63Ø DATA 141,52,3 : REM STA $334
64Ø :: REM MAKE-SPACE
65Ø DATA 164,254 : REM LDY $FE
66Ø DATA 185,Ø,197 : REM LDA $C5ØØ,Y
67Ø DATA 172,52,3 : REM LDY $334
68Ø DATA 153,Ø,197 : REM STA $C5ØØ,Y
69Ø DATA 2Ø6,52,3 : REM DEC $334
7ØØ DATA 198,254 : REM DEC $FE
71Ø DATA 2Ø2 : REM DEX
72Ø DATA 2Ø8,237 : REM BNE $ED
73Ø :: REM INSERT-SUBSTRING
74Ø DATA 169,Ø : REM LDA #$ØØ
75Ø DATA 133,254 : REM STA $FE
76Ø DATA 166,252 : REM LDX $FC
77Ø :: REM TRANSFER
78Ø DATA 164,254 : REM LDY $FE
79Ø DATA 185,Ø,198 : REM LDA $C6ØØ,Y
8ØØ DATA 164,253 : REM LDY $FE
81Ø DATA 153,Ø,197 : REM STA $C5ØØ,Y
82Ø DATA 23Ø,253 : REM INC $FD
83Ø DATA 23Ø,254 : REM INC $FE
84Ø DATA 2Ø2 : REM DEX
85Ø DATA 2Ø8,239 : REM BNE $EF
86Ø DATA 173,53,3 : REM LDA $335
87Ø DATA 2Ø8,3 : REM BNE $Ø3
88Ø :: REM GOOD
89Ø DATA 24 : REM CLC
9ØØ DATA 144,1 : REM BCC $Ø1
91Ø :: REM ERROR
92Ø DATA 56 : REM SEC
93Ø :: REM FINISH
94Ø DATA 96 : REM RTS
95Ø :
96Ø REM ** GET MAIN STRING AND STORE AT
    $C5ØØ **
97Ø PRINT CHR$(147)
98Ø INPUT"MAIN STRING";B$
99Ø FOR LOOP=1 TO LEN(B$)
1ØØØ TEMP$=MID$(B$,LOOP,1)

```

```

1010 B=ASC(TEMP$)
1020 POKE MAIN+LOOP-1,B
1030 NEXT LOOP
1040 :
1050 REM ** GET SUBSTRING AND STORE AT $C600 **
1060 INPUT"SUB STRING";C$
1070 FOR LOOP=1 TO LEN(C$)
1080 TEMP$=MID$(C$,LOOP,1)
1090 B=ASC(TEMP$)
1100 POKE SUB+LOOP-1,B
1110 NEXT LOOP
1120 :
1130 REM ** GET INSERTION INDEX **
1140 INPUT"INSERTION INDEX"; X
1150 :
1160 REM ** POKE VALUES INTO ZERO PAGE **
1170 POKE 251,LEN(B$)
1180 POKE 252,LEN(C$)
1190 POKE 253,X
1200 :
1210 SYS CODE
1220 :
1230 REM ** READ FINAL STRING **
1240 COUNT=LEN(B$)+LEN(C$)-1
1250 FOR LOOP=0 TO COUNT
1260 Z=PEEK(MAIN+LOOP)
1270 PRINT CHR$(Z);
1280 NEXT LOOP

```

The program begins by checking the length bytes to ensure that no null strings are present (lines 150 to 200) and then sums the two lengths to obtain the final length. If the addition results in the Carry flag being set (line 250), the total length will exceed 256 bytes and, as a result, the inserted substring will be truncated (lines 310 to 390).

If the insertion index is greater than the length of the string, the substring is actually concatenated on to the end of the main string. This evaluation is performed through lines 400 to 530. Before inserting the substring, all characters to the left of the index must be shuffled up through memory to make space for it. These calculations are carried out in lines 550 to 650, ready for the shuffling process (lines 660 to 740). Inserting the substring now involves simply copying it from its buffer into the space opened up for it

(lines 750 to 870), the X register being used as the characters-moved counter.

Finally, the error flag is restored and the Carry flag conditioned to signal any errors.

### Line-by-line

A line-by-line description of Program 11 follows:

line 130 : clear indexing register  
line 140 : clear error flag  
line 150 : get substring length  
line 160 : branch to ZERO-LENGTH if Z=0  
line 170 : otherwise carry on  
line 180 : entry for ZERO-LENGTH  
line 190 : get offset  
line 200 : branch to ERROR if Z=1  
line 210 : entry for CHECK  
line 220 : clear Carry flag  
line 230 : get substring length  
line 240 : add it to main string length  
line 250 : branch to CUT-OFF if greater than 256  
line 260 : is it maximum length?  
line 270 : branch to CALC-LENGTH if  
line 280 : it is equal to or greater than  
line 290 : entry for CUT-OFF  
line 300 : get the maximum length allowed  
line 310 : set Carry flag  
line 320 : subtract length of string  
line 330 : branch to ERROR if  
line 340 : length is equal to or greater than string  
line 350 : save characters free  
line 360 : set error flag  
line 380 : entry for CALC-LENGTH  
line 390 : get main string length  
line 400 : is offset within string?  
line 410 : branch to NO-PROBLEMS if it is  
line 420 : else place substring  
line 430 : at end of main string  
line 440 : save X in offset  
line 450 : and flag the error

line 460 : in error flag byte  
line 470 : clear Carry flag  
line 480 : get length of string  
line 490 : calculate total length  
line 500 : and save result  
line 510 : jump to INSERT-SUBSTRING  
line 520 : entry for NO-PROBLEMS  
line 530 : set Carry flag  
line 540 : get length of substring  
line 550 : subtract offset  
line 560 : move index into X  
line 570 : increment index  
line 580 : get length  
line 590 : save in source  
line 600 : clear Carry flag  
line 610 : find total length  
line 620 : save result  
line 630 : and for index  
line 640 : entry for MAKE-SPACE  
line 650 : get source index  
line 660 : get byte from main  
line 670 : get offset into string  
line 680 : move byte along  
line 690 : decrement both indexes  
line 710 : decrement counter  
line 720 : branch to MAKE-SPACE until done  
line 730 : entry for INSERT-SUBSTRING  
line 740 : clear accumulator  
line 750 : and source  
line 760 : get counter  
line 770 : entry for TRANSFER  
line 780 : get index  
line 790 : get byte from substring  
line 800 : get offset into main string  
line 810 : and place byte in main  
line 820 : increment both indexes  
line 840 : do until substring inserted  
line 850 : branch to TRANSFER  
line 860 : get error flag  
line 870 : branch to ERROR

line 88∅ : entry for GOOD  
line 89∅ : signal no error  
line 90∅ : branch to FINISH  
line 91∅ : entry for ERROR  
line 92∅ : denote error  
line 93∅ : entry for FINISH  
line 94∅ : return to calling routine

# 6 Printing Print!

Every machine code program sooner or later requires text to be printed on to the screen. In most instances, this is a fairly simple process and often involves merely indexing into an ASCII string table and printing the characters, using one of the Operating System calls, until either a RETURN character or zero byte is encountered. Program 12 uses this method.

## Program 12

```
10 REM ** PRINT STRING FROM MEMORY **
20 CODE=49152
30 FOR LOOP=0 TO 13
40 READ BYTE
50 POKE CODE+LOOP,BYTE
60 NEXT LOOP
70 :
80 REM ** M/C DATA **
90 :: REM STRING-PRINT
100 DATA 162,0 : REM LDX #00
110 :: REM NEXT-CHARACTER
120 DATA 189,0,197 : REM LDA $C500,X
130 DATA 32,210,255 : REM JSR $FFD2
140 DATA 232 : REM INX
150 DATA 201,13 : REM CMP #0D
160 DATA 208,245 : REM BNE $F5
170 DATA 96 : REM RTS
180 :
190 REM ** GET STRING TO BE PRINTED **
200 STRING=50432
```



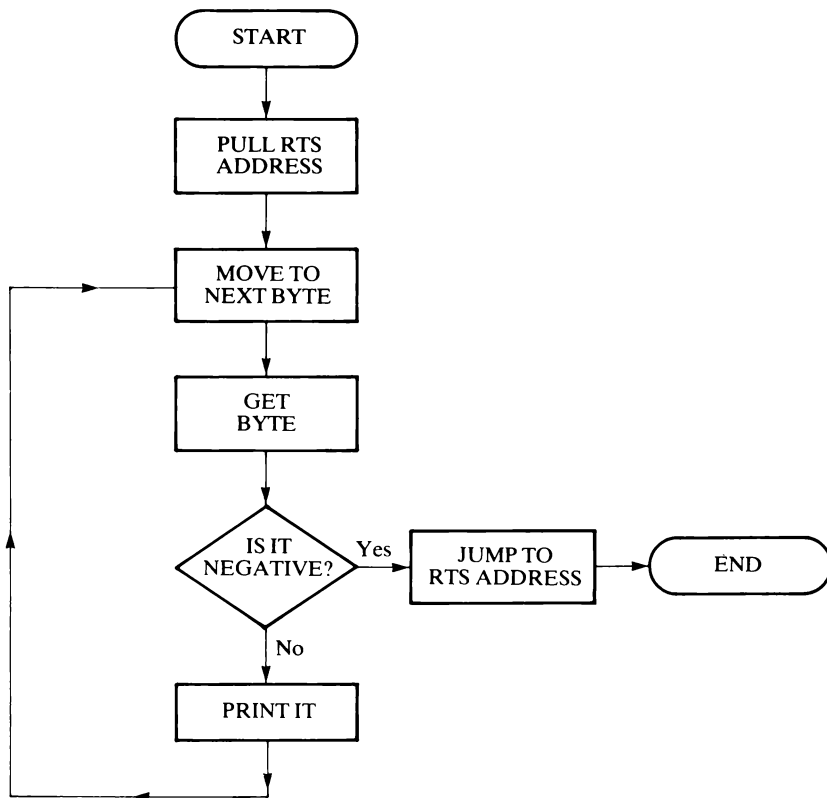


Figure 6.1 Printing embedded code flowchart

```

21Ø PRINT CHR$(147)
22Ø INPUT "INPUT STRING :";A$
23Ø FOR LOOP=1 TO LEN(A$)
24Ø TEMP$=MID$(A$,LOOP,1)
25Ø B=ASC(TEMP$)
26Ø POKE STRING+LOOP-1,B
27Ø NEXT LOOP
28Ø PRINT:PRINT
29Ø PRINT"YOUR STRING WAS AS FOLLOWS :";
30Ø SYS CODE
  
```

Here, a string buffer is located at \$C500 (50432) and the requirement for printing the string is that it must be terminated with an ASCII RETURN character, \$0D. The program begins by initializing an index, the X register (line 100), and loading the byte at \$C500+X into the accumulator. This is printed using the Kernal's CHROUT routine, the index is incremented and then the accumulator's contents are compared to see whether the character just output was a RETURN (line 150). If not, the loop branches back and the next character is sought.

Program 13 shows how several strings may be printed to the screen using a loop similar to that described above. The number of strings for printing may be variable, the desired number being passed into the routine via the Y register. The string data has been entered using the DATA statement. If a large amount of string data is to be stored, and the amount to be printed at any one time varied, a vectored address should be used to access the table. Positioning of the text on the screen can be performed by embedding the relative number of RETURNS and spaces into the DATA, or more neatly by using the Kernal's PLOT routine to set the X and Y tab co-ordinates.

### Program 13

```

1Ø REM ** PRINT Y NUMBER OF STRINGS **
2Ø CODE=49152
3Ø FOR LOOP=Ø TO 18
4Ø READ BYTE
5Ø POKE CODE+LOOP,BYTE
6Ø NEXT LOOP
7Ø :
8Ø REM ** M/C DATA **
9Ø DATA 162,Ø : REM LDX #ØØØ
1ØØ DATA 16Ø,4 : REM LDY #Ø4
11Ø :: REM NEXT-CHARACTER
12Ø DATA 189,Ø,197 : REM LDA $C5ØØ,X
13Ø DATA 32,21Ø,255 : REM JSR $FFD2
14Ø DATA 232 : REM INX
15Ø DATA 2Ø1,13 : REM CMP #ØD
16Ø DATA 2Ø8,245 : REM BNE $F5
17Ø DATA 136 : REM DEY
18Ø DATA 2Ø8,242 : REM BNE $F2
19Ø DATA 96 : REM RTS
2ØØ :
21Ø REM ** SET UP FOUR SIMPLE STRINGS **
22Ø STRING=5Ø432
23Ø FOR LOOP=Ø TO 31
24Ø READ BYTE
25Ø POKE STRING+LOOP,BYTE
26Ø NEXT
27Ø :
28Ø REM ** ASCII DATA **

```

```

290 DATA 32,65,65,65,65,65,65,13
300 DATA 32,32,66,66,66,66,66,13
310 DATA 32,32,32,67,67,67,67,13
320 DATA 32,32,32,32,68,68,68,13

```

The final program in this chapter shows the way I find easiest to store and print character strings, stowing them directly within the machine code. The two main advantages of this method are that the string is inserted directly at the point it is needed, avoiding the need to calculate indexes into look-up tables, and that because it manipulates its own address it is fully relocatable.

#### Program 14

```

10 REM ** ASCII STRING OUTPUT ROUTINE **
20 CODE=49152
30 FOR LOOP=0 TO 26
40 READ BYTE
50 POKE CODE+LOOP,BYTE
60 NEXT LOOP
70 :
80 REM ** M/C DATA **
90 DATA 104 : REM PLA
100 DATA 133,251 : REM STA $FB
110 DATA 104 : REM PLA
120 DATA 133,252 : REM STA $FC
130 :: REM REPEAT
140 DATA 160,0 : REM LDY #$0
150 DATA 230,251 : REM INC $FB
160 DATA 208,2 : REM BNE $02
170 DATA 230,252 : REM INC $FC
180 :: REM OVER
190 DATA 177,251 : REM LDA ($FB),Y
200 DATA 48,6 : REM BMI $06
210 DATA 32,210,255 : REM JSR $FFD2
220 DATA 76,6,192 : REM JMP $C006
230 :: REM FINISH
240 DATA 108,251,0 : REM JMP ($FB)
250 :
260 REM ** DEMO ROUTINE LOCATED AT $C200 **
270 DEMO=49664

```

```

280 FOR LOOP=0 TO 38
290 READ BYTE
300 POKE DEMO+LOOP,BYTE
310 NEXT LOOP
320 :
330 REM ** DEMO M/C DATA **
340 DATA 169,147 : REM LDA #93
350 DATA 32,210,255 : REM JSR $FFD2
360 DATA 32,0,192 : REM JSR $C000
370 REM ** NOW STORE ASCII CODES FOR PRINTING **
380 DATA 13 : REM CARRIAGE-RETURN
390 DATA 83,84,82,73,78,71,83,32
: REM STRINGS<SPACE>
400 DATA 87,73,84,72,73,78,32
: REM WITHIN<SPACE>
410 DATA 77,65,67,72,73,78,69,32
: REM MACHINE<SPACE>
420 DATA 67,79,68,69,33
: REM CODE!
430 DATA 234 : REM NOP
440 DATA 96 : REM RTS
450 :
460 SYS DEMO

```

The ASCII character string is placed in memory by leaving the machine code assembly (line 360) and POKEing the ASCII codes of the string directly into successive memory locations (lines 380 to 420).

For this routine to work, it is imperative that the first byte following the string is a negative byte—that is, one with bit 7 set. The opcode for NOP, \$EA, is ideal for this purpose as it has its most significant bit set (\$EA=11101010) and its only effect is to cause a very short delay.

The ASCII print routine is just 27 bytes in length and it should be called as a subroutine immediately before the string is encountered (line 360). On entry into the subroutine, the first four operations pull the return address from the stack and save it in a zero page vector at \$FB and \$FC. These bytes are then incremented by one to point at the byte following the subroutine call.

Because the string data follows on immediately after the ASCII print subroutine call, post-indexed indirect addressing can be used to load the first string character into the accumulator (line 190). The string terminating negative byte is tested for (line 200), and if not found the byte is printed with a CHROUT call. A JMP to

REPEAT is then performed and the loop reiterated. When the negative byte is encountered, and the branch of line 200 succeeds, an indirect jump (line 240) via the current vectored address is executed, returning control back to the calling machine code at the end of the ASCII string.

### Line-by-line

A line-by-line description of Program 14 follows:

```
line 90 : set low byte RTS address
line 100 : save in $FB
line 110 : get high byte RTS address
line 120 : save in $FC
line 130 : entry for REPEAT
line 140 : initialize index to zero
line 150 : increment low byte of vectored address
line 160 : branch to OVER if not zero
line 170 : else increment page value
line 180 : entry for OVER
line 190 : get byte from within program
line 200 : if negative, branch to FINISH
line 210 : else print it
line 220 : jump to REPEAT
line 230 : entry for FINISH
line 240 : jump back into main program
line 340 : load accumulator with clear screen code
line 350 : and print it
line 360 : call string printing routine at $C000
line 380 : ASCII code for RETURN
line 390 : ASCII string 'STRINGS '
line 400 : ASCII string 'WITHIN '
line 410 : ASCII string 'MACHINE '
line 420 : ASCII string 'CODE!'
line 430 : negative byte
line 440 : back to BASIC
```

# 7 A Bubble of Sorts

Any program written to handle quantities of data will, at some time, require the data in a data table to be sorted into ascending or descending order. Several algorithms are available to facilitate this manipulation of data, of which the bubble sort is perhaps the simplest to implement in BASIC or machine code.

The technique involves moving through the data list and comparing pairs of bytes. If the first byte is smaller than the next byte in the list, the next pair of bytes is sought. If, on the other hand, the second byte is less than the first, the two bytes are swapped. This procedure is repeated until a pass is executed in which no elements are exchanged, so all are in ascending order. Program 15 is the BASIC version of such a bubble sort.

## Program 15

```
10 REM ** BASIC BUBBLE SORT **
20 TABLE=828
30 FOR LOOP=0 TO 19
40 READ BYTE
50 POKE TABLE+LOOP,BYTE
60 NEXT LOOP
70 :
80 REM ** BUBBLE-UP ROUTINE **
90 FOR BUBBLE=0 TO 19
100 TEMP=BUBBLE
110 :
120 IF PEEK(TABLE+TEMP)>PEEK(TABLE+(TEMP-1))
    THEN GOTO 180
130 HOLD=PEEK(TABLE+TEMP)
140 POKE TABLE+TEMP,PEEK(TABLE+(TEMP-1))
```

```

15Ø POKE TABLE+(TEMP-1),HOLD
16Ø TEMP=TEMP-1
17Ø IF TEMP<>Ø THEN GOTO 12Ø
18Ø NEXT
19Ø :
20Ø REM ** DATA FOR SORTING **
21Ø DATA 1,255,67,89,12Ø
22Ø DATA 6,2ØØ,85,45,199
23Ø DATA Ø,123,77,98,231
24Ø DATA 9,234,99,98,1ØØ
25Ø :
26Ø REM ** PRINT SORTED DATA **
27Ø FOR LOOP=Ø TO 19
28Ø PRINT PEEK(TABLE+LOOP)
29Ø NEXT LOOP

```

The data bytes for sorting are held within the four data lines from 21Ø to 24Ø and these are read into a memory array called TABLE. The sorting procedure is performed through lines 9Ø to 18Ø, line 12Ø checking to see if a swap is required. If a swap is unnecessary, GOTO 18Ø is executed and the swap routine bypassed. If it is required, however, the GOTO statement is not encountered, and the swap is performed in lines 13Ø to 16Ø. The byte currently being pointed to is PEEKed into the variable HOLD (line 13Ø) and the next byte is PEEKed and then POKEd into the location immediately before it (line 14Ø). The swap is completed by POKing the value of HOLD into the now 'vacant' location. The variable TEMP is used to keep track of the number of passes through the loop.

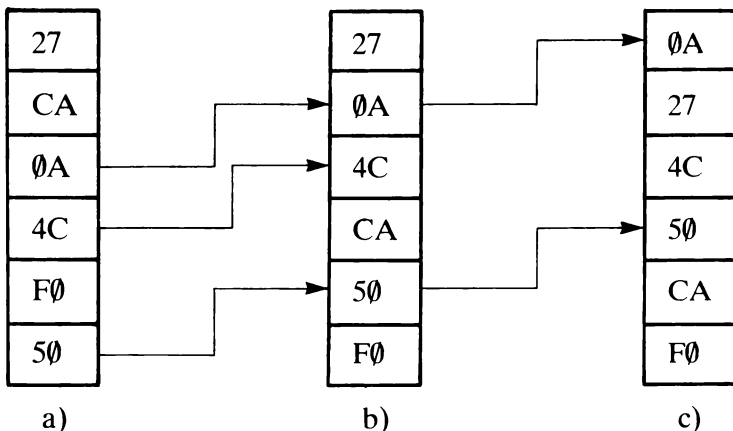


Figure 7.1 Numbers bubbling up

Figure 7.1 illustrates how small numbers bubble up through a data list using this sorting method. In this example, the data list consists of six numbers 27, CA, 0A, 4C, F0 and 50 (Figure 7.1a). After the first pass of the bubble sort three swaps have occurred (Figure 7.1b), thus:

1. 27 < CA therefore no change.
2. CA > 0A therefore swap items.
3. CA > 4C therefore swap items.
4. CA < F0 therefore no change.
5. F0 > 50 therefore swap items.

The next pass through the data list produces the ordered list of Figure 7.1c in which just two swaps occurred, as follows:

1. 27 > 0A therefore swap items.
2. 27 < 4C therefore no change.
3. 4C < 50 therefore no change.
4. CA > 50 therefore swap items.
5. CA < F0 therefore no change.

All the data elements are now in their final order, so the next pass through the list will have no effect. We can signal this by using an exchange flag to indicate whether the last pass produced any swaps, the sort routine exiting when the flag is cleared. This detail is included in the BASIC loader listed below as Program 16.

### Program 16

```

10 REM *** BUBBLE SORT ***
20 CODE=49152
30 TABLE=50432
40 FOR LOOP=0 TO 44
50 READ BYTE
60 POKE CODE+LOOP,BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 DATA 206,52,3 : REM DEC $334
110 :: REM BUBBLE-LOOP
120 DATA 160,0 : REM LDY #$00
130 DATA 140,53,3 : REM STY $335
140 DATA 174,52,3 : REM LDX $334
150 :: REM LOOP

```



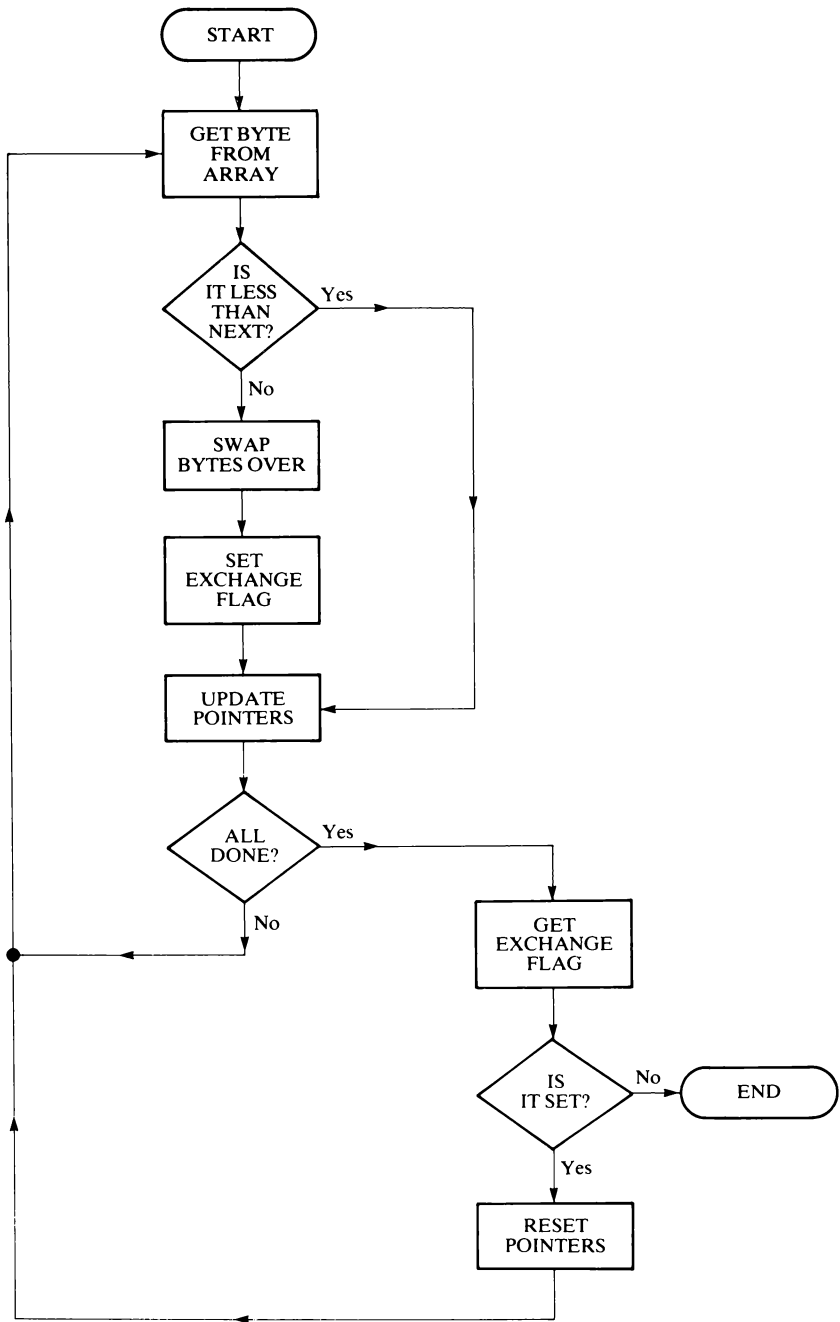


Figure 7.2 Bubble sort flowchart

```

160 DATA 177,253      : REM LDA ($FD),Y
170 DATA 209,251     : REM CMP ($FB),Y
180 DATA 176,13      : REM BCS $0D
190 DATA 72          : REM PHA
200 DATA 177,251     : REM LDA ($FB),Y
210 DATA 145,253     : REM STA ($FD),Y
220 DATA 104         : REM PLA
230 DATA 145,251     : REM STA ($FB),Y
240 DATA 169,1       : REM LDA #01
250 DATA 141,53,3    : REM STA $335
260 ::                REM SECOND-FIRST
270 DATA 200         : REM INY
280 DATA 202         : REM DEX
290 DATA 208,233     : REM BNE $E9
300 DATA 173,53,3    : REM LDA $335
310 DATA 240,5       : REM BEQ $05
320 DATA 206,52,3    : REM DEC $334
330 DATA 208,215     : REM BNE $D7
335 ::                REM FINISH
340 DATA 96          : REM RTS
350 :
360 REM ** SET UP VECTORS **
370 REM $FB=$C500, $FD=$C501
380 POKE 251,0 : POKE 252,197
390 POKE 253,1 : POKE 253,197
400 :
410 REM ** SET UP SCREEN AND ARRAY **
420 PRINT CHR$(147)
430 PRINT "**** MACHINE CODE BUBBLE SORT ****"
440 PRINT:PRINT
450 INPUT"NUMBER OF ELEMENTS IN ARRAY ";N
460 POKE 820,N        : REM LENGTH OF ARRAY
                        AT $334
470 FOR LOOP=0 TO N-1
480 PRINT"INPUT ELEMENT ";LOOP+1;
490 INPUT A
500 POKE TABLE+LOOP,A
510 NEXT LOOP
520 :
530 REM ** CALL CODE THEN PRINT SORTED TABLE **

```

```

54Ø SYS CODE
55Ø PRINT"SORTED VALUES ARE AS FOLLOWS"
56Ø FOR LOOP=Ø TO N-1
57Ø PRINT PEEK(TABLE+LOOP)
58Ø NEXT LOOP

```

After POKeIng the machine code data into memory at \$CØØØ, two zero page vectors are created to hold the address of the TABLE and TABLE+1 (lines 37Ø to 39Ø). The program then requests (in BASIC!) the number of elements in the array, which should be a series of integer values less than 256. These are then POKEd into memory (lines 45Ø to 51Ø). The machine code begins by decrementing the length of array byte by one (line 1ØØ), because the last element in the array will have no element beyond it to swap with. The swap flag is then cleared (line 13Ø) and the main loop entered using the X register to count the iterations.

The LOOP begins by loading the data byte into the accumulator (line 16Ø) and comparing it with the one immediately preceding it. If the byte+1 is greater than the byte, the Carry flag will be set and no swap required, in which case the branch to SECOND-FIRST is executed (line 18Ø).

If a swap is required, the second byte is saved, pushing it on to the hardware stack. The first byte is then transferred to the second byte's position (lines 2ØØ and 21Ø) and the accumulator is restored from the stack and transferred to the position of the first byte (lines 22Ø to 23Ø). To denote that a swap has occurred, the swap flag is set (lines 24Ø and 25Ø). The index and counters are then adjusted (lines 27Ø and 28Ø) and the loop continues until all the array elements have been compared. Upon completion of a full pass through the array, the swap flag is checked. If it is clear, no exchanges took place during the last pass, so the data list is now ordered and the sort finished (line 3ØØ and 31Ø). If the flag is set, the length of array byte is decremented and the procedure repeated once more (lines 32Ø and 33Ø). On return from the SYS call, the now ordered list is printed out to the screen.

### Line-by-line

A line-by-line description of Program 16 now follows:

```

line 1ØØ : subtract one from the length of the array
line 11Ø : entry for BUBBLE-LOOP
line 12Ø : initialize indexing register
line 13Ø : clear the swap flag
line 14Ø : get the array size into the X register to act as a loop
           counter

```

line 150 : entry for LOOP  
line 160 : get the byte at the byte+1 position  
line 170 : compare it with the previous byte  
line 180 : branch to SECOND-FIRST if the second byte  
          (byte+1) is larger than the first (byte)  
line 190 : save accumulator on hardware stack  
line 200 : get first byte at 'byte' position  
line 210 : place in current location (byte+1)  
line 220 : restore accumulator  
line 230 : and complete swap of bytes  
line 240 : load accumulator with 1  
line 250 : and set the swap flag to denote that a swap has been  
          performed  
line 260 : entry for SECOND-FIRST  
line 270 : move index on to next byte  
line 280 : decrement loop counter  
line 290 : branch to LOOP until done  
line 300 : get the swap flag into the accumulator  
line 310 : if clear, branch to FINISH  
line 320 : decrement outer counter  
line 330 : branch to BUBBLE-LOOP until all done  
line 335 : entry to FINISH  
line 340 : back to calling routine

## Projects

Rewrite the BASIC sections of the program to make it a complete machine code routine.

Adapt the sorting routine to handle 16-bit numbers.

## 8 Software Stack

One of the criticisms of the 6510 processor is that it has a very limited set of operation instructions—only 56, though addressing modes extend this to 152 functions. With some thought, however, it is possible to implement operations present on other processors, such as the Z80 or 6809, and build up a set of very useful sub-routines which can ultimately be strung together to perform quite sophisticated operations, as well as making the conversion of programs written for other processors much easier.

The routine described below mimics an instruction in the 6809 instruction set which allows the contents of up to eight registers to be pushed on to a stack in memory. This stack is often known as the user stack. I said ‘up to eight registers’, because the ones to be pushed can be selected, this being determined by the bit pattern of the byte after the user stack subroutine call. But more of that in a moment. First, which registers are we going to push? Obviously all the processor registers: the Program Counter, Status register, accumulator, and Index registers. The three remaining ones, we will implement as three two-byte ‘psuedo-registers’ from the user area of zero page. These are:

```
PR1  :  $80 and $81
PR2  :  $82 and $83
PR3  :  $84 and $84
```

This now enables us to save the contents of these locations when required.

As already stated, the byte after the user stack subroutine call determines by its bit pattern which registers are to be pushed, as follows:

```
bit 0 : pseudo-register 1
bit 1 : pseudo-register 2
bit 2 : pseudo-register 3
```

bit 3 : Y register  
 bit 4 : X register  
 bit 5 : accumulator  
 bit 6 : Status register  
 bit 7 : Program Counter

The rule here is that if the bit is set, the related register is pushed. Thus the instructions:

```
JSR USER-STACK
.BYTE $FF
```

would push all registers on to the user stack, the embedded byte being \$FF or 11111111. Alternatively, the coding:

```
JSR USER-STACK
.BYTE $1E
```

where \$1E = ~~000~~11110 would push only the accumulator, Status and Index registers. Perhaps at this point a question is running through your mind: 'won't the embedded byte cause my program to crash?'. That's true on face value, but what we do is get the user stack coding to move the Program Counter on one byte, to pass over it, as Program 17 shows:

### Program 17

```

10 REM ** USER STACK **
20 CODE=49152
30 FOR LOOP=0 TO 116
40 READ BYTE
50 POKE CODE+LOOP,BYTE
60 NEXT LOOP
70 :
80 REM ** M/C DATA **
90 DATA 8 : REM PHP
100 DATA 72 : REM PHA
110 DATA 138,72 : REM TXA : PHA
120 DATA 152,72 : REM TYA : PHA
130 DATA 186 : REM TSX
140 DATA 160,6 : REM LDY #06
150 :: REM PUSH-ZERO-PAGE
160 DATA 185,138,0 : REM LDA 008A,Y
  
```

```

17Ø DATA 72 : REM PHA
18Ø DATA 136 : REM DEY
19Ø DATA 2Ø8,249 : REM BNE $F9
2ØØ DATA 254,5,1 : REM INC $1Ø5,X
21Ø DATA 189,5,1 : REM LDA $1Ø5,X
22Ø DATA 133,139 : REM STA $8B
23Ø DATA 2Ø8,3 : REM BNE $Ø3
24Ø DATA 254,6,1 : REM INC $1Ø6,X
25Ø :: REM PC-LOW
26Ø DATA 189,6,1 : REM LDA $1Ø6,X
27Ø DATA 133,14Ø : REM STA $8C
28Ø DATA 169,135 : REM LDA #87
29Ø DATA 133,141 : REM STA $8D
3ØØ DATA 177,139 : REM LDA ($8B),Y
31Ø DATA 133,142 : REM STA $8E
32Ø DATA 169,8 : REM LDA #Ø8
33Ø DATA 133,143 : REM STA $8F
34Ø DATA 136 : REM DEY
35Ø DATA 198,252 : REM DEC $FC
36Ø :: REM ROTATE-BYTE
37Ø DATA 38,142 : REM ROL $8E
38Ø DATA 144,16 : REM BCC $1Ø
39Ø DATA 189,6,1 : REM LDA $1Ø6,X
4ØØ DATA 145,251 : REM STA ($FB),Y
41Ø DATA 136 : REM DEY
42Ø DATA 36,141 : REM BIT $8D
43Ø DATA 16,6 : REM BPL $Ø6
44Ø DATA 189,5,1 : REM LDA $1Ø5,X
45Ø DATA 145,251 : REM STA ($FB),Y
46Ø DATA 136 : REM DEY
47Ø :: REM BIT-CLEAR
48Ø DATA 2Ø2 : REM DEX
49Ø DATA 38,141 : REM ROL $8D
5ØØ DATA 144,1 : REM BCC $Ø1
51Ø DATA 2Ø2 : REM DEX
52Ø :: REM OVER
53Ø DATA 198,143 : REM DEC $8F
54Ø DATA 2Ø8,226 : REM BNE $E2
55Ø DATA 56 : REM SEC
56Ø DATA 152 : REM TYA

```

```

57Ø DATA 1Ø1,251 : REM ADC $FB
58Ø DATA 133,251 : REM STA $FB
59Ø DATA 144,2 : REM BCC $Ø2
6ØØ DATA 23Ø,252 : REM INC $FC
61Ø :: REM CLEAR-STACK
62Ø DATA 162,Ø : REM LDX #Ø
63Ø :: REM REPEAT
64Ø DATA 1Ø4 : REM PLA
65Ø DATA 149,139 : REM STA $8B,X
66Ø DATA 232 : REM INX
67Ø DATA 224,6 : REM CPX #$$Ø6
68Ø DATA 2Ø8,248 : REM BNE $F8
69Ø DATA 1Ø4,168 : REM PLA : TAY
7ØØ DATA 1Ø4,17Ø : REM PLA : TAX
71Ø DATA 1Ø4 : REM PLA
72Ø DATA 4Ø : REM PLP
73Ø DATA 96 : REM RTS
74Ø :: REM TEST-ROUTINE
75Ø DATA 169,24Ø : REM LDA #$$Ø
76Ø DATA 162,15 : REM LDX #$$ØF
77Ø DATA 16Ø,255 : REM LDY #$$F
78Ø DATA 32,Ø,192 : REM JSR $CØØØ
79Ø DATA 255 : REM EMBEDDED-BYTE
8ØØ DATA 96 : REM RTS
81Ø :
82Ø REM ** SET UP ZERO PAGE AND FREE RAM **
83Ø PRINT CHR$(147)
84Ø POKE 251,12 POKE 252,197
85Ø FOR N=139 TO 144 : POKE N,N : NEXT
86Ø FOR N=5Ø432 TO 5Ø44Ø : POKE N,Ø : NEXT
87Ø :
88Ø SYS 49258 : REM SYS TEST-ROUTINE
89Ø :
9ØØ REM ** READ RESULTS **
91Ø FOR LOOP=5Ø432 TO 5Ø443
92Ø READ NAME$
93Ø PRINT NAME$;
94Ø PRINT PEEK(LOOP)
95Ø NEXT LOOP
96Ø :

```



```

97Ø DATA "ZERO PAGE  ", "ZERO PAGE+1"
98Ø DATA "ZERO PAGE+2", "ZERO PAGE+3"
99Ø DATA "ZERO PAGE+4", "ZERO PAGE+5"
1ØØØ DATA "Y REGISTER ", "X REGISTER "
1Ø1Ø DATA "ACCUMULATOR", "STATUS      "
1Ø2Ø DATA "PC LOW      ", "PC HIGH     "

```

The problem to solve next is that of where to place the user stack. This will depend on your own requirements, so to make the whole thing flexible, a vectored address in the bytes at \$FB and \$FC contains the stack address. In the program listed above, this is \$C512 (line 84Ø). The vectored address is, in fact, the address + 12. This is because the stack is pushed in reverse (decreasing) order.

When executed, the coding first pushes all the processor registers on to the hardware stack and moves the stack pointer across into the X register (lines 9Ø to 14Ø). Next, the six zero page pseudo-registers are pushed there (lines 15Ø to 19Ø). The return address from the subroutine call is then incremented on the stack, using the contents of the X register (stack pointer) to access it (lines 2ØØ to 24Ø). The two bytes that form the RTS address are copied into pseudo-register 1 (now safely on the hardware stack) to form a vector through which the embedded data byte can be loaded into the accumulator and then saved for use in zero page (lines 25Ø to 31Ø).

In line 28Ø, a pre-defined byte was loaded into the accumulator and saved in zero page. This byte holds a bit code that will inform the program as to whether the register being pulled from the hardware stack for transfer to the software stack is one or two bytes long. The byte value, \$87, is 1ØØØØ111 in binary and the set bits correspond to the two-byte registers, the Program Counter and the three pseudo-registers. By rotating this byte left after each pull operation and using the BIT operation, the Negative flag can be tested to see if a further pull is needed. All this and the copy hardware stack/push software stack is handled by lines 32Ø to 55Ø.

Finally, the registers and pseudo-registers are restored to their original values (lines 62Ø to 73Ø). The test routine between lines 75Ø and 8ØØ shows the way the program is used. When run, the test procedure produces the following output on the screen:

```

ZERO PAGE      139
ZERO PAGE+1    14Ø
ZERO PAGE+2    141
ZERO PAGE+3    142
ZERO PAGE+4    143
ZERO PAGE+5    144
Y REGISTER     255
X REGISTER     15

```

ACCUMULATOR	240
STATUS	176
PC LOW	115
PC HIGH	192

As can be seen, the zero page bytes contain the values POKEd into them by the FOR...NEXT loop of line 830 while the accumulator and Index registers display their seeded values (lines 750 to 770). The Program Counter holds  $192 * 256 + 115$ , or \$C073, which was the point in the program where its contents were pushed at line 780.

This program could be extended to provide a routine to perform a pull user stack, to copy the contents of a software stack into the processor and pseudo-registers.

### Line-by-line

A line-by-line description of Program 17 follows:

```

line 90 : save all processor registers on hardware stack
line 140 : move stack pointer into X for index
line 150 : entry for PUSH-ZERO-PAGE
line 160 : get zero page byte
line 170 : push on to hardware stack
line 180 : decrement index
line 190 : branch to PUSH-ZERO-PAGE until done
line 200 : increment low byte of RTS address
line 210 : get it from stack
line 220 : and save in zero page
line 230 : if not equal branch to PC-LOW
line 240 : else increment page byte of RTS address
line 250 : entry for PC-LOW
line 260 : get high byte of RTS address
line 270 : and save it to form vector
line 280 : get bit code to indicate register size
line 290 : and save it
line 300 : get embedded code after subroutine call
line 310 : and save it
line 320 : eight bits in embedded byte to test
line 330 : save bit count
line 340 : decrement index to $FF
line 350 : decrement high byte of vectored address at $FB

```

line 360 : entry for ROTATE-BYTE  
line 370 : move next coded bit into Carry flag  
line 380 : if bit clear skip it, branch to BIT-CLEAR  
line 390 : otherwise get byte from stack  
line 400 : save it on user stack  
line 410 : decrement index  
line 420 : is it a two byte register?  
line 430 : no, so branch to BIT-CLEAR  
line 440 : yes, so get the second byte from the stack  
line 450 : and save it on the user stack  
line 460 : decrement index  
line 470 : entry for BIT-CLEAR  
line 480 : decrement hardware stack index  
line 490 : move bit of register code into Carry flag  
line 500 : if clear, branch to OVER  
line 510 : else decrement hardware stack index  
line 520 : entry for OVER  
line 530 : decrement bit counter  
line 540 : and repeat until all done  
line 550 : set Carry flag  
line 560 : move user stack pointer into accumulator  
line 570 : add to low byte of address  
line 580 : and save  
line 590 : branch to CLEAR-STACK if carry is clear  
line 600 : else increment high byte of address  
line 610 : entry for CLEAR-STACK  
line 620 : initialize X register  
line 630 : entry for REPEAT  
line 640 : pull byte from stack  
line 650 : and restore zero page  
line 660 : increment index  
line 670 : all bytes restored?  
line 680 : no, branch to REPEAT  
line 690 : yes, restore all registers  
line 730 : back to calling routine  
line 740 : entry for TEST-ROUTINE  
line 750 : seed registers  
line 780 : call user stack routine  
line 790 : embedded byte  
line 800 : back to BASIC

## BINARY INS AND OUTS

Sometimes when printing the values of registers, it is necessary to have their binary representation—for example, in the case of the Status register, because we are concerned with the state of the particular bits within it, rather than the overall value of the contents. Program 18 provides a short routine which produces such a binary output from a decimal input. This could easily be adapted for use within a program such as the software stack given above.

### Program 18

```
10 REM ** PRINT ACCUMULATOR AS A **
20 REM ** BINARY NUMBER **
30 CODE=49152
40 FOR LOOP=0 TO 17
50 READ BYTE
60 POKE CODE+LOOP,BYTE
70 NEXT LOOP
80 :
90 REM ** M/C DATA **
100 DATA 162,0 : REM LDX #$08
110 DATA 72 : REM PHA
120 :: REM NEXT-BIT
130 DATA 104 : REM PLA
140 DATA 10 : REM ASL A
150 DATA 72 : REM PHA
160 DATA 169,48 : REM LDA #$30
170 DATA 105,0 : REM ADC #$00
180 DATA 32,210,255 : REM JSR $FFD2
190 DATA 202 : REM DEX
200 DATA 208,243 : REM BNE $F3
210 DATA 104 : REM PLA
220 DATA 96 : REM RTS
230 :
240 REM ** SET UP DEMO RUN **
250 REM LDA $FB : JSR $C000 : RTS
260 POKE 820,165 : POKE 821,251
270 POKE 822,32 : POKE 823,0
280 POKE 824,192 : POKE 825,96
290 PRINT CHR$(147) PRINT
300 INPUT "INPUT A NUMBER ";A$
```

```

31Ø A=VAL(A$)
32Ø POKE 251,A
33Ø PRINT"BINARY VALUE IS :";
34Ø SYS 82Ø

```

### Line-by-line

The following line-by-line description should make the program's operation clear. It is simply moving each bit of the accumulator in turn into the Carry flag position, using the arithmetic shift left operation (see Figure 8.1) and adding its value to the ASCII code for Ø, i.e.

accumulator=48+carry

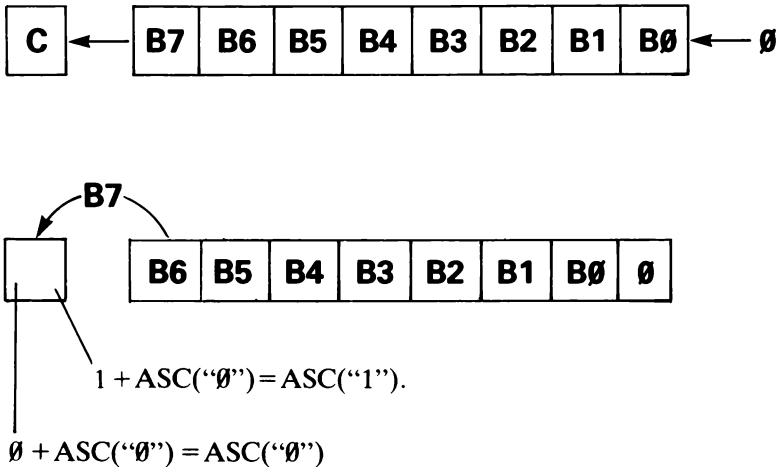


Figure 8.1 Arithmetic shift left

If the Carry flag is clear, the result will be  $48 + \text{Ø} = 48$ , so the CHROUT routine will print a Ø. On the other hand, if the Carry flag is set, the result of the addition will be  $48 + 1 = 49$ , so a 1 will be printed by CHROUT.

```

line 1ØØ : eight bits in a byte
line 11Ø : push accumulator on to stack
line 12Ø : entry for NEXT-BIT

```

```

line 130 : restore accumulator
line 140 : shift bit 7 into carry
line 150 : save shifted accumulator on stack
line 160 : get ASCII code for 0
line 170 : add carry
line 180 : print either 0 or 1
line 190 : decrement bit counter
line 200 : do NEXT-BIT until complete
line 210 : pull stack to balance push
line 220 : back to BASIC

```

## COME IN

By reversing this process, it is possible to input a number directly into the accumulator in binary form as Program 19 shows. The program scans the keyboard for a pressed 1 or 0 key and the Carry flag is set or cleared respectively. A copy of the accumulator, initially cleared, is kept on the hardware stack and restored each time round to rotate the carry bit into it using the rotate left operation (see Figure 8.2). The loop is executed eight times, once for each bit, and on completion, the accumulator holds the hexadecimal value of the binary number.

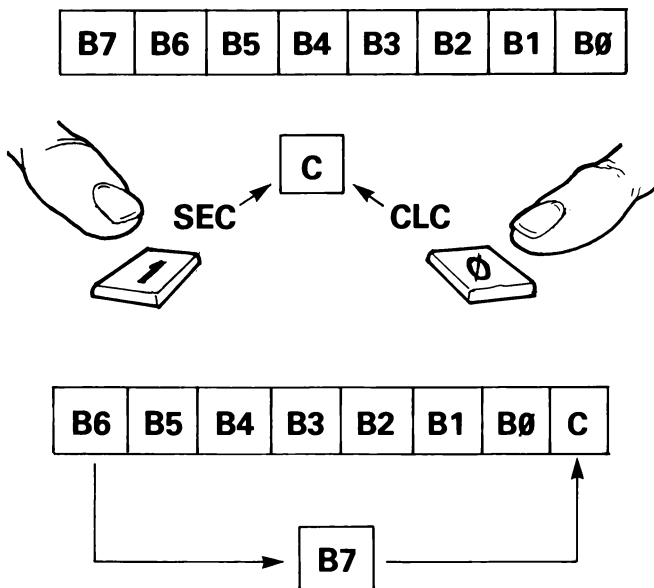


Figure 8.2 Input a number directly into the accumulator

## Program 19

```
1Ø REM ** INPUT A HEX NUMBER IN BINARY FORM **
2Ø CODE=49152
3Ø FOR LOOP=Ø TO 41
4Ø READ BYTE
5Ø POKE CODE+LOOP,BYTE
6Ø NEXT LOOP
7Ø :
8Ø REM ** M/C DATA **
9Ø DATA 162,8 : REM LDX #$Ø8
1ØØ DATA 169,Ø : REM LDA #$ØØ
11Ø DATA 72 : REM PHA
12Ø DATA 24 : REM CLC
13Ø :: REM MAINLOOP
14Ø DATA 134,243 : REM STX $FD
15Ø :: REM LOOP
16Ø DATA 32,228,255 : REM JSR $FFE4
17Ø DATA 24Ø,251 : REM BEQ $FB
18Ø DATA 2Ø1,49 : REM CMP #$31
19Ø DATA 24Ø,7 : REM BEQ $Ø7
2ØØ DATA 2Ø1,48 : REM CMP #$3Ø
21Ø DATA 2Ø8,243 : REM BNE $F3
22Ø DATA 24 : REM CLC
23Ø DATA 144,1 : REM BCC $Ø1
24Ø :: REM SET
25Ø DATA 56 : REM SEC
26Ø :: REM OVER
27Ø DATA 8 : REM PHP
28Ø DATA 32,21Ø,255 : REM JSR $FFD2
29Ø DATA 4Ø : REM PLP
3ØØ DATA 1Ø4 : REM PLA
31Ø DATA 42 : REM ROL A
32Ø DATA 72 : REM PHA
33Ø DATA 166,253 : REM LDX $FD
34Ø DATA 2Ø2 : REM DEX
35Ø DATA 2Ø8,224 : REM BNE $EØ
36Ø DATA 1Ø4 : REM PLA
37Ø DATA 133,251 : REM STA $FB
38Ø DATA 96 : REM RTS
```

```

39Ø :
40Ø PRINT CHR$(147)
41Ø PRINT
42Ø PRINT"INPUT YOUR BINARY NUMBER :";
43Ø SYS CODE
44Ø PRINT PEEK(251)

```

### Line-by-line

A line-by-line explanation of Program 19 now follows:

```

line 9Ø : eight bits to read
line 10Ø : clear accumulator—shift register
line 11Ø : push it on to stack
line 12Ø : clear the Carry flag
line 13Ø : entry for MAINLOOP
line 14Ø : save X register
line 15Ø : entry for LOOP
line 16Ø : jump to GETIN
line 17Ø : if null, branch to LOOP
line 18Ø : is it ASC"1"?
line 19Ø : yes, branch to SET
line 20Ø : is it ASC"Ø"?
line 21Ø : no, branch to LOOP
line 22Ø : yes, clear Carry flag
line 23Ø : and force branch to OVER
line 24Ø : entry for SET
line 25Ø : set Carry flag
line 26Ø : entry for OVER
line 27Ø : save Carry flag on stack
line 28Ø : print Ø or 1
line 29Ø : restore Carry flag
line 30Ø : restore accumulator
line 31Ø : move Carry flag into bit Ø
line 32Ø : save accumulator
line 33Ø : restore bit count
line 34Ø : decrement it by one
line 35Ø : branch to MAINLOOP until all done
line 36Ø : restore accumulator
line 37Ø : save in zero page
line 38Ø : back to BASIC

```



## **Project**

Convert the software stack program to print the binary values of each register upon completion.

Modify it further to allow register values to be seeded into the software stack test routine, using the binary input routine. Note that you should only attempt seeding the accumulator and Index registers. Why?

# 9 Move, Fill and Dump

## MOVE IT!

The ability to move blocks of memory around within the bounds of the memory map is a necessity. When manipulating hi-resolution graphics, for example, large blocks of memory need to be moved around quickly and smoothly. The program could also be used to relocate sections of machine code rather than rewriting the assembler that created them—assuming, of course, that your code has been designed to make it portable.

At first sight, it may seem that the simplest method of moving a block of memory is to take the first byte to be moved and store it at the destination address, take the second byte and place it at the destination address + 1, and so forth. There would be no problem here if the destination address was outside the source address, but consider what would happen if the destination address was within the bounds to be searched by the source address—that is, the two regions overlapped. Figure 9.1 illustrates the problem using this straightforward method to move a block of five bytes forward by just a single byte, relocating the five bytes from \$C500 to \$C501.

Using the obvious method, the first character, 'S', is moved from \$C500 to \$C501 thereby overwriting the 'A'. The program then takes the next character at location START+1 (\$C501), the 'S' that has just been written there, and places it at START+2 (\$C502)

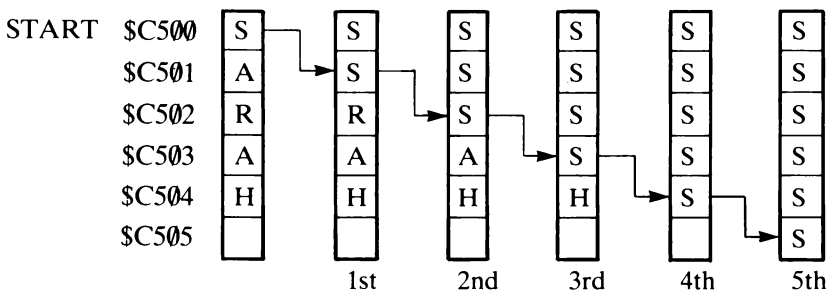


Figure 9.1 The overwriting move sequence

overwriting the 'R'. As you can see, the end result is SSSSS—the whole block is full of 'S's—not the required effect!

To avoid this problem, the MOVE routine acts 'intelligently' and if it calculates that an overwrite would occur, performs the movement of bytes in the reverse order, starting at the highest address and moving down the memory map as Figure 9.2 shows.

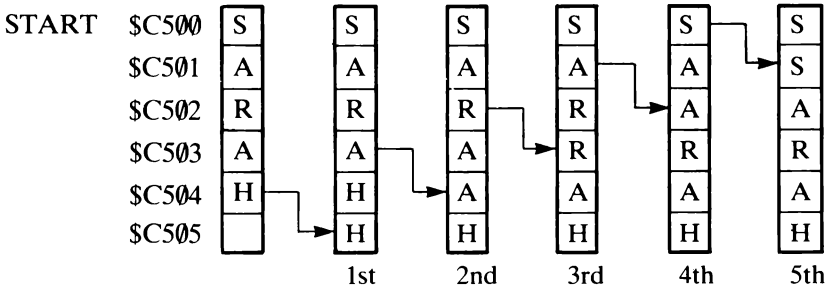


Figure 9.2 The correct move sequence

### Program 20

```

1Ø REM ** MEMORY BLOCK MOVE ROUTINE **
2Ø REM ** 1Ø9 BYTES LONG WHEN ASSEMBLED **
3Ø REM ** PLUS 5 DATA BYTES IN ZERO PAGE **
4Ø CODE=49152
5Ø FOR LOOP=Ø TO 1Ø8
6Ø READ BYTE
7Ø POKE CODE+LOOP,BYTE
8Ø NEXT LOOP
9Ø :
1ØØ REM ** M/C DATA **
11Ø DATA 56 : REM SEC
12Ø DATA 165,251 : REM LDA $FB
13Ø DATA 229,253 : REM SBC $FD
14Ø DATA 17Ø : REM TAX
15Ø DATA 165,252 : REM LDA $FC
16Ø DATA 229,254 : REM SBC $FE
17Ø DATA 168 : REM TAY
18Ø DATA 138 : REM TXA
19Ø DATA 2Ø5,52,3 : REM CMP $334
2ØØ DATA 152 : REM TYA
21Ø DATA 237,53,3 : REM SBC $335

```

```

220 DATA 176,2      : REM BCS $02
230 DATA 144,35     : REM BCC $23
240 ::              REM MOVE-LEFT
250 DATA 160,0      : REM LDY #000
260 DATA 174,53,3  : REM LDX $335
270 DATA 240,14     : REM BEQ $0E
280 ::              REM LEFT-COMPLETE-PAGES
290 DATA 177,253   : REM LDA ($FD),Y
300 DATA 145,251   : REM STA ($FB),Y
310 DATA 200        : REM INY
320 DATA 208,249   : REM BNE $F9
330 DATA 230,254   : REM INC $FE
340 DATA 230,252   : REM INC $FC
350 DATA 202        : REM DEX
360 DATA 208,242   : REM BNE $F2
370 ::              REM LEFT-PARTIAL-PAGE
380 DATA 174,52,3  : REM LDX $334
390 DATA 240,8      : REM BEQ $08
400 ::              REM LAST-LEFT
410 DATA 177,253   : REM LDA ($FD),Y
420 DATA 145,251   : REM STA ($FB),Y
430 DATA 200        : REM INY
440 DATA 202        : REM DEX
450 DATA 208,248   : REM BNE $F8
460 ::              REM EXIT
470 DATA 96         : REM RTS
480 :
490 ::              REM MOVE-RIGHT
500 DATA 24         : REM CLC
510 DATA 173,53,3  : REM LDA $335
520 DATA 72         : REM PHA
530 DATA 101,254   : REM ADC $FE
540 DATA 133,254   : REM STA $FE
550 DATA 24         : REM CLC
560 DATA 104        : REM PLA
570 DATA 101,252   : REM ADC $FC
580 DATA 133,252   : REM STA $FC
590 DATA 172,52,3  : REM LDY $334
600 DATA 240,9      : REM BEQ $09

```

```

61Ø :: REM TRANSFER
62Ø DATA 136 : REM DEY
63Ø DATA 177,253 : REM LDA ($FD),Y
64Ø DATA 145,251 : REM STA ($FB),Y
65Ø DATA 192,Ø : REM CPY #ØØØ
66Ø DATA 2Ø8,247 : REM BNE $F7
67Ø :: REM RIGHT-COMPLETE-PAGES
68Ø DATA 174,53,3 : REM LDX $335
69Ø DATA 24Ø,221 : REM BEQ $DD
7ØØ :: REM UPDATE
71Ø DATA 198,254 : REM DEC $FE
72Ø DATA 198,252 : REM DEC $FC
73Ø :: REM PAGE
74Ø DATA 136 : REM DEY
75Ø DATA 177,253 : REM LDA ($FD),Y
76Ø DATA 145,251 : REM STA ($FB),Y
77Ø DATA 192,Ø : REM CPY #ØØØ
78Ø DATA 2Ø8,247 : REM BNE $F7
79Ø DATA 2Ø2 : REM DEX
8ØØ DATA 2Ø8,24Ø : REM BNE $FØ
81Ø DATA 96 : REM RTS
82Ø :
83Ø REM ** SET UP VARIABLES **
84Ø PRINT CHR$(147)
85Ø PRINT" *** MEMORY MOVER V1.1 ***"
86Ø INPUT"START ADDRESS ";S
87Ø INPUT"DESTINATION ";D
88Ø INPUT"LENGTH IN BYTES ";L
89Ø :
9ØØ S1=INT(S/256) : S2=S-(S1*256)
91Ø D1=INT(D/256) : D2=D-(D1*256)
92Ø L1=INT(L/256) : L2=L-(L1*256)
93Ø :
94Ø POKE 251,D2 : POKE 252,D1
95Ø POKE 253,S2 : POKE 254,S1
96Ø POKE 82Ø,L2 : POKE 821,L1
97Ø :
98Ø REM ** SET UP DEMO **
99Ø FOR N=Ø TO 15
1ØØØ POKE 828+N,N

```

```

1010 POKE 900+N,0
1020 NEXT N
1030 :
1040 SYS CODE
1050 :
1060 REM ** PRINT THE RESULTS! **
1070 FOR N=0 TO 15
1080 PRINT PEEK(828+N);"      ";
1090 PRINT PEEK(900+N)
1100 NEXT N

```

Bytes reserved:

```

251-252      ($FB-$FC)   : Destination vector
253-254      ($FD-$FE)   : Source vector
820-821      ($334-$335) : Length of block to be
                        moved

```

When run, the BASIC test requests three inputs: the START address of the memory block to be moved, its DESTINATION address and its LENGTH in bytes. All values should be entered as decimal values. Thus, to move a 1K block of memory from 49152 to 56000, the values to input are:

```

START ADDRESS      : 49152
DESTINATION        : 56000
LENGTH IN BYTES   : 1024

```

For reasons already explained, the coding begins by ascertaining whether a left-move or a right-move operation is required. It calculates this (lines 110 to 210) by subtracting the source address from the destination address. If the result is less than the number of bytes to be moved, overwriting would occur using the MOVE-LEFT routine, so the MOVE-RIGHT coding is called (line 230). If the memory locations do not overlap, the quicker MOVE-LEFT routine is selected (line 220). For further description purposes we will examine the MOVE-LEFT routine (lines 240 to 470).

Memory movement is performed in two phases: complete memory pages are first relocated, and then any remaining bytes in the final partial page are moved. These details are held in the length of block bytes \$334 and \$335.

The routine begins by loading the number of pages to be moved into the X register (line 260), branching to LEFT-PARTIAL-PAGE if it is zero (line 280). Transfer of data bytes is completed using post-indexed indirect addressing through the zero page vectors. When all the whole pages have been transferred, any

remaining bytes are transferred by the LEFT-PARTIAL-PAGE loop (lines 370 to 450).

The MOVE-RIGHT routine is similar in operation, except that it starts at the highest memory location referenced and moves down through memory, the highest address of the source and destination being calculated in lines 500 to 650.

### **Line-by-line**

A line-by-line description of Program 20 now follows:

line 110 : set Carry flag  
line 120 : get low byte destination address  
line 130 : subtract low byte source address  
line 140 : transfer result into X register  
line 150 : get high byte destination address  
line 160 : subtract high byte source address  
line 170 : save result in X register  
line 180 : restore result of low byte subtraction  
line 190 : compare it with low byte of length  
line 200 : restore result of high byte subtraction  
line 210 : subtract high byte of length from it  
line 220 : if Carry flag set, branch to MOVE-LEFT  
line 230 : else branch to MOVE-RIGHT  
line 240 : entry for MOVE-LEFT  
line 250 : initialize index  
line 260 : get number of pages to be moved  
line 270 : if zero, branch to LEFT-PARTIAL-PAGE  
line 280 : entry for LEFT-COMPLETE-PAGES  
line 290 : get source byte  
line 300 : store at destination  
line 310 : increment index  
line 320 : branch to LEFT-COMPLETE-PAGES until page  
done  
line 330 : increment source page  
line 340 : increment destination page  
line 350 : decrement page counter  
line 360 : branch to LEFT-COMPLETE-PAGES until all moved  
line 370 : entry for LEFT-PARTIAL-PAGE  
line 380 : get number of bytes on page to be moved  
line 390 : if zero, branch to EXIT

line 400 : entry for LAST-LEFT  
line 410 : get source byte  
line 420 : store at destination  
line 430 : increment index  
line 440 : decrement byte count  
line 450 : branch to LAST-LEFT until done  
line 460 : entry for EXIT  
line 470 : back to BASIC  
line 490 : entry for MOVE-RIGHT  
line 500 : clear Carry flag  
line 510 : get number of pages to be moved  
line 520 : save on stack  
line 530 : add it to source high byte  
line 540 : and save result  
line 550 : reclear Carry flag  
line 560 : get length high byte off stack  
line 570 : add it to destination high byte  
line 580 : and save the result  
line 590 : get low byte of length into Y register  
line 600 : branch to RIGHT-COMPLETE-PAGES if zero  
line 610 : entry for TRANSFER  
line 620 : decrement index  
line 630 : get source byte  
line 640 : and copy to destination  
line 650 : is Y = 0?  
line 660 : no, branch to TRANSFER  
line 670 : entry for RIGHT-COMPLETE-PAGES  
line 680 : get number of pages to be moved  
line 690 : if zero, branch to EXIT  
line 700 : entry for UPDATE  
line 710 : decrement number of pages to do  
line 720 : and also destination  
line 730 : entry for PAGE  
line 740 : decrement index  
line 750 : get source byte  
line 760 : copy to destination  
line 770 : is Y = 0?  
line 780 : no, branch to PAGE



line 790 : decrement page counter  
line 800 : if not zero, branch to UPDATE  
line 810 : return to BASIC

## FILL

Program 21 provides the BASIC loader listing to implement a memory FILL routine, which is particularly useful for clearing sections of RAM with a pre-determined value.

### Program 21

```
10 REM ** MEMORY FILL ROUTINE **
20 REM ** 30 BYTES LONG WHEN ASSEMBLED **
30 REM ** PLUS 5 DATA BYTES IN ZERO PAGE **
40 CODE=49152
50 FOR LOOP=0 TO 30
60 READ BYTE
70 POKE CODE+LOOP,BYTE
80 NEXT LOOP
90 :
100 REM ** M/C DATA **
110 DATA 165,255 : REM LDA $FF
120 DATA 166,252 : REM LDX $FC
130 DATA 240,12 : REM BEQ $0C
140 DATA 160,0 : REM LDY #$00
150 :: REM COMPLETE-PAGE
160 DATA 145,253 : REM STA ($FD),Y
170 DATA 200 : REM INY
180 DATA 208,251 : REM BNE $FB
190 DATA 230,254 : REM INC $FE
200 DATA 202 : REM DEX
210 DATA 208,246 : REM BNE $F6
220 :: REM PARTIAL-PAGE
230 DATA 166,251 : REM LDX $FB
240 DATA 240,8 : REM BEQ $08
250 DATA 160,0 : REM LDY #$00
260 :: REM AGAIN
270 DATA 145,253 : REM STA ($FD),Y
```

```

280 DATA 200      : REM INY
290 DATA 202      : REM DEX
300 DATA 208,250  : REM BNE $FA
310 ::            : REM FINISH
320 DATA 96       : REM RTS
330 :
340 REM ** GET DETAILS **
350 PRINT CHR$(147)
360 INPUT"FILL DATA      :";F
370 INPUT"START ADDRESS  :";S
380 INPUT"NUMBER OF BYTES :";L
390 :
400 S1=INT(S/256)  : S2=S-(S1*256)
410 L1=INT(L/256)  : L2=L-(L1*256)
420 :
430 POKE 251,L2    : POKE 252,L1
440 POKE 253,S2    : POKE 254,S1
450 POKE 255,F
460 :
470 SYS CODE

```

Bytes reserved:

```

251-252    ($FB-$FC) : number of bytes to be filled
253-254    ($FD-$FE) : start of address of bytes to be
                    filled
255        ($FF)      : value to fill with

```

When executed, the machine code expects to find the fill value, the start address and the amount of memory to be filled, in five zero page bytes of memory from \$FB. Input of each of these is handled by a few lines of BASIC from line 360. To clear a 1K block of RAM from \$C500 with zero, the following information should be entered in response to the 64's prompt:

```

FILL DATA      : 0
START ADDRESS   : 49152
NUMBER OF BYTES : 1024

```

The FILL routine works in a similar manner to the MOVE routine described above, dealing with whole and partial pages separately. The main fill loop is embodied in lines 150 to 300.

## Line-by-line

A line-by-line description of the program now follows:

```
line 11Ø : get data with which to fill
line 12Ø : get number of complete pages to be filled
line 13Ø : if zero, branch to PARTIAL-PAGE
line 14Ø : initialize index
line 15Ø : entry for COMPLETE-PAGE
line 16Ø : fill byte
line 17Ø : increment index
line 18Ø : branch to COMPLETE-PAGE until all of page is
           done
line 19Ø : increment page
line 20Ø : decrement page counter
line 21Ø : branch to COMPLETE-PAGE until all pages are
           filled
line 22Ø : entry for PARTIAL-PAGE
line 23Ø : get number of bytes left to be filled
line 24Ø : if zero, branch to FINISH
line 25Ø : else clear index
line 26Ø : entry for AGAIN
line 27Ø : fill byte
line 28Ø : increment index
line 29Ø : decrement bytes left to do count
line 30Ø : branch to AGAIN until all filled
line 31Ø : entry for FINISH
line 32Ø : back to BASIC
```

## A MEMORY DUMP

A hex and ASCII dump of memory can be extremely useful, not only within machine code programs, but also when used from a BASIC program. Most often it provides information about the way a program is manipulating numeric and string variables and tables. Figure 9.3 shows the type of dump produced by the routine: twenty-four lines of eight bytes each. The example shows some text stored in memory. Each line starts with the current address, followed by the eight bytes stored in memory from that point. The far right of the listing provides the ASCII equivalents of each byte. Any non-ASCII character (that is, one greater than \$7F) or control code (those less than \$20) is represented by a full stop.

```

C108 : 54 68 69 73 20 69 73 20 This is
C110 : 61 20 73 69 6D 70 6C 65 a simple
C118 : 20 65 78 61 6D 70 6C 65 example
C120 : 20 6F 66 20 68 6F 77 20 of how
C128 : 74 68 65 20 8D 64 75 6D the .dum
C130 : 70 20 72 6F 75 74 69 6E p routin
C138 : 65 20 66 6F 72 20 74 68 e for th
C140 : 65 20 43 6F 6D 6D 6F 64 e Commod
C148 : 6F 72 65 20 36 34 20 8D ore 64 .
C150 : 77 6F 72 6B 73 2E 0D 54 works..T
C158 : 68 65 20 64 75 6D 70 20 he dump
C160 : 63 61 6E 20 62 65 20 64 can be d
C168 : 69 76 69 64 65 64 20 69 ived i
C170 : 6E 74 6F 20 74 68 72 65 nto thre
C178 : 65 20 8D 73 65 63 74 69 e .secti
C180 : 6F 6E 73 2E 20 54 68 65 ons. The
C188 : 20 66 69 72 73 74 20 63 first c
C190 : 6F 6C 75 6D 6E 20 6C 69 olumn li
C198 : 73 74 73 20 74 68 65 20 sts the
C1A0 : 8D 73 74 61 72 74 20 61 .start a
C1A8 : 64 64 72 65 73 73 20 6F ddress o
C1B0 : 66 20 74 68 65 20 62 6C f the bl
C1B8 : 6F 63 6B 2E 20 54 68 65 ock. The
C1C0 : 20 73 65 63 6F 6E 64 20 second
C1C8 : 8D 63 6F 6C 75 6D 6E 20 .column
C1D0 : 69 73 20 69 6E 20 66 61 is in fa
C1D8 : 63 74 20 74 68 65 20 68 ct the h
C1E0 : 65 78 61 64 65 63 69 6D exadecim
C1E8 : 61 6C 20 8D 76 61 6C 75 al .valu
C1F0 : 65 73 20 6F 66 20 65 69 es of ei
C1F8 : 67 68 74 20 62 79 74 65 ght byte
C200 : 73 20 66 72 6F 6D 20 74 s from t
C208 : 68 69 73 20 8D 61 64 64 his .add
C210 : 72 65 73 73 2E 20 46 69 ress. Fi
C218 : 6E 61 6C 6C 79 20 74 68 nally th
C220 : 65 20 6C 61 73 74 20 63 e last c
C228 : 6F 6C 75 6D 6E 20 8D 64 olumn .d
C230 : 65 70 69 63 74 73 20 74 epicts t
C238 : 68 65 20 41 53 43 49 49 he ASCII
C240 : 20 76 61 6C 75 65 73 20 values
C248 : 6F 66 20 74 68 65 73 65 of these
C250 : 20 8D 62 79 74 65 73 2E .bytes.
C258 : 20 75 6E 6C 65 73 73 20 unless
C260 : 74 68 65 20 62 79 74 65 the byte
C268 : 20 69 73 20 6E 6F 6E 2D is non-
C270 : 41 53 43 49 49 20 8D 77 ASCII .w
C278 : 68 69 63 68 20 69 73 20 hich is
C280 : 74 68 65 6E 20 64 69 73 then dis
C288 : 70 6C 61 79 65 64 20 61 played a
C290 : 73 20 61 2C 66 75 6C 6C s a full
C298 : 20 73 74 6F 70 21 0D 00 stop!..
C2A0 : 00 4C 00 C9 A9 FF 85 22 .L....."
C2A8 : 08 60 00 00 00 00 00 00 .^.....

```

Figure 9.3 Memory dump

As it stands, the routine requires three zero page data bytes, two for the start address and one for the number of eight byte lines to be dumped. The routine also employs the ADDRESS-PRINT and HEXPRINT routines discussed earlier.

## Program 22

```

100 REM ** DUMP LINES OF 8 BYTES OF **
200 REM ** MEMORY IN HEX AND ASCII **
300 CODE=49152
400 FOR LOOP=0 TO 111
500 READ BYTE
600 POKE CODE+LOOP,BYTE
700 NEXT LOOP
800 :
900 REM ** M/C DATA **
1000 DATA 32,71,192 : REM JSR $C047
1100 :: REM HEX-BYTES
1200 DATA 162,7 : REM LDX #$07
1300 DATA 160,0 : REM LDY #$00
1400 REM HEX-LOOP
1500 DATA 177,251 : REM LDA ($FB),Y
1600 DATA 32,90,192 : REM JSR $C05A
1700 DATA 32,66,192 : REM JSR $C042
1800 DATA 200 : REM INY
1900 DATA 202 : REM DEX
2000 DATA 16,244 : REM BPL $F4
2100 DATA 32,66,192 : REM JSR $C042
2200 :: REM ASCII-BYTES
2300 DATA 162,7 REM LDX #$07
2400 DATA 160,0 : REM LDY #$00
2500 :: REM ASCII-LOOP
2600 DATA 177,251 : REM LDA ($FB),Y
2700 DATA 201,32 : REM CMP #$20
2800 DATA 48,4 : REM BMI $04
2900 DATA 201,128 : REM CMP #$80
3000 DATA 144,2 : REM BCC $02
3100 :: REM FULL-STOP
3200 DATA 169,46 REM LDA #$2E
3300 :: REM LEAP-FROG

```

```

340 DATA 32,210,255 : REM JSR $FFD2
350 DATA 200 : REM INY
360 DATA 202 : REM DEX
370 DATA 16,237 : REM BPL $ED
380 DATA 169,13 : REM LDA #00
390 DATA 32,210,255 : REM JSR $FFD2
400 DATA 24 : REM CLC
410 DATA 165,251 : REM LDA $FB
420 DATA 105,8 : REM ADC #08
430 DATA 133,251 : REM STA $FB
440 DATA 144,2 : REM BCC 02
450 DATA 230,252 : REM INC $FC
460 :: REM NO-CARRY
470 DATA 198,254 : REM DEC $FE
480 DATA 208,191 : REM BNE $BF
490 DATA 96 : REM RTS
500 :: REM SPACE
510 DATA 169,32 : REM LDA #20
520 DATA 76,210,255 : REM JMP $FFD2
530 :: REM ADDRESS-PRINT
540 DATA 162,251 : REM LDX $FB
550 DATA 181,1 : REM LDA 1,X
560 DATA 32,90,192 : REM JSR $C05A
570 DATA 181,0 : REM LDA 0,X
580 DATA 32,90,192 : REM JSR $C05A
590 DATA 32,66,192 : REM JSR $C042
600 DATA 32,66,192 : REM JSR $C042
610 DATA 96 : REM RTS
620 :: REM HEXPRINT
630 DATA 72 : REM PHA
640 DATA 74,74 : REM LSR A : LSR A
650 DATA 74,74 : REM LSR A : LSR A
660 DATA 32,99,192 : REM JSR $C063
670 DATA 104 : REM PLA
680 :: REM FIRST
690 DATA 41,15 : REM AND #0F
700 DATA 201,10 : REM CMP #0A
710 DATA 144,2 : REM BCC 02
720 DATA 105,6 : REM ADC #06
730 :: REM OVER

```

```

74Ø DATA 1Ø5,48      : REM  ADC #Ø3Ø
75Ø DATA 76,21Ø,255 : REM  JMP $FFD2
76Ø :
77Ø REM ** INPUT DETAILS FOR DUMP **
78Ø PRINT CHR$(147)
79Ø INPUT"DUMP START ADDRESS ";A
8ØØ HIGH=INT(A/256)
81Ø LOW=A-(HIGH*256)
82Ø POKE 251,LOW : POKE 252,HIGH
83Ø INPUT"NUMBER OF LINES (2Ø/SCREEN) ";B
84Ø POKE 254,B
85Ø SYS CODE

```

The program's operation is quite simple, using the X register to count the bytes as they are printed across the screen using HEXPRINT (lines 12Ø to 21Ø). The second section of code (lines 22Ø to 37Ø) is responsible for printing either the ASCII character contained in the byte, or a full stop if an unprintable character or a control code is encountered. The final section of code moves the cursor down one line and increments the address counter. The whole loop is repeated until the line count reaches zero.

### Line-by-line

A line-by-line description of the Program 22 now follows:

```

line 1ØØ : print start address of current line
line 11Ø : entry for HEX-BYTES
line 12Ø : eight bytes to do (Ø-7)
line 13Ø : clear index
line 14Ø : entry for HEX-LOOP
line 15Ø : get byte through vectored address
line 16Ø : print it as two hex digits
line 17Ø : print a space
line 18Ø : increment index
line 19Ø : decrement bit count
line 2ØØ : branch to HEX-LOOP until all done
line 21Ø : print a space
line 22Ø : entry for ASCII-BYTES
line 23Ø : eight bytes to redo
line 24Ø : set index
line 25Ø : entry for ASCII-LOOP

```

line 260 : get byte through vectored address  
 line 270 : is it less than ASC“ ”?  
 line 280 : yes, branch to FULL-STOP  
 line 290 : is it greater than 128?  
 line 300 : no, branch to LEAP-FROG  
 line 310 : entry for FULL-STOP  
 line 320 : get ASC“.” into accumulator  
 line 330 : entry for LEAP-FROG  
 line 340 : print accumulator’s contents  
 line 350 : increment index  
 line 360 : decrement bit count  
 line 370 : branch to ASCII-LOOP until all done  
 line 380 : get ASCII code for RETURN  
 line 390 : print new line  
 line 400 : clear Carry flag  
 line 410 : get low byte of address  
 line 420 : add 8 to it  
 line 430 : save result  
 line 440 : if no carry, branch to NO-CARRY  
 line 450 : else increment high byte of address  
 line 460 : entry for NO-CARRY  
 line 470 : decrement line counter  
 line 480 : branch to start at \$C000 until all lines done  
 line 490 : return to BASIC  
 line 500 : entry to SPACE  
 line 510 : get ASCII code for space  
 line 520 : print it and return through jump  
 line 530 : entry to ADDRESS-PRINT  
 line 540 : load index into X register  
 line 550 : get high byte of address  
 line 560 : print it as two hex digits  
 line 570 : get low byte of address  
 line 580 : print it as two hex digits  
 line 590 : print a space  
 line 600 : print a second space  
 line 610 : return to main program  
 line 620 : entry to HEXPRINT  
 line 630 : save accumulator on stack  
 line 640 : move high nibble into low nibble position  
 line 660 : call FIRST subroutine



line 670 : restore accumulator to do low byte  
line 680 : entry for FIRST  
line 690 : mask off high nibble  
line 700 : is it less than 10?  
line 710 : yes, so jump OVER  
line 720 : add 7 to convert to A-F  
line 730 : entry to OVER  
line 740 : add 48 to convert to ASCII code  
line 750 : print it and return

# 10 Hi-res Graphics

The Commodore 64 can support hi-resolution graphics. However, as you are no doubt aware, setting up the hi-res screen prior to using it can be a rather long-winded process, requiring several lines of BASIC text. In fact, four routines are normally required:

1. Move start of BASIC user area and set position for hi-res screen.
2. Clear screen memory.
3. Select screen colour and clear to that colour.
4. Reselect normal character mode.

All of these can be performed quite simply at machine level, and the routines for each follow. They can be compiled as DATA at the end of a graphics program, poked into memory at RUN time and executed via a SYS call. This does have one of the original disadvantages, in that a large chunk of program is required. However, the main advantage is speed, particularly in clearing the screen. Alternatively, any of these routines would make an admirable addition to the Wedge Operating System, allowing it to be called by name from within your programs. Suitable command names might be:

- @MOVEBAS : move BASIC program area to make room for hi-res screen
- @HIRES : select hi-res screen
- @CLEAR : clear hi-res screen
- @GCOL : clear to graphics colour specified in a dedicated byte
- @MODE : select normal character mode

Let us now examine each command in turn.

## A BASIC MOVE

You may be wondering why we should bother to move the BASIC program area at all—why not just position the hi-res screen midway in memory? The reason for the careful positioning of the routine is as a matter of safety—placing the hi-res screen above the BASIC program area could lead to it being corrupted, especially if it is being used in conjunction with the program, because adding a line or two to the program could cause it to extend into the hi-res screen. Making sure the BASIC program fits in is no real safeguard either, as variables, strings and arrays all eat up memory at an incredible rate, and these could find their way into the screen memory. All these problems can be avoided by moving the start of BASIC up enough bytes to allow the hi-res screen to be tucked in underneath.

To do this requires a machine code program. The *Programmer's Reference Guide* lists five vectors associated with BASBAS (that's my mnemonic for BASIC's base!), as follows:

\$2B-\$2C	TXTTAB	:	start of BASIC text
\$2D-\$2E	VARTAB	:	start of BASIC variables
\$2F-\$30	ARYTAB	:	start of BASIC arrays
\$31-\$32	STREND	:	end of BASIC arrays+1
\$281-\$282	MEMSTR	:	bottom of memory

To move BASIC, each of these vectors must be reset to point to the new start area and the first three bytes of the new start area must be cleared to keep the Kernal happy.

Program 23 performs each of these functions. The address of the new BASIC area is \$4000, which allows room for the hi-res screen plus 32 sprites.

### Program 23

```
10 REM ** MOVE BASIC PROGRAM AREA START **
20 REM ** UP TO 16348 TO FREE HI-RES SCREEN **
30 :
40 CODE=49152
50 FOR LOOP=0 TO 39
60 READ BYTE
70 POKE CODE+LOOP, BYTE
80 NEXT LOOP
90 :
100 REM ** M/C DATA **
110 DATA 169,0 : REM LDA #000
```

```

12Ø DATA 141,2,64 : REM STA $4ØØ2
13Ø DATA 141,1,64 : REM STA $4ØØ1
14Ø DATA 141,Ø,64 : REM STA $4ØØØ
15Ø DATA 141,129,2 : REM STA $Ø281
16Ø DATA 169,64 : REM LDA #$4Ø
17Ø DATA 133,44 : REM STA $2C
18Ø DATA 133,46 : REM STA $2E
19Ø DATA 133,48 : REM STA $3Ø
2ØØ DATA 133,5Ø : REM STA $32
21Ø DATA 141,13Ø,2 : REM STA $Ø282
22Ø DATA 169,1 : REM LDA #$Ø1
23Ø DATA 133,43 : REM STA $2B
24Ø DATA 169,3 : REM LDA #$Ø3
25Ø DATA 133,45 : REM STA $2D
26Ø DATA 133,47 : REM STA $2F
27Ø DATA 133,49 : REM STA $31
28Ø DATA 96 : REM RTS

```

### Line-by-line

A line-by-line description of Program 23 follows:

```

line 11Ø : initialize accumulator
line 12Ø : and clear first four bytes of new program area
line 15Ø : set low byte of MEMSTR (bottom of memory pointer)
line 16Ø : load high byte of new program area address into
           accumulator
line 17Ø : set high byte of TXTTAB
line 18Ø : set high byte of VARTAB
line 19Ø : set high byte of ARYTAB
line 2ØØ : set high byte of STREND
line 21Ø : set high byte of MEMSTR
line 22Ø : load accumulator with 1
line 23Ø : store in low byte of TXTTAB
line 24Ø : load accumulator with 3
line 25Ø : set low bytes of all vectored addresses

```

## SELECTING HI-RES

Before selecting the hi-resolution screen mode, it is necessary to point the VIC chip to the start of screen memory. This is done by writing to the VIC Memory Control register located at \$D018 (57272). The actual location is controlled by the condition of bits 3, 2 and 1. Table 10.1 details their settings for various addresses.

**Table 10.1**

Bit code	Value	Address selected
xxxx000x	0	0-2047 (\$0000-\$07FF)
xxxx001x	2	2048-4095 (\$0800-\$0FFF)
xxxx010x	4	4096-6143 (\$1000-\$17FF)
xxxx011x	6	6144-8191 (\$1800-\$1FFF)
xxxx100x	8	8192-10239 (\$2000-\$27FF)
xxxx101x	10	10240-12287 (\$2800-\$2FFF)
xxxx110x	12	12288-14335 (\$3000-\$37FF)
xxxx111x	14	14336-16383 (\$3800-\$3FFF)

You can see from the table that the screen memory may be moved around in 2K block steps. An 'x' in each of the other bits denotes that these bits may be in either state. However, remember that these bits are controlling other aspects of the VIC's function, so that any reprogramming of bits 3, 2 and 1 must preserve the other bits. This is best done with the logical OR function. Looking at Table 10.1 we can see that bit 3 must be set to point the Memory Control register at location 8192. In BASIC this would simplify to:

```
100 A=PEEK(53727) : REM GET VALUE
110 A=A OR 8 : REM SET BIT 3
120 POKE 53727,A : REM REPROGRAM
```

which translates to assembler as:

```
LDA #08
ORA $D018
STA $D018
```

Now that the hi-res screen has been defined, it can be switched in by setting bit 5 of the VIC Control register at \$D011 (53265).

Again, the other bits in the register must be preserved, so the byte must be ORed with 32 (00100000 binary). In BASIC this is:

```
130 A=PEEK(53265) : REM GET VALUE
140 A=A OR 32      : REM SET BIT 5
150 POKE 53265,A  : REM REPROGRAM
```

and in assembler:

```
LDA #$20
ORA $D011
STA $D011
```

## A CLEAR VIEW

Once hi-res mode has been selected, it will be filled with junk (often referred to as garbage). To clear this, each location must in turn be POKEd with zero. A BASIC program to do this would take the form:

```
200 SB=8192
210 FOR L=SB TO SB+7999
220 POKE L,0
230 NEXT L
```

Previously, in normal character mode, locations 1024 to 2023 were used to control which character was displayed—for example, POKing a 1 into location 1024 would make a letter A appear at the top left hand corner of the screen. When in hi-res mode, this area of memory is used to hold the colour information of that byte. Note that the colour information does not now come from the colour memory—colour details are taken directly from the hi-res screen itself. The high nibble of the byte (that is, bits 4 to 7) holds the colour code of any bit that is set in that 8 by 8 bit matrix, while the lower nibble (bits 3 to 0) holds the colour of any bits that are clear in the same area.

To clear the hi-res screen to black ink on green paper in BASIC we would use:

```
240 FOR C=1024 TO 2023
250 POKE C,13
260 NEXT C
```

If all the above BASIC program lines were to be combined and RUN, the resulting hi-res screen would take around 20 seconds to construct—a bit slow, you'll agree! Program 24 provides the

machine code equivalent. Note that the value assigned to CODE is 49408 and NOT 49152 as we have been using previously. This is to allow the program to be used in conjunction with the MOVEBAS program described earlier. After you have entered and RUN MOVEBAS, try this one for an instant hi-res screen!

### Program 24

```

10 REM ** HI-RES GRAPHICS SCREEN SET AND
    CLEAR **
20 CODE=49408
30 FOR LOOP=0 TO 105
40 READ BYTE
50 POKE CODE+LOOP,BYTE
60 NEXT LOOP
70 :
80 REM ** M/C DATA **
85 :: REM SELECT-HI-RES
90 DATA 169,8 : REM LDA #00
100 DATA 13,24,208 : REM ORA $D018
110 DATA 141,24,208 : REM STA $D018
120 DATA 169,32 : REM LDA #$20
130 DATA 13,17,208 : REM ORA $D011
140 DATA 141,17,208 : REM STA $D011
150 :: REM CLEAR-SCREEN-MEMORY
160 DATA 169,0 : REM LDA #00
170 DATA 133,251 : REM STA $FC
180 DATA 169,32 : REM LDA #$20
190 DATA 133,252 : REM STA $FC
200 DATA 169,64 : REM LDA #$40
210 DATA 133,253 : REM STA $FD
220 DATA 169,63 : REM LDA #$3F
230 DATA 133,254 : REM STA $FE
240 :: REM IN
250 DATA 165,252 : REM LDA $FC
260 DATA 197,254 : REM CMP $FE
270 DATA 208,9 : REM BNE $09
280 DATA 165,251 : REM LDA $FB
290 DATA 197,253 : REM CMP $FD
300 DATA 208,3 : REM BNE $03

```

```

31Ø DATA 76,62,192 : REM JMP $CØ3E
32Ø :: REM CLEAR
33Ø DATA 16Ø,Ø : REM LDY #ØØØ
34Ø DATA 169,Ø : REM LDA #ØØØ
35Ø DATA 145,251 : REM STA ($FB),Y
36Ø DATA 23Ø,251 : REM INC $FB
37Ø DATA 2Ø8,231 : REM BNE $E7
38Ø DATA 23Ø,252 : REM INC $FC
39Ø DATA 56 : REM SEC
40Ø DATA 176,226 : REM BCS $E2
41Ø :
42Ø :: REM COLOUR
43Ø DATA 169,Ø : REM LDA #ØØØ
44Ø DATA 133,251 : REM STA $FB
45Ø DATA 169,4 : REM LDA #ØØ4
46Ø DATA 133,252 : REM STA $FC
47Ø DATA 169,231 : REM LDA #$E7
48Ø DATA 133,253 : REM STA $FD
49Ø DATA 169,7 : REM LDA #ØØ7
50Ø DATA 133,254 : REM STA $FE
51Ø :: REM CIN
52Ø DATA 165,252 : REM LDA $FC
53Ø DATA 197,254 : REM CMP $FE
54Ø DATA 2Ø8,7 : REM BNE $Ø7
55Ø DATA 165,251 : REM LDA $FB
56Ø DATA 197,253 : REM CMP $FD
57Ø DATA 2Ø8,1 : REM BNE $Ø1
58Ø DATA 96 : REM RTS
59Ø :: REM GREEN
60Ø DATA 16Ø,Ø : REM LDY #ØØØ
61Ø DATA 169,13 : REM LDA #ØØD
62Ø DATA 145,251 : REM STA ($FB),Y
63Ø DATA 23Ø,251 : REM INC $FB
64Ø DATA 2Ø8,233 : REM BNE $E9
65Ø DATA 23Ø,252 : REM INC $FC
66Ø DATA 56 : REM SEC
67Ø DATA 176,228 : REM BCS $E4

```



## Line-by-line

A line-by-line description of Program 24 follows:

line 90 : load accumulator with mask 00001000  
line 100 : force bit 3 to select 8196 as bit map start address  
line 110 : and program VIC Memory Control register  
line 120 : load accumulator with mask 00100000  
line 130 : force bit 5 to select bit map mode  
line 140 : and program CIC Control register  
line 150 : entry for bit map CLEAR-SCREEN-MEMORY  
routine  
line 160 : set up vector to point to screen start address \$2000  
line 200 : set up vector to point to screen end address \$403F  
line 240 : entry for IN  
line 250 : get high byte current address  
line 260 : is it same as high byte end address?  
line 270 : no, so branch to CLEAR  
line 280 : yes, get low byte current address  
line 290 : is it same as low byte end address  
line 300 : no, so branch to CLEAR  
line 310 : yes, all done jump to COLOUR  
line 320 : entry for CLEAR  
line 330 : initialize index  
line 340 : clear accumulator  
line 350 : clear byte of screen memory  
line 360 : increment low byte of current screen address  
line 370 : branch to IN if no carry over  
line 380 : increment high byte  
line 390 : set Carry flag  
line 400 : force branch to IN  
line 420 : entry for COLOUR  
line 430 : set up vector to point to start of colour memory  
line 470 : set up vector to point to end of colour memory  
line 510 : entry for CIN  
line 520 : get high byte of current address  
line 530 : is it the same as high byte end address?  
line 540 : no, branch to GREEN

line 55Ø : get low byte of current address  
line 56Ø : is it the same as the low byte end address?  
line 57Ø : no, branch to GREEN  
line 58Ø : back to calling routine  
line 59Ø : entry for GREEN  
line 60Ø : clear indexing register  
line 61Ø : get code for green into accumulator  
line 62Ø : POKE it into colour memory  
line 63Ø : increment low byte of current address  
line 64Ø : branch to CIN if no carry over  
line 65Ø : increment high byte  
line 66Ø : set Carry flag  
line 67Ø : and force branch to CIN

# Appendix 1: 6510 Complete Instruction Set

<b>ADC</b>		<b>Add with carry</b>		<b>NZCV</b>
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>	
Immediate	\$69	2	2	
Zero page	\$65	2	3	
Zero page,X	\$75	2	4	
Absolute	\$6D	3	4	
Absolute,X	\$7D	3	4 or 5	
Absolute,Y	\$79	3	4 or 5	
(Indirect,X)	\$61	2	6	
(Indirect),Y	\$71	2	5	

<b>AND</b>		<b>AND with accumulator</b>		<b>NZ</b>
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>	
Immediate	\$29	2	2	
Zero page	\$25	2	3	
Zero page,X	\$35	2	4	
Absolute	\$2D	3	4	
Absolute,X	\$3D	3	4 or 5	
Absolute,Y	\$39	3	4 or 5	
(Indirect,X)	\$21	2	6	
(Indirect),Y	\$31	2	5	

<b>ASL</b>	Shift left		<b>NZC</b>
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Accumulator	\$0A	1	2
Zero page	\$06	2	5
Zero page,X	\$16	2	6
Absolute	\$0E	3	6
Absolute,X	\$1E	3	7

<b>BCC</b>	Branch if C = 0		Flags unaltered
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Relative	\$90	2	3 or 2

<b>BCS</b>	Branch if C = 1		Flags unaltered
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Relative	\$B0	2	3 or 2

<b>BEQ</b>	Branch if Z = 1		Flags unaltered
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Relative	\$F0	2	3 or 2

---

<b>BIT</b>			Z,N,V
------------	--	--	-------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Zero page	\$24	2	3
Absolute	\$2C	3	4

---



---

<b>BMI</b>	Branch if N = 1		Flags unaltered
------------	-----------------	--	-----------------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Relative	\$30	2	3 or 2

---



---

<b>BNE</b>	Branch if Z = 0		Flags unaltered
------------	-----------------	--	-----------------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Relative	\$D0	2	3 or 2

---



---

<b>BPL</b>	Branch if N = 0		Flags unaltered
------------	-----------------	--	-----------------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Relative	\$10	2	3 or 2

---

<b>BRK</b>	Break			B flag = 1
<i>Address mode</i>		<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied		\$00	1	7
<b>BVC</b>	Branch if V = 0			Flags unaltered
<i>Address mode</i>		<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Relative		\$50	2	3 or 2
<b>BVS</b>	Branch if V = 1			Flags unaltered
<i>Address mode</i>		<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Relative		\$70	2	3 or 2
<b>CLC</b>	Clear Carry flag			C flag = 0
<i>Address mode</i>		<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied		\$18	1	2

<b>CLD</b>	Clear Decimal flag	D flag = 0	
<i>Address mode</i> Implied	<i>Op-code</i> \$D8	<i>Bytes</i> 1	<i>Cycles</i> 2
<b>CLI</b>	Clear Interrupt flag	I flag = 0	
<i>Address mode</i> Implied	<i>Op-code</i> \$58	<i>Bytes</i> 1	<i>Cycles</i> 2
<b>CLV</b>	Clear Overflow flag	V flag = 0	
<i>Address mode</i> Implied	<i>Op-code</i> \$B8	<i>Bytes</i> 1	<i>Cycles</i> 2
<b>CMP</b>	Compare accumulator	NZC	
<i>Address mode</i> Immediate	<i>Op-code</i> \$C9	<i>Bytes</i> 2	<i>Cycles</i> 2
Zero page	\$C5	2	3
Zero page,X	\$D5	2	4
Absolute	\$CD	3	4
Absolute,X	\$DD	3	4 or 5
Absolute,Y	\$D9	3	4 or 5
(Indirect,X)	\$C1	2	6
(Indirect),Y	\$D1	2	5 or 6

---

<b>CPX</b>	Compare X register		<b>NZC</b>
------------	--------------------	--	------------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$E0	2	2
Zero page	\$E4	2	3
Absolute	\$EC	3	4

---



---

<b>CPY</b>	Compare Y register		<b>NZC</b>
------------	--------------------	--	------------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$C0	2	2
Zero page	\$C4	2	3
Absolute	\$CC	3	4

---



---

<b>DEC</b>	Decrement memory		<b>NZ</b>
------------	------------------	--	-----------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Zero page	\$C6	2	5
Zero page,X	\$D6	2	6
Absolute	\$CE	3	6
Absolute,X	\$DE	3	7

---



---

<b>DEX</b>	Decrement X register		<b>NZ</b>
------------	----------------------	--	-----------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$CA	1	2

---



---

<b>DEY</b>	Decrement Y register		NZ
------------	----------------------	--	----

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$88	1	2

---



---

<b>EOR</b>	Exclusive-OR		NZ
------------	--------------	--	----

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$ 49	2	2
Zero page	\$ 45	2	3
Zero page,X	\$ 55	2	4
Absolute	\$ 4D	3	4
Absolute,X	\$ 5D	3	4 or 5
Absolute,Y	\$ 59	3	4 or 5
(Indirect,X)	\$ 41	2	6
(Indirect),Y	\$ 51	2	5

---



---

<b>INC</b>	Increment memory		NZ
------------	------------------	--	----

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Zero page	\$E6	2	5
Zero page,X	\$F6	2	6
Absolute	\$EE	3	6
Absolute,X	\$FE	3	7

---

---

<b>INX</b>	Increment X register		NZ
------------	----------------------	--	----

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$E8	1	2

---



---

<b>INY</b>	Increment Y register		NZ
------------	----------------------	--	----

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$C8	1	2

---



---

<b>JMP</b>	Jump		Flags unaltered
------------	------	--	-----------------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Absolute	\$4C	3	3
Indirect	\$6C	3	5

---



---

<b>JSR</b>	Jump to subroutine		Flags unaltered
------------	--------------------	--	-----------------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Absolute	\$20	3	6

---

---

<b>LDA</b>	Load accumulator		NZ
------------	------------------	--	----

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$A9	2	2
Zero page	\$A5	2	3
Zero page,X	\$B5	2	4
Absolute	\$AD	3	4
Absolute,X	\$BD	3	4 or 5
Absolute,Y	\$B9	3	4 or 5
(Indirect,X)	\$A1	2	6
(Indirect),Y	\$B1	2	5 or 6

---



---

<b>LDX</b>	Load X register		NZ
------------	-----------------	--	----

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$A2	2	2
Zero page	\$A6	2	3
Zero page,Y	\$B6	2	4
Absolute	\$AE	3	4
Absolute,Y	\$BE	3	4 or 5

---



---

<b>LDY</b>	Load Y register		NZ
------------	-----------------	--	----

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$A0	2	2
Zero page	\$A4	2	3
Zero page,X	\$B4	2	4
Absolute	\$AC	3	4
Absolute,X	\$BC	3	4 or 5

---

---

<b>LSR</b>	Logical shift right	<b>N = 0,ZC</b>	
------------	---------------------	-----------------	--

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Accumulator	\$4A	1	2
Zero page	\$46	2	5
Zero page,X	\$56	2	6
Absolute	\$4E	3	6
Absolute,X	\$5E	3	7

---



---

<b>NOP</b>	No operation	<b>Flags unaltered</b>	
------------	--------------	------------------------	--

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$EA	1	2

---



---

<b>ORA</b>	Inclusive OR	<b>NZ</b>	
------------	--------------	-----------	--

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$09	2	2
Zero page	\$05	2	3
Zero page,X	\$15	2	4
Absolute	\$0D	3	4
Absolute,X	\$1D	3	4 or 5
Absolute,Y	\$19	3	4 or 5
(Indirect,X)	\$01	2	6
(Indirect),Y	\$11	2	5

---

---

<b>PHA</b>	Push accumulator	Flags unaltered	
------------	------------------	-----------------	--

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$48	1	3

---



---

<b>PHP</b>	Push Status register	Flags unaltered	
------------	----------------------	-----------------	--

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$08	1	3

---



---

<b>PLA</b>	Pull accumulator	NZ	
------------	------------------	----	--

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$68	1	4

---



---

<b>PLP</b>	Pull Status register	Flags as status	
------------	----------------------	-----------------	--

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$28	1	4

---

<b>ROL</b>	<b>Rotate left</b>	<b>NZC</b>	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Accumulator	\$2A	1	2
Zero page	\$26	2	5
Zero page,X	\$36	2	6
Absolute	\$2E	3	6
Absolute,X	\$3E	3	7

<b>ROR</b>	<b>Rotate right</b>	<b>NZC</b>	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Accumulator	\$6A	1	2
Zero page	\$66	2	5
Zero page,X	\$76	2	6
Absolute	\$6E	3	6
Absolute,X	\$7E	3	7

<b>RTI</b>	<b>Return from interrupt</b>	<b>Flags as pulled</b>	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$40	1	6

<b>RTS</b>	<b>Return from subroutine</b>	<b>Flags unaltered</b>	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$ 60	1	6

<b>SBC</b>	<b>Subtract from accumulator</b>	<b>NZCV</b>	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Immediate	\$ E9	2	2
Zero page	\$ E5	2	3
Zero page,X	\$ F5	2	4
Absolute	\$ ED	3	4
Absolute,X	\$ FD	3	4 or 5
Absolute,Y	\$ F9	3	4 or 5
(Indirect,X)	\$ E1	2	6
(Indirect),Y	\$ F1	2	5 or 6

<b>SEC</b>	<b>Set Carry flag</b>	<b>C = 1</b>	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$ 38	1	2

<b>SED</b>	<b>Set Decimal flag</b>	<b>D = 1</b>	
<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$ F8	1	2

---

<b>SEI</b>	Set Interrupt flag		I = 1
------------	--------------------	--	-------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$ 78	1	2

---



---

<b>STA</b>	Store accumulator		Flags unaltered
------------	-------------------	--	-----------------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Zero page	\$ 85	2	3
Zero page,X	\$ 95	2	4
Absolute	\$ 8D	3	4
Absolute,X	\$ 9D	3	5
Absolute,Y	\$ 99	3	5
(Indirect,X)	\$ 81	2	6
(Indirect),Y	\$ 91	2	6

---



---

<b>STX</b>	Store X register		Flags unaltered
------------	------------------	--	-----------------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Zero page	\$ 86	2	3
Zero page,Y	\$ 96	2	4
Absolute	\$ 8E	3	4

---

.



---

<b>STY</b>	Store Y register	<b>Flags unaltered</b>	
------------	------------------	------------------------	--

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Zero page	\$84	2	3
Zero page,X	\$94	2	4
Absolute	\$8C	3	4

---



---

<b>TAX</b>	Transfer accumulator to X	<b>NZ</b>	
------------	---------------------------	-----------	--

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$AA	1	2

---



---

<b>TAY</b>	Transfer accumulator to Y	<b>NZ</b>	
------------	---------------------------	-----------	--

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$A8	1	2

---



---

<b>TSX</b>	Transfer Stack Pointer to X	<b>NZ</b>	
------------	-----------------------------	-----------	--

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
Implied	\$BA	1	2

---

---

<b>TXA</b>	Transfer X to accumulator		<b>NZ</b>
------------	---------------------------	--	-----------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
<b>Implied</b>	<b>\$8A</b>	<b>1</b>	<b>2</b>

---



---

<b>TXS</b>	Transfer X to Stack Pointer		<b>Flags unaltered</b>
------------	-----------------------------	--	------------------------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
<b>Implied</b>	<b>\$9A</b>	<b>1</b>	<b>2</b>

---



---

<b>TYA</b>	Transfer Y to accumulator		<b>NZ</b>
------------	---------------------------	--	-----------

---

<i>Address mode</i>	<i>Op-code</i>	<i>Bytes</i>	<i>Cycles</i>
<b>Implied</b>	<b>\$98</b>	<b>1</b>	<b>2</b>

---

# Appendix 2: 6510 Opcodes

All numbers are hexadecimal.

00	BRK implied	1C	Future expansion
01	ORA (zero page, X)	1D	ORA absolute, X
02	Future expansion	1E	ASL absolute, X
03	Future expansion	1F	Future expansion
04	Future expansion	20	JSR absolute
05	ORA zero page	21	AND (zero page, X)
06	ASL zero page	22	Future expansion
07	Future expansion	23	Future expansion
08	PHP implied	24	BIT zero page
09	ORA #immediate	25	AND zero page
0A	ASL accumulator	26	ROL zero page
0B	Future expansion	27	Future expansion
0C	Future expansion	28	PLP implied
0D	ORA absolute	29	AND #immediate
0E	ASL absolute	2A	ROL accumulator
0F	Future expansion	2B	Future expansion
10	BPL relative	2C	BIT absolute
11	ORA (zero page), Y	2D	AND absolute
12	Future expansion	2E	ROL absolute
13	Future expansion	2F	Future expansion
14	Future expansion	30	BMI relative
15	ORA zero page, X	31	AND (zero page), Y
16	ASL zero page, X	32	Future expansion
17	Future expansion	33	Future expansion
18	CLC implied	34	Future expansion
19	ORA absolute, Y	35	AND zero page, X
1A	Future expansion	36	ROL zero page, X
1B	Future expansion	37	Future expansion

38 SEC implied  
39 AND absolute, Y  
3A Future expansion  
3B Future expansion  
3C Future expansion  
3D AND absolute, X  
3E ROL absolute, X  
3F Future expansion  
40 RTI implied  
41 EOR (zero page, X)  
42 Future expansion  
43 Future expansion  
44 Future expansion  
45 EOR zero page  
46 LSR zero page  
47 Future expansion  
48 PHA implied  
49 EOR #immediate  
4A LSR accumulator  
4B Future expansion  
4C JMP absolute  
4D EOR absolute  
4E LSR absolute  
4F Future expansion  
50 BVC relative  
51 EOR (zero page), Y  
52 Future expansion  
53 Future expansion  
54 Future expansion  
55 EOR zero page, X  
56 LSR zero page, X  
57 Future expansion  
58 CLI implied  
59 EOR absolute, Y  
5A Future expansion  
5B Future expansion  
5C Future expansion

5D EOR absolute, X  
5E LSR absolute, X  
5F Future expansion  
60 RTS implied  
61 ADC (zero page, X)  
62 Future expansion  
63 Future expansion  
64 Future expansion  
65 ADC zero page  
66 ROR zero page  
67 Future expansion  
68 PLA implied  
69 ADC #immediate  
6A ROR accumulator  
6B Future expansion  
6C JMP (indirect)  
6D ADC absolute  
6E ROR absolute  
6F Future expansion  
70 BVS relative  
71 ADC (zero page), Y  
72 Future expansion  
73 Future expansion  
74 Future expansion  
75 ADC zero page, X  
76 ROR zero page, X  
77 Future expansion  
78 SEI implied  
79 ADC absolute, Y  
7A Future expansion  
7B Future expansion  
7C Future expansion  
7D ADC absolute, X  
7E ROR absolute, X  
7F Future expansion  
80 Future expansion  
81 STA (zero page, X)

82 Future expansion  
83 Future expansion  
84 STY zero page  
85 STA zero page  
86 STX zero page  
87 Future expansion  
88 DEY implied  
89 Future expansion  
8A TXA implied  
8B Future expansion  
8C STY absolute  
8D STA absolute  
8E STX absolute  
8F Future expansion  
90 BCC relative  
91 STA (zero page), Y  
92 Future expansion  
93 Future expansion  
94 STY zero page, X  
95 STA zero page, X  
96 STX zero page, Y  
97 Future expansion  
98 TYA implied  
99 STA absolute, Y  
9A TXS implied  
9B Future expansion  
9C Future expansion  
9D STA absolute, X  
9E Future expansion  
9F Future expansion  
A0 LDY #immediate  
A1 LDA (zero page, X)  
A2 LDX #immediate  
A3 Future expansion  
A4 LDY zero page  
A5 LDA zero page  
A6 LDX zero page

A7 Future expansion  
A8 TAY implied  
A9 LDA #immediate  
AA TAX implied  
AB Future expansion  
AC LDY absolute  
AD LDA absolute  
AE LDX absolute  
AF Future expansion  
B0 BCS relative  
B1 LDA (zero page), Y  
B2 Future expansion  
B3 Future expansion  
B4 LDY zero page, X  
B5 LDA zero page, X  
B6 LDX zero page, Y  
B7 Future expansion  
B8 CLV implied  
B9 LDA absolute, Y  
BA TSX implied  
BB Future expansion  
BC LDY absolute, X  
BD LDA absolute, X  
BE LDX absolute, Y  
BF Future expansion  
C0 CPY #immediate  
C1 CMP (zero page, X)  
C2 Future expansion  
C3 Future expansion  
C4 CPY zero page  
C5 CMP zero page  
C6 DEC zero page  
C7 Future expansion  
C8 INY implied  
C9 CMP #immediate  
CA DEX implied  
CB Future expansion

CC	CPY absolute	E6	INC zero page
CD	CMP absolute	E7	Future expansion
CE	DEC absolute	E8	INX implied
CF	Future expansion	E9	SBC #immediate
D0	BNE relative	EA	NOP implied
D1	CMP (zero page), Y	EB	Future expansion
D2	Future expansion	EC	CPX absolute
D3	Future expansion	ED	SBC absolute
D4	Future expansion	EE	INC absolute
D5	CMP zero page, X	EF	Future expansion
D6	DEC zero page, X	F0	BEQ relative
D7	Future expansion	F1	SBC (zero page), Y
D8	CLD implied	F2	Future expansion
D9	CMP absolute, Y	F3	Future expansion
DA	Future expansion	F4	Future expansion
DB	Future expansion	F5	SBC zero page, X
DC	Future expansion	F6	INC zero page, X
DD	CMP absolute, X	F7	Future expansion
DE	DEC absolute, X	F8	SED implied
DF	Future expansion	F9	SBC absolute, Y
E0	CPX #immediate	FA	Future expansion
E1	SBC (zero page, X)	FB	Future expansion
E2	Future expansion	FC	Future expansion
E3	Future expansion	FD	SBC absolute, X
E4	CPX zero page	FE	INC absolute, X
E5	SBC zero page	FF	Future expansion

# Appendix 3: Commodore 64 Memory Map

Kernal Operating System ROM	FFFF
Colour RAM	DC00
VIC and SID	D800
'Free' RAM	D000
BASIC interpreter ROM	C000
VSP cartridge ROM	A000
Program area	8000
Screen memory	800
Kernal vectors and flags	400
Input buffers	300
Stack	200
Zero page	100
	00

# Appendix 4: Branch Calculators

The branch calculators are used to give branch values in hex. First, count the number of bytes you need to branch. Then locate this number in the centre of the appropriate table, and finally, read off the high and low hex nibbles from the side column and top row respectively.

*Example* For a backward branch of 16 bytes:

Locate 16 in the centre of Table A4.1 (bottom row), then read off high nibble (#F) and low nibble (#0) to give displacement value (#F0).

**Table A4.1 Backward branch calculator**

LSD \ MSD	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
8	128	127	126	125	124	123	122	121	120	119	118	117	116	115	114	113
9	112	111	110	109	108	107	106	105	104	103	102	101	100	99	98	97
A	96	95	94	93	92	91	90	89	88	87	86	85	84	83	82	81
B	80	79	78	77	76	75	74	73	72	71	70	69	68	67	66	65
C	64	63	62	61	60	59	58	57	56	55	54	53	52	51	50	49
D	48	47	46	45	44	43	42	41	40	39	38	37	36	35	34	33
E	32	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17
F	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1

**Table A4.2 Forward branch calculator**

LSD \ MSD	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
1	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
2	32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	47
3	48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63
4	64	65	66	67	68	69	70	71	72	73	74	75	76	77	78	79
5	80	81	82	83	84	85	86	87	88	89	90	91	92	93	94	95
6	96	97	98	99	100	101	102	103	104	105	106	107	108	109	110	111
7	112	113	114	115	116	117	118	119	120	121	122	123	124	125	126	127



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*Other titles of interest*

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